

low level

TEX

lowlevel

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## Colofon

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## **2 Conditionals**



low level

TEX

conditionals

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## 2.1 Preamble

### 2.1.1 Introduction

You seldom need the low level conditionals because there are quite some so called support macros available in ConT<sub>E</sub>Xt. For instance, when you want to compare two values (or more accurate: sequences of tokens), you can do this:

```
\doifelse {foo} {bar} {
  the same
} {
  different
}
```

But if you look in the ConT<sub>E</sub>Xt code, you will see that often we use primitives that start with `\if` in low level macros. There are good reasons for this. First of all, it looks familiar when you also code in other languages. Another reason is performance but that is only true in cases where the snippet of code is expanded very often, because T<sub>E</sub>X is already pretty fast. Using low level T<sub>E</sub>X can also be more verbose, which is not always nice in a document source. But, the most important reason (for me) is the layout of the code. I often let the look and feel of code determine the kind of coding. This also relates to the syntax highlighting that I am using, which is consistent for T<sub>E</sub>X, MetaPost, Lua, etc. and evolved over decades. If code looks bad, it probably is bad. Of course this doesn't mean all my code looks good; you're warned. In general we can say that I often use `\if...` when coding core macros, and `\doifelse...` macros in (document) styles and modules.

In the sections below I will discuss the low level conditions in T<sub>E</sub>X. For the often more convenient ConT<sub>E</sub>Xt wrappers you can consult the source of the system and support modules, the wiki and/or manuals.

Some of the primitives shown here are only available in Lua $\TeX$ , and some only in Lua-Meta $\TeX$ . We could do without them for decades but they were added to these engines because of convenience and, more important, because then made for nicer code. Of course there's also the fun aspect. This manual is not an invitation to use these very low level primitives in your document source. The ones that probably make most sense are `\ifnum`, `\ifdim` and `\ifcase`. The others are often wrapped into support macros that are more convenient.

## 2.1.2 Number and dimensions

Numbers and dimensions are basic data types in  $\TeX$ . When you enter one, a number is just that but a dimension gets a unit. Compare:

```
1234
1234pt
```

If you also use MetaPost, you need to be aware of the fact that in that language there are not really dimensions. The `post` part of the name implies that eventually a number becomes a PostScript unit which represents a base point (`bp`) in  $\TeX$ . When in MetaPost you entry `1234pt` you actually multiply 1234 by the variable `pt`. In  $\TeX$  on the other hand, a unit like `pt` is one of the keywords that gets parsed. Internally dimensions are also numbers and the unit (keyword) tells the scanner what multiplier to use. When that multiplier is one, we're talking of scaled points, with the unit `sp`.

```
\the\dimexpr 12.34pt \relax
\the\dimexpr 12.34sp \relax
\the\dimexpr 12.99sp \relax
\the\dimexpr 1234sp \relax
\the\numexpr 1234 \relax
```

```
12.34pt
0.00018pt
0.00018pt
0.01883pt
1234
```

When we serialize a dimension it always shows the dimension in points, unless we serialize it as number.

```
\scratchdimen1234sp
\number\scratchdimen
\the\scratchdimen
```

1234  
0.01883pt

When a number is scanned, the first thing that is taken care of is the sign. In many cases, when  $\text{T}_{\text{E}}\text{X}$  scans for something specific it will ignore spaces. It will happily accept multiple signs:

```
\number +123
\number +++123
\number + + + 123
\number +--+123
\number --123
\number ---123
```

123  
123  
123  
123  
123  
-123

Watch how the negation accumulates. The scanner can handle decimal, hexadecimal and octal numbers:

```
\number -123
\number -"123
\number -'123
```

-123  
-291  
-83

A dimension is scanned like a number but this time the scanner checks for upto three parts: an either or not signed number, a period and a fraction. Here no number means zero, so the next is valid:

```
\the\dimexpr . pt \relax
\the\dimexpr 1. pt \relax
\the\dimexpr .1pt \relax
\the\dimexpr 1.1pt \relax
```

0.0pt  
1.0pt

0.1pt

1.1pt

Again we can use hexadecimal and octal numbers but when these are entered, there can be no fractional part.

```
\the\dimexpr 16 pt \relax
```

```
\the\dimexpr "10 pt \relax
```

```
\the\dimexpr '20 pt \relax
```

16.0pt

16.0pt

16.0pt

The reason for discussing numbers and dimensions here is that there are cases where when  $\TeX$  expects a number it will also accept a dimension. It is good to know that for instance a macro defined with `\chardef` or `\mathchardef` also is treated as a number. Even normal characters can be numbers, when prefixed by a ``` (backtick).

The maximum number in  $\TeX$  is 2147483647 so we can do this:

```
\scratchcounter2147483647
```

but not this

```
\scratchcounter2147483648
```

as it will trigger an error. A dimension can be positive and negative so there we can do at most:

```
\scratchdimen 1073741823sp
```

```
\scratchdimen1073741823sp
```

```
\number\scratchdimen
```

```
\the\scratchdimen
```

```
\scratchdimen16383.99998pt
```

```
\number\scratchdimen
```

```
\the\scratchdimen
```

1073741823

16383.99998pt

1073741823

16383.99998pt

We can also do this:

```
\scratchdimen16383.99999pt
\number\scratchdimen
\the\scratchdimen

1073741823
16383.99998pt
```

but the next one will fail:

```
\scratchdimen16383.9999999pt
```

Just keep in mind that T<sub>E</sub>X scans both parts as number so the error comes from checking if those numbers combine well.

```
\ifdim 16383.99999 pt = 16383.99998 pt the same \else different \fi
\ifdim 16383.999979 pt = 16383.999980 pt the same \else different \fi
\ifdim 16383.999987 pt = 16383.999991 pt the same \else different \fi
```

Watch the difference in dividing, the / rounds, while the : truncates.

```
the same
the same
the same
```

You need to be aware of border cases, although in practice they never really are a problem:

```
\ifdim \dimexpr16383.99997 pt/2\relax = \dimexpr 16383.99998 pt/2\relax
the same \else different
\fi
\ifdim \dimexpr16383.99997 pt:2\relax = \dimexpr 16383.99998 pt:2\relax
the same \else different
\fi
```

```
different
the same
```

```
\ifdim \dimexpr1.99997 pt/2\relax = \dimexpr 1.99998 pt/2\relax
the same \else different
\fi
\ifdim \dimexpr1.99997 pt:2\relax = \dimexpr 1.99998 pt:2\relax
the same \else different
```

**\fi**different  
the same**\ifdim \dimexpr**1.999999 pt/2**\relax = \dimexpr** 1.999995 pt/2**\relax**  
the same **\else** different**\fi****\ifdim \dimexpr**1.999999 pt:2**\relax = \dimexpr** 1.999995 pt:2**\relax**  
the same **\else** different**\fi**the same  
the same

This last case demonstrates that at some point the digits get dropped (still assuming that the fraction is within the maximum permitted) so these numbers then are the same. Anyway, this is not different in other programming languages and just something you need to be aware of.

## 2.2 T<sub>E</sub>X primitives

### 2.2.1 \if

I seldom use this one. Internally T<sub>E</sub>X stores (and thinks) in terms of tokens. If you see for instance `\def` or `\dimen` or `\hbox` these all become tokens. But characters like A or @ also become tokens. In this test primitive all non-characters are considered to be the same. In the next examples this is demonstrated.

```
[\if AB yes\else nop\fi]
[\if AA yes\else nop\fi]
[\if CDyes\else nop\fi]
[\if CCyes\else nop\fi]
[\if\dimen\font yes\else nop\fi]
[\if\dimen\font yes\else nop\fi]
```

Watch how spaces after the two characters are kept: [nop] [ yes] [nop] [yes] [yes] [yes]. This primitive looks at the next two tokens but when doing so it expands. Just look at the following:

```
\def\AA{AA}%
```

```
\def\AB{AB}%
[\if\AA yes\else nop\fi]
[\if\AB yes\else nop\fi]
```

We get: [yes] [nop].

### 2.2.2 \ifcat

In T<sub>E</sub>X characters (in the input) get interpreted according to their so called catcodes. The most common are letters (alphabetic) and other (symbols) but for instance the backslash has the property that it starts a command, the dollar signs trigger math mode, while the curly braced deal with grouping. If for instance either or not the ampersand is special (for instance as column separator in tables) depends on the macro package.

```
[\ifcat AB yes\else nop\fi]
[\ifcat AA yes\else nop\fi]
[\ifcat CDyes\else nop\fi]
[\ifcat CCyes\else nop\fi]
[\ifcat Clyes\else nop\fi]
[\ifcat\dimen\font yes\else nop\fi]
[\ifcat\dimen\font yes\else nop\fi]
```

This time we also compare a letter with a number: [ yes] [ yes] [yes] [yes] [nop] [yes] [yes]. In that case the category codes differ (letter vs other) but in this test comparing the letters result in a match. This is a test that is used only once in ConT<sub>E</sub>Xt and even that occasion is dubious and will go away.

You can use \noexpand to prevent expansion:

```
\def\A{A}%
\let\B B%
\def\C{D}%
\let\D D%
[\ifcat\noexpand\A Ayes\else nop\fi]
[\ifcat\noexpand\B Byes\else nop\fi]
[\ifcat\noexpand\C Cyes\else nop\fi]
[\ifcat\noexpand\C Dyes\else nop\fi]
[\ifcat\noexpand\D Dyes\else nop\fi]
```

We get: [nop] [yes] [nop] [nop] [yes], so who still thinks that T<sub>E</sub>X is easy to understand for a novice user?



### 2.2.3 `\ifnum`

This condition compares its argument with another one, separated by an `<`, `=` or `>` character.

```
\ifnum\scratchcounter<0
  less than
\else\ifnum\scratchcounter>0
  more than
\else
  equal to
\fi zero
```

This is one of these situations where a dimension can be used instead. In that case the dimension is in scaled points.

```
\ifnum\scratchdimen<0
  less than
\else\ifnum\scratchdimen>0
  more than
\else
  equal to
\fi zero
```

Of course this equal treatment of a dimension and number is only true when the dimension is a register or box property.

### 2.2.4 `\ifdim`

This condition compares one dimension with another one, separated by an `<`, `=` or `>` sign.

```
\ifdim\scratchdimen<0pt
  less than
\else\ifdim\scratchdimen>0pt
  more than
\else
  equal to
\fi zero
```

While when comparing numbers a dimension is a valid quantity but here you cannot mix them: something with a unit is expected.

### 2.2.5 `\ifodd`

This one can come in handy, although in ConT<sub>E</sub>Xt it is only used in checking for an odd of even page number.

```
\scratchdimen 3sp
\scratchcounter4

\ifodd\scratchdimen very \else not so \fi odd
\ifodd\scratchcounter very \else not so \fi odd
```

As with the previously discussed `\ifnum` you can use a dimension variable too, which is then interpreted as representing scaled points. Here we get:

```
very odd
not so odd
```

### 2.2.6 `\ifvmode`

This is a rather trivial check. It takes no arguments and just is true when we're in vertical mode. Here is an example:

```
\hbox{\ifvmode\else\par\fi\ifvmode v\else h\fi mode}
```

We're always in horizontal mode and issuing a `\par` inside a horizontal box doesn't change that, so we get: `hmode`.

### 2.2.7 `\ifhmode`

As with `\ifvmode` this one has no argument and just tells if we're in vertical mode.

```
\vbox {
  \noindent \ifhmode h\else v\fi mode
  \par
  \ifhmode h\else \noindent v\fi mode
}
```

You can use it for instance to trigger injection of code, or prevent that some content (or command) is done more than once:

```
hmode
vmode
```

## 2.2.8 `\ifmmode`

Math is something very  $\TeX$  so naturally you can check if you're in math mode. here is an example of using this test:

```
\def\enforcemath#1{\ifmmode#1\else$ #1 $\fi}
```

Of course in reality macros that do such things are more advanced than this one.

## 2.2.9 `\ifinner`

```
\def\ShowMode
  {\ifhmode      \ifinner inner \fi hmode
  \else\ifvmode \ifinner inner \fi vmode
  \else\ifmmode \ifinner inner \fi mmode
  \else        \ifinner inner \fi unset
  \fi\fi\fi}
```

```
\ShowMode \ShowMode
```

```
\vbox{\ShowMode}
```

```
\hbox{\ShowMode}
```

```
$_ShowMode$
```

```
$$\ShowMode$$
```

The first line has two tests, where the first one changes the mode to horizontal simply because a text has been typeset. Watch how display math is not inner.

```
vmode hmode
inner vmode
inner hmode
innermmode
innermmode
```

By the way, moving the `\ifinner` test outside the branches (to the top of the macro) won't work because once the word `inner` is typeset we're no longer in vertical mode, if we were at all.

### 2.2.10 `\ifvoid`

A box is one of the basic concepts in  $\TeX$ . In order to understand this primitive we present four cases:

```
\setbox0\hbox{}           \ifvoid0 void \else content \fi
\setbox0\hbox{123}       \ifvoid0 void \else content \fi
\setbox0\hbox{} \box0    \ifvoid0 void \else content \fi
\setbox0\hbox to 10pt{} \ifvoid0 void \else content \fi
```

In the first case, we have a box which is empty but it's not void. It helps to know that internally an hbox is actually an object with a pointer to a linked list of nodes. So, the first two can be seen as:

```
hlist -> [nothing]
hlist -> 1 -> 2 -> 3 -> [nothing]
```

but in any case there is a hlist. The third case puts something in a hlist but then flushes it. Now we have not even the hlist any more; the box register has become void. The last case is a variant on the first. It is an empty box with a given width. The outcome of the four lines (with a box flushed in between) is:

```
content
content
```

```
void
content
```

So, when you want to test if a box is really empty, you need to test also its dimensions, which can be up to three tests, depending on your needs.

```
\setbox0\emptybox           \ifvoid0 void\else content\fi
\setbox0\emptybox          \wd0=10pt \ifvoid0 void\else content\fi
\setbox0\hbox to 10pt {}    \ifvoid0 void\else content\fi
\setbox0\hbox               {} \wd0=10pt \ifvoid0 void\else content\fi
```

Setting a dimension of a void (empty) box doesn't make it less void:

```
void
void
content
content
```

### 2.2.11 `\ifhbox`

This test takes a box number and gives true when it is an hbox.

### 2.2.12 `\ifvbox`

This test takes a box number and gives true when it is an vbox. Both a `\vbox` and `\vtop` are vboxes, the difference is in the height and depth and the baseline. In a `\vbox` the last line determines the baseline

```
vbox or vtop
```

```
vtop or vbox
```

And in a `\vtop` the first line takes control:

```
vbox or vtop
```

```
vtop or vbox
```

but, once wrapped, both internally are just vlists.

### 2.2.13 `\ifx`

This test is actually used a lot in `ConTeXt`: it compares two token(list)s:

```

\ifx a b Y\else N\fi
\ifx ab Y\else N\fi
\def\A {a}\def\B{b}\ifx \A\B Y\else N\fi
\def\A{aa}\def\B{a}\ifx \A\B Y\else N\fi
\def\A {a}\def\B{a}\ifx \A\B Y\else N\fi

```

Here the result is: “NNNNY”. It does not expand the content, if you want that you need to use an `\edef` to create two (temporary) macros that get compared, like in:

```
\edef\TempA{...}\edef\TempB{...}\ifx\TempA\TempB ...\else ...\fi
```

### 2.2.14 `\ifeof`

This test checks if a the pointer in a given input channel has reached its end. It is also true when the file is not present. The argument is a number which relates to the `\openin` primitive that is used to open files for reading.

### 2.2.15 `\iftrue`

It does what it says: always true.

### 2.2.16 `\iffalse`

It does what it says: always false.

### 2.2.17 `\ifcase`

The general layout of an `\ifcase` tests is as follows:

```
\ifcase<number>
  when zero
\or
  when one
\or
  when two
\or
  ...
\else
  when something else
\fi
```

As in other places a number is a sequence of signs followed by one or more digits

## 2.3 $\varepsilon$ -T<sub>E</sub>X primitives

### 2.3.1 `\ifdefined`

This primitive was introduced for checking the existence of a macro (or primitive) and with good reason. Say that you want to know if `\MyMacro` is defined? One way to do that is:

```
\ifx\MyMacro\undefined
  {\bf undefined indeed}
\fi
```

This results in: **undefined indeed**, but is this macro really undefined? When T<sub>E</sub>X scans your source and sees a the escape character (the forward slash) it will grab the next

characters and construct a control sequence from it. Then it finds out that there is nothing with that name and it will create a hash entry for a macro with that name but with no meaning. Because `\undefined` is also not defined, these two macros have the same meaning and therefore the `\ifx` is true. Imagine that you do this many times, with different macro names, then your hash can fill up. Also, when a user defined `\undefined` you're suddenly get a different outcome.

In order to catch the last problem there is the option to test directly:

```
\ifdefined\MyOtherMacro \else
  {\bf also undefined}
\fi
```

This (or course) results in: **also undefined**, but the macro is still sort of defined (with no meaning). The next section shows how to get around this.

### 2.3.2 `\ifcurname`

A macro is often defined using a ready made name, as in:

```
\def\OhYes{yes}
```

The name is made from characters with catcode letter which means that you cannot use for instance digits or underscores unless you also give these characters that catcode, which is not that handy in a document. You can however use `\curname` to define a control sequence with any character in the name, like:

```
\expandafter\def\curname Oh Yes : 1\endcurname{yes}
```

Later on you can get this one with `\curname`:

```
\curname Oh Yes : 1\endcurname
```

However, if you say:

```
\curname Oh Yes : 2\endcurname
```

you won't get some result, nor a message about an undefined control sequence, but the name triggers a define anyway, this time not with no meaning (undefined) but as equivalent to `\relax`, which is why

```
\expandafter\ifx\curname Oh Yes : 2\endcurname\relax
  {\bf relaxed indeed}
```

```
\fi
```

is the way to test its existence. As with the test in the previous section, this can deplete the hash when you do lots of such tests. The way out of this is:

```
\ifcsname Oh Yes : 2\endcsname \else
  {\bf unknown indeed}
\fi
```

This time there is no hash entry created and therefore there is not even an undefined control sequence.

In Lua<sub>T</sub><sub>E</sub><sub>X</sub> there is an option to return false in case of a messy expansion during this test, and in LuaMeta<sub>T</sub><sub>E</sub><sub>X</sub> that is default. This means that tests can be made quite robust as it is pretty safe to assume that names that make sense are constructed from regular characters and not boxes, font switches, etc.

### 2.3.3 `\iffontchar`

This test was also part of the  $\varepsilon$ -<sub>T</sub><sub>E</sub><sub>X</sub> extensions and it can be used to see if a font has a character.

```
\iffontchar\font`A
  {\em This font has an A!}
\fi
```

And, as expected, the outcome is: “*This font has an A!*”. The test takes two arguments, the first being a font identifier and the second a character number, so the next checks are all valid:

```
\iffontchar\font`A yes\else nop\fi\par
\iffontchar\nullfont`A yes\else nop\fi\par
\iffontchar\textfont0`A yes\else nop\fi\par
```

In the perspective of LuaMeta<sub>T</sub><sub>E</sub><sub>X</sub> I considered also supporting `\fontid` but it got a bit messy due to the fact that this primitive expands in a different way so this extension was rejected.

### 2.3.4 `\unless`

You can negate the results of a test by using the `\unless` prefix, so for instance you can replace:



```

\ifdim\scratchdimen=10pt
  \dosomething
\else\ifdim\scratchdimen<10pt
  \dosomething
\fi\fi

```

by:

```

\unless\ifdim\scratchdimen>10pt
  \dosomething
\fi

```

An `\unless` makes little sense when used with `\ifcase` but contrary to the other engines we don't error or it; we just give a warning. Some conditionals internally use a case so there we can actually provide a variant:

```

\ifcase 1 \relax zero \or one \or two \else else \fi = one \par
\ifcase 2 \relax zero \or one \or two \else else \fi = two \par

```

```

\unless\ifcase 1 \relax zero \or one \or two \else else \fi % warning
\unless\ifcase 2 \relax zero \or one \or two \else else \fi % warning

```

```

\ifchkdim1pt\or yes \else nop \fi = yes \par
\ifchkdim2 \or nop \else yes \fi = yes \par

```

```

\unless\ifchkdim1pt\or nop \else yes \fi = yes \par
\unless\ifchkdim2 \or yes \else nop \fi = yes \par

```

The `\ifchkdim`, `\ifchkdimension`, `\ifchknum`, `\ifchknumber` and `\ifparameter` are supported.

```

one = one
two = two
one two
yes = yes
yes = yes
yes = yes
yes = yes

```

## 2.4 Lua<sub>T</sub><sub>E</sub>X primitives

### 2.4.1 `\ifincsname`

As it had no real practical usage it might get dropped in LuaMeta<sub>T</sub><sub>E</sub>X, so it will not be discussed here.

### 2.4.2 `\ifprimitive`

As it had no real practical usage due to limitations, this one is not available in LuaMeta<sub>T</sub><sub>E</sub>X so it will not be discussed here. If really needed you can use `\ifflags`.

### 2.4.3 `\ifabsnum`

This test is inherited from pdf<sub>T</sub><sub>E</sub>X and behaves like `\ifnum` but first turns a negative number into a positive one.

### 2.4.4 `\ifabsdim`

This test is inherited from pdf<sub>T</sub><sub>E</sub>X and behaves like `\ifdim` but first turns a negative dimension into a positive one.

### 2.4.5 `\ifcondition`

This is not really a test but in order to unstand that you need to know how <sub>T</sub><sub>E</sub>X internally deals with tests.

```
\ifdimen\scratchdimen>10pt
  \ifdim\scratchdimen<20pt
    result a
  \else
    result b
  \fi
\else
  result c
\fi
```

When we end up in the branch of “result a” we need to skip two `\else` branches after we're done. The `\if..` commands increment a level while the `\fi` decrements a level.

The `\else` needs to be skipped here. In other cases the true branch needs to be skipped till we end up at the right `\else`. When doing this skipping,  $\TeX$  is not interested in what it encounters beyond these tokens and this skipping (therefore) goes real fast but it does see nested conditions and doesn't interpret grouping related tokens.

A side effect of this is that the next is not working as expected:

```
\def\ifmorethan{\ifdim\scratchdimen>}
\def\iflessthan{\ifdim\scratchdimen<}

\ifmorethan10pt
  \iflessthan20pt
    result a
  \else
    result b
  \fi
\else
  result c
\fi
```

The `\iflessthan` macro is not seen as an `\if...` so the nesting gets messed up. The solution is to fool the scanner in thinking that it is. Say we have:

```
\scratchdimen=25pt

\def\ifmorethan{\ifdim\scratchdimen>}
\def\iflessthan{\ifdim\scratchdimen<}
```

and:

```
\ifcondition\ifmorethan10pt
  \ifcondition\iflessthan20pt
    result a
  \else
    result b
  \fi
\else
  result c
\fi
```

When we expand this snippet we get: “result b” and no error concerning a failure in locating the right `\fi`'s. So, when scanning the `\ifcondition` is seen as a valid

`\if...` but when the condition is really expanded it gets ignored and the `\ifmorethan` has better come up with a match or not.

In this perspective it is also worth mentioning that nesting problems can be avoided this way:

```
\def\WhenTrue {something \iftrue ...}
\def\WhenFalse{something \iffalse ...}

\ifnum\scratchcounter>123
  \let\next\WhenTrue
\else
  \let\next\WhenFalse
\fi
\next
```

This trick is mentioned in The  $\TeX$ book and can also be found in the plain  $\TeX$  format. A variant is this:

```
\ifnum\scratchcounter>123
  \expandafter\WhenTrue
\else
  \expandafter\WhenFalse
\fi
```

but using `\expandafter` can be quite intimidating especially when there are multiple in a row. It can also be confusing. Take this: an `\ifcondition` expects the code that follows to produce a test. So:

```
\def\ifwhatever#1%
  {\ifdim#1>10pt
   \expandafter\iftrue
  \else
   \expandafter\iffalse
  \fi}

\ifcondition\ifwhatever{10pt}
  result a
\else
  result b
\fi
```

This will not work! The reason is in the already mentioned fact that when we end up in the greater than 10pt case, the scanner will happily push the `\iftrue` after the `\fi`, which is okay, but when skipping over the `\else` it sees a nested condition without matching `\fi`, which makes it fail. I will spare you a solution with lots of nasty tricks, so here is the clean solution using `\ifcondition`:

```
\def\truecondition {\iftrue}
\def\falsecondition{\iffalse}

\def\ifwhatever#1%
  {\ifdim#1>10pt
   \expandafter\truecondition
  \else
   \expandafter\falsecondition
  \fi}

\ifcondition\ifwhatever{10pt}
  result a
\else
  result b
\fi
```

It will be no surprise that the two macros at the top are predefined in `ConTeXt`. It might be more of a surprise that at the time of this writing the usage in `ConTeXt` of this `\ifcondition` primitive is rather minimal. But that might change.

As a further teaser I'll show another simple one,

```
\def\HowOdd#1{\unless\ifnum\numexpr ((#1):2)*2\relax=\numexpr#1\relax}

\ifcondition\HowOdd{1}very \else not so \fi odd
\ifcondition\HowOdd{2}very \else not so \fi odd
\ifcondition\HowOdd{3}very \else not so \fi odd
```

This renders:

```
very odd
not so odd
very odd
```

The code demonstrates several tricks. First of all we use `\numexpr` which permits more complex arguments, like:

```
\ifcondition\HowOdd{4+1}very \else not so \fi odd
```

```
\ifcondition\HowOdd{2\scratchcounter+9}very \else not so \fi odd
```

Another trick is that we use an integer division (the `:`) which is an operator supported by LuaMetaTeX.

## 2.5 LuaMetaTeX primitives

### 2.5.1 `\ifnum` and `ifdim`

These have been extended with a few more operators. For instance, we can use a negation:

```
\ifnum 10 > 5 Y\else N\fi
\ifnum 10 !> 5 Y\else N\fi
```

Results in: YN. A bitwise comparison is possible too:

```
\ifnum "02 & 2 Y\else N\fi
\ifnum "02 & 4 Y\else N\fi
\ifnum "02 !& 8 Y\else N\fi
```

yields: YNY. You can also use the Unicode variants `€`, `ℳ`, `≠`, `≤`, `≥`, `≠`, and `≠`.

### 2.5.2 `\iffloat`

This is a test for a float, much like a test for a `dimen` without unit.

### 2.5.3 `\ifabsfloat`

This is a test for a float, much like a test for a `dimen` without unit.

### 2.5.4 `\ifintervalnum`

This is a test for equality of two numbers within an interval, as in:

```
\ifintervalnum 1 2 1 Y\else N\fi
\ifintervalnum 1 3 1 Y\else N\fi
\ifintervalnum 100 102 1 Y\else N\fi
\ifintervalnum 100 102 3 Y\else N\fi
```

which results in: YNNY.

### 2.5.5 `\ifintervaldim`

This is a test for equality of two dimensions within an interval, as in:

```
\ifintervaldim 1pt 2pt 1pt Y\else N\fi
\ifintervaldim 1pt 3pt 1pt Y\else N\fi
\ifintervaldim 100pt 102pt 1pt Y\else N\fi
\ifintervaldim 100pt 102pt 3pt Y\else N\fi
```

We get: YNNY.

### 2.5.6 `\ifintervalfloat`

This is a test for a float, much like a test for a `dimen` without unit.

### 2.5.7 `\ifdimexpression`

This is a boolean checker so the comparison is done as part of the expression, as in:

```
\ifdimexpression{10pt > (4pt + 8pt)}Y\else N\fi
```

### 2.5.8 `\ifnumexpression`

This is a boolean checker so the comparison is done as part of the expression, as in:

```
\ifnumexpression{10 > (4 + 8)}Y\else N\fi
```

### 2.5.9 `\ifcmpnum`

This one is part of a set of three tests that all are a variant of a `\ifcase` test. A simple example of the first test is this:

```
\ifcmpnum 123 345 less \or equal \else more \fi
```

The test scans for two numbers, which of course can be registers or expressions, and sets the case value to 0, 1 or 2, which means that you then use the normal `\or` and `\else` primitives for follow up on the test.

### 2.5.10 `\ifchknum`

This test scans a number and when it's okay sets the case value to 1, and otherwise to 2. So you can do the next:

```
\ifchknum 123\or good \else bad \fi
\ifchknum bad\or good \else bad \fi
```

An error message is suppressed and the first `\or` can be seen as a sort of recovery token, although in fact we just use the fast scanner mode that comes with the `\ifcase`: because the result is 1 or 2, we never see invalid tokens.

In order to avoid another scan the a valid result it is made available in `\lastchknumber`.

### 2.5.11 `\ifchknumber`

This one is a more rigorous variant of `\ifchknum` and doesn't like trailing non numeric crap.

### 2.5.12 `\ifchknumexpr`

This test goes a bit further and accepts an expression.

```
\ifchknumexpr 123 + 45\or good \else bad \fi
```

As with the other checkers, if there is a valid result it is available in `\lastchknumber`.

### 2.5.13 `\ifnumval`

A sort of combination of the previous two is `\ifnumval` which checks a number but also if it's less, equal or more than zero:

```
\ifnumval 123\or less \or equal \or more \else error \fi
\ifnumval bad\or less \or equal \or more \else error \fi
```

You can decide to ignore the bad number or do something that makes more sense. Often the to be checked value will be the content of a macro or an argument like `#1`.

### 2.5.14 `\ifcmpdim`

This test is like `\ifcmpnum` but for dimensions.

### 2.5.15 `\ifchkdim`

This test is like `\ifchknum` but for dimensions. The last checked value is available as `\lastchknumber`.



### 2.5.16 `\ifchkdimension`

This one is a more rigorous variant of `\ifchkdim` and doesn't like trailing rubbish.

### 2.5.17 `\ifchkdimexpr`

This test is like `\ifchknumexpr` but for dimensions. The last checked value is available as `\lastchkdimension`.

### 2.5.18 `\ifdimval`

This test is like `\ifnumval` but for dimensions. The last checked value is available as `\lastchkdim`

### 2.5.19 `\iftok`

Although this test is still experimental it can be used. What happens is that two to be compared ‘things’ get scanned for. For each we first gobble spaces and `\relax` tokens. Then we can have several cases:

1. When we see a left brace, a list of tokens is scanned up to the matching right brace.
2. When a reference to a token register is seen, that register is taken as value.
3. When a reference to an internal token register is seen, that register is taken as value.
4. When a macro is seen, its definition becomes the to be compared value.
5. When a number is seen, the value of the corresponding register is taken

An example of the first case is:

```
\iftok {abc} {def}%
...
\else
...
\fi
```

The second case goes like this:

```
\iftok\scratchtoksone\scratchtokstwo
...
\else
...
\fi
```

Case one and four mixed:

```
\iftok{123}\TempX
...
\else
...
\fi
```

The last case is more a catch: it will issue an error when no number is given. Eventually that might become a bit more clever (depending on our needs.)

### 2.5.20 \ifzeronum, \ifzerodim, \ifzerofloat

The names of these three tells what they do: checking for a zero value.

```
(\ifzerodim 10pt\norelax A\orelse\ifzerodim 0pt\norelax B\else C\fi)
(\ifzeronum 10 \norelax A\orelse\ifzeronum 0 \norelax B\else C\fi)
(\ifzerofloat 10.0\norelax A\orelse\ifzerofloat 0.0\norelax B\else C\fi)
```

Here we use the `\norelax` to get rid of trailing spaces: (B) (B) (B).

### 2.5.21 \ifhaschar, \ifhastok, \ifhastoks, \ifhasxtoks

These checkers can be used to identify a (sequence) of token(s) in a given token list. Their working can best be shown with a few examples:

```
\ifhaschar c {abcd}Y\else N\fi
\ifhastok c {abcd}Y\else N\fi
\ifhastoks {c}{abcd}Y\else N\fi
\ifhasxtoks {c}{abcd}Y\else N\fi

\def\abcd{abcd}

\ifhaschar c {\abcd}Y\else N\fi
\ifhastok c {\abcd}Y\else N\fi
\ifhastoks {c}{\abcd}Y\else N\fi
\ifhasxtoks {c}{\abcd}Y\else N\fi

\ifhaschar c {a{bc}d}Y\else N\fi
\ifhastok c {a{bc}d}Y\else N\fi
\ifhastoks {c}{a{bc}d}Y\else N\fi
\ifhasxtoks {c}{a{bc}d}Y\else N\fi
```

```
\def\abcd{a{bc}d}
```

```
\ifhaschar c {\abcd}Y\else N\fi
```

```
\ifhastok c {\abcd}Y\else N\fi
```

```
\ifhastoks {c}{\abcd}Y\else N\fi
```

```
\ifhasxtoks {c}{\abcd}Y\else N\fi
```

YYYY

NNNY

NYYY

NNNY

The `\ifhaschar` test will not descend into a braced sublist. The `x` variants expand the list before comparison.

### 2.5.22 `\ifcstok`

There is a subtle difference between this one and `\iftok`: spaces and `\relax` tokens are skipped but nothing gets expanded. So, when we arrive at the to be compared ‘things’ we look at what is there, as-is.

### 2.5.23 `\iffrozen`

*This is an experimental test.* Commands can be defined with the `\frozen` prefix and this test can be used to check if that has been the case.

### 2.5.24 `\ifprotected`

Commands can be defined with the `\protected` prefix (or in `ConTeXt`, for historic reasons, with `\unexpanded`) and this test can be used to check if that has been the case.

### 2.5.25 `\ifarguments`

This conditional can be used to check how many arguments were matched. It only makes sense when used with macros defined with the `\tolerant` prefix and/or when the sentinel `\ignorearguments` after the arguments is used. More details can be found in the lowlevel macros manual.

### 2.5.26 `\ifrelax`

The following tests all return the same: YYY; it is a shortcut for `\ifx ... \relax` that looks nicer in code.

```

\ifrelax\relax           Y\else N\fi
\ifrelax\norelax        Y\else N\fi
\expandafter\ifrelax\c
sname ReLaX\endcsname Y\else N\fi

```

### 2.5.27 `\ifempty`

This is again a shortcut, this time for `\ifx ... \empty` assuming that `\empty` is defined as being nothing. Instead of a token you can also pass a list, so here we get YNY.

```

\ifempty{}              Y\else N\fi
\ifempty{!}             Y\else N\fi
\ifempty\empty         Y\else N\fi

```

### 2.5.28 `\iflastnamedcs`

This test is part of the `\csname` repertoire and uses the last valid result from such a command.

```

\def\Hello{upper}
\def\hello{lower}
\ifcsname Hello\endcsname
  \iflastnamedcs\hello
    world
  \orelse\iflastnamedcs\Hello
    World
\fi
\fi

```

Here the ‘Hello’ test result in ‘World’. It is an example of a follow up test, most likely used in user interfacing.

### 2.5.29 `\ifboolean`

Another new one is the following: it tests a number for being zero or not. As with any primitive that scans for a number, it accepts a braced expression too.

```
(\ifboolean 0 T\else F\fi)
(\ifboolean 1 T\else F\fi)
(\ifboolean {(2 * 4) < 5} T\else F\fi)
(\ifboolean \dimexpression{(1em > 20pt) or (1ex > 15pt)} T\else F\fi)
(\ifboolean \dimexpression{(1em > 3pt) and (1ex < 3pt)} T\else F\fi)
```

We get: (F) (T) (F) (F) (F).

### 2.5.30 \iflist

The \ifvoid test doesn't really test for a box being empty, which is why we have an additional primitive. Compare the following:

```
\setbox0\hbox{}
\setbox2\hbox{!}
\setbox4\emptybox % \box\voidbox
\setbox8\box6

\wd0 10pt \wd2 10pt \wd4 10pt \wd6 10pt

[\ifvoid0 Y\else N\fi \iflist0 Y\else N\fi \the\wd0] % empty hbox
[\ifvoid2 Y\else N\fi \iflist2 Y\else N\fi \the\wd2] % hbox with content
[\ifvoid4 Y\else N\fi \iflist4 Y\else N\fi \the\wd4] % no box
[\ifvoid6 Y\else N\fi \iflist6 Y\else N\fi \the\wd6] % no box
```

The result demonstrates that we check if there is any content at all, independent of dimensions or the presence of a wrapping list node.

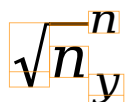
```
[NN10.0pt] [NY10.0pt] [YN0.0pt] [YN0.0pt]
```

### 2.5.31 \ifcramped

This test relates to math and in particular to four of the eight states:

```
\im {
  \sqrt
    {\ifcramped\mathstyle y\else n\fi}
  ^ {\ifcramped\mathstyle y\else n\fi}
  _ {\ifcramped\mathstyle y\else n\fi}
}
```

Because a math formula is first read and then processed in several passes you need to be aware of this state not always being easily predictable because there can be a delay between that read and successive treatments.



### 2.5.32 `\ifmathparameter`

The next example demonstrates what this test provides:

```
[\ifmathparameter\Umathextrasubspace \displaystyle zero\or set\else
  unset\fi]
[\ifmathparameter\Umathaccentbaseheight\displaystyle zero\or set\else
  unset\fi]
[\ifmathparameter\Umathaccentbasedepth \displaystyle zero\or set\else
  unset\fi]
```

There are three possible outcomes; here we get: [zero] [set] [set]. In LuaMetaT<sub>E</sub>X we have more math parameters than in LuaT<sub>E</sub>X, and some are set in font specific so called ‘goodie’ files.

### 2.5.33 `\ifmathstyle`

Here you need to keep in mind that you test the style that is set when T<sub>E</sub>X scans for formula. Processing happens afterwards and then styles can change.

```
{\ifmathstyle D\or D'\or T\or T'\or S\or S'\or SS\or SS'\else ?\fi}
\im{\ifmathstyle D\or D'\or T\or T'\or S\or S'\or SS\or SS'\else ?\fi}
\dm{\ifmathstyle D\or D'\or T\or T'\or S\or S'\or SS\or SS'\else ?\fi}
```

We get: ? *T D*. The odd values are cramped.

### 2.5.34 `\ifinalignment`

This test is an experimental one:

```
\halign \bgroup
  \aligncontent
  \aligntab
  \aligncontent
```

```

\cr
one \aligntab \ifinalignment two\else three\fi \cr
\noalign{\ifinalignment yes\else no\fi}
one \aligntab \hbox{\ifinalignment two\else three\fi} \cr
\egroup

\hbox{\ifinalignment two\else three\fi}

```

We get:

```

one two
yes
one two
three

```

### 2.5.35 \ifinsert

This primitive checks if an insert box has content. Usage depends on the macro package so for instance in ConT<sub>E</sub>Xt, after `\footnote{A note.}` you can actually check it with:

```
\setupheadertexts[\ifinsert\namedinsertionnumber{footnote} Y\else N\fi]
```

You pass the number of a insert class and in this example the content, set by the page builder, hasn't yet been flushed.

### 2.5.36 \ifflags

This one related to interfacing. When a macro is defined, one can apply several prefixes to that macro. Some of these prefixes result in a specific kind of macro, for instance a protected, tolerant, tolerant protected, or regular macro. When a macro is defined global, its (internal) level value indicates that. In addition macros, or actually any control sequence, also the built-in ones, can have a set of flags. Some, have consequences in the engine, so for instance an untraced macro will present itself as a primitive, without details that clutter a log. Other flags get meaning when the overload protection mechanisms are enabled.

Testing flags can give some insight but in ConT<sub>E</sub>Xt there is little reason to use this test other than for illustrative purposes. Take this definition

```
\global\protected\def\Foo{Foo}
```

This macro is internally represented as follows; here we used `\meaning` is:

```
\global \protected \def \Foo {Foo}
```

When we use `\meaning` we get:

```
protected macro:Foo
```

With `\meaningfull` we get:

```
global protected macro:Foo
```

Here is how you can test what properties and flags are set.

```
\ifflags\Foo\global    global    \fi
\ifflags\Foo\protected protected \fi
\ifflags\Foo\tolerant  tolerant  \fi
```

We only show a few tests here:

```
global protected
```

Instead of a prefix you can also pass a number:

```
\ifflags\relax\primitiveflagcode primitive \fi
\ifflags\relax\permanentflagcode permanent \fi
```

```
primitive
```

In `ConTEXt` many macros are defined as permanent which in terms of overload protection has the same impact. Relevant flag values are available in `tex.getflagvalues()` but in `ConTEXt` we prefer predefined constants:

```
\aliasedflagcode, \conditionalflagcode, \constantflagcode, \deferredflagcode,
\frozenflagcode, \globalflagcode, \immediateflagcode, \immutableflagcode,
\inheritedflagcode, \instanceflagcode, \mutableflagcode, \noalignedflagcode,
\overloadedflagcode, \permanentflagcode, \primitiveflagcode, \protectedflagcode,
\semiprotectedflagcode, \tolerantflagcode, \untracedflagcode, \valueflagcode
```

### 2.5.37 `\ifparameters`

This is an `\ifcase` where the number is the number of parameters passed to the current macro. Of course, when used in a macro one should be aware of the fact that another macro call will change this number.



### 2.5.38 `\ifparameter`

This test checks if a parameter has been set, and it's used as follows:

```
\ifparameter#4\or set\else unset\fi
```

because `#4` is actually a reference it refers to the parameter in the current macro and is not influenced by nested macro calls which makes it more reliable than a `\ifparameters` test.

### 2.5.39 `\orelse`

This is not really a test primitive but it does act that way. Say that we have this:

```
\ifdim\scratchdimen>10pt
  case 1
\else\ifdim\scratchdimen<20pt
  case 2
\else\ifcount\scratchcounter>10
  case 3
\else\ifcount\scratchcounter<20
  case 4
\fi\fi\fi\fi
```

A bit nicer looks this:

```
\ifdim\scratchdimen>10pt
  case 1
\orelse\ifdim\scratchdimen<20pt
  case 2
\orelse\ifcount\scratchcounter>10
  case 3
\orelse\ifcount\scratchcounter<20
  case 4
\fi
```

We stay at the same level. Sometimes a more flat test tree had advantages but if you think that it gives better performance then you will be disappointed. The fact that we stay at the same level is compensated by a bit more parsing, so unless you have millions such cases (or expansions) it might make a bit of a difference. As mentioned, I'm a bit sensitive for how code looks so that was the main motivation for introducing it. There is a companion `\orunless` continuation primitive.

A rather neat trick is the definition of `\quitcondition`:

```
\def\quitcondition{\orelse\iffalse}
```

This permits:

```
\ifdim\scratchdimen>10pt  
  case 1a  
  \quitcondition  
  case 4b  
\fi
```

where, of course, the quitting normally is the result of some intermediate extra test. But let me play safe here: beware of side effects.

### 2.5.40 `\orunless`

This is the negated variant of `\orelse`.

## 2.6 For the brave

### 2.6.1 Full expansion

If you don't understand the following code, don't worry. There is seldom much reason to go this complex but obscure T<sub>E</sub>X code attracts some users so ...

When you have a macro that has for instance assignments, and when you expand that macro inside an `\edef`, these assignments are not actually expanded but tokenized. In LuaMetaT<sub>E</sub>X there is a way to apply these assignments without side effects and that feature can be used to write a fully expandable user test. For instance:

```
\def\truecondition {\iftrue}  
\def>falsecondition{\iffalse}  
  
\def\fontwithidhaschar#1#2%  
  {\beginlocalcontrol  
    \scratchcounter\numexpr\fontid\font\relax  
    \setfontid\numexpr#1\relax  
    \endlocalcontrol  
    \iffontchar\font\numexpr#2\relax  
    \beginlocalcontrol
```

```

\setfontid\scratchcounter
\endlocalcontrol
\expandafter\truecondition
\else
\expandafter\falsecondition
\fi}

```

The `\iffontchar` test doesn't handle numeric font id, simply because at the time it was added to  $\varepsilon$ -TeX, there was no access to these id's. Now we can do:

```

\edef\foo{\fontwithidhaschar{1} {75}yes\else nop\fi} \meaning\foo
\edef\foo{\fontwithidhaschar{1}{999}yes\else nop\fi} \meaning\foo

[\ifcondition\fontwithidhaschar{1} {75}yes\else nop\fi]
[\ifcondition\fontwithidhaschar{1}{999}yes\else nop\fi]

```

These result in:

```

macro:yes
macro:nop

```

```

[yes]
[nop]

```

If you remove the `\immediateassignment` in the definition above then the typeset results are still the same but the meanings of `\foo` look different: they contain the assignments and the test for the character is actually done when constructing the content of the `\edef`, but for the current font. So, basically that test is now useless.

## 2.6.2 User defined if's

There is a `\newif` macro that defines three other macros:

```

\newif\ifOnMyOwnTerms

```

After this, not only `\ifOnMyOwnTerms` is defined, but also:

```

\OnMyOwnTermstrue
\OnMyOwnTermsfalse

```

These two actually are macros that redefine `\ifOnMyOwnTerms` to be either equivalent to `\iftrue` and `\iffalse`. The (often derived from plain TeX) definition of `\newif` is a

bit if a challenge as it has to deal with removing the `if` in order to create the two extra macros and also make sure that it doesn't get mixed up in a catcode jungle.

In ConTEXt we have a variant:

```
\newconditional\MyConditional
```

that can be used with:

```
\settrue\MyConditional
\setfalse\MyConditional
```

and tested like:

```
\ifconditional\MyConditional
  ...
\else
  ...
\fi
```

This one is cheaper on the hash and doesn't need the two extra macros per test. The price is the use of `\ifconditional`, which is *not* to confused with `\ifcondition` (it has bitten me already a few times).

## 2.7 Relaxing

When TEX scans for a number or dimension it has to check tokens one by one. On the case of a number, the scanning stops when there is no digit, in the case of a dimension the unit determine the end of scanning. In the case of a number, when a token is not a digit that token gets pushed back. When digits are scanned a trailing space or `\relax` is pushed back. Instead of a number of dimension made from digits, periods and units, the scanner also accepts registers, both the direct accessors like `\count` and `\dimen` and those represented by one token. Take these definitions:

```
\newdimen\MyDimenA \MyDimenA=1pt \dimen0=\MyDimenA
\newdimen\MyDimenB \MyDimenB=2pt \dimen2=\MyDimenB
```

I will use these to illustrate the side effects of scanning. Watch the spaces in the result.

First I show what effect we want to avoid. When second argument contains a number (digits) the zero will become part of it so we actually check `\dimen00` here.

```
\def\whatever#1#2%
```

```
{\ifdim#1=#20\else1\fi}
```

```
\whatever{1pt}{2pt}          [macro:1]
\whatever{1pt}{1pt}          [macro:0]
\whatever{\dimen 0}{\dimen 2} [macro:1]
\whatever{\dimen 0}{\dimen 0} [macro:]
\whatever\MyDimenA\MyDimenB  [macro:1]
\whatever\MyDimenA\MyDimenB  [macro:1]
```

The solution is to add a space but watch how that one can end up in the result:

```
\def\whatever#1#2%
  {\ifdim#1=#2 0\else1\fi}
```

```
\whatever{1pt}{2pt}          [macro:1]
\whatever{1pt}{1pt}          [macro:0]
\whatever{\dimen 0}{\dimen 2} [macro:1]
\whatever{\dimen 0}{\dimen 0} [macro:0]
\whatever\MyDimenA\MyDimenB  [macro:1]
\whatever\MyDimenA\MyDimenB  [macro:1]
```

A variant is using `\relax` and this time we get this token retained in the output.

```
\def\whatever#1#2%
  {\ifdim#1=#2\relax0\else1\fi}
```

```
\whatever{1pt}{2pt}          [macro:1]
\whatever{1pt}{1pt}          [macro:\relax 0]
\whatever{\dimen 0}{\dimen 2} [macro:1]
\whatever{\dimen 0}{\dimen 0} [macro:\relax 0]
\whatever\MyDimenA\MyDimenB  [macro:1]
\whatever\MyDimenA\MyDimenB  [macro:1]
```

A solution that doesn't have side effects of forcing the end of a number (using a space or `\relax` is one where we use expressions. The added overhead of scanning expressions is taken for granted because the effect is what we like:

```
\def\whatever#1#2%
  {\ifdim\dimexpr#1\relax=\dimexpr#2\relax0\else1\fi}
```

```
\whatever{1pt}{2pt}          [macro:1]
\whatever{1pt}{1pt}          [macro:0]
\whatever{\dimen 0}{\dimen 2} [macro:1]
```

```

\whatever{\dimen 0}{\dimen 0} [macro:0]
\whatever\MyDimenA\MyDimenB [macro:1]
\whatever\MyDimenA\MyDimenB [macro:1]

```

Just for completeness we show a more obscure trick: this one hides assignments to temporary variables. Although performance is okay, it is the least efficient one so far.

```

\def\whatever#1#2%
  {\beginlocalcontrol
   \MyDimenA#1\relax
   \MyDimenB#2\relax
  \endlocalcontrol
  \ifdim\MyDimenA=\MyDimenB0\else1\fi}

\whatever{1pt}{2pt} [macro:1]
\whatever{1pt}{1pt} [macro:0]
\whatever{\dimen 0}{\dimen 2} [macro:1]
\whatever{\dimen 0}{\dimen 0} [macro:0]
\whatever\MyDimenA\MyDimenB [macro:1]
\whatever\MyDimenA\MyDimenB [macro:1]

```

It is kind of a game to come up with alternatives but for sure those involve dirty tricks and more tokens (and runtime). The next can be considered a dirty trick too: we use a special variant of `\relax`. When a number is scanned it acts as `\relax`, but otherwise it just is ignored and disappears.

```

\def\whatever#1#2%
  {\ifdim#1=#2\norelax0\else1\fi}

\whatever{1pt}{2pt} [macro:1]
\whatever{1pt}{1pt} [macro:0]
\whatever{\dimen 0}{\dimen 2} [macro:1]
\whatever{\dimen 0}{\dimen 0} [macro:0]
\whatever\MyDimenA\MyDimenB [macro:1]
\whatever\MyDimenA\MyDimenB [macro:1]

```

## 2.7 Colofon

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Support	<a href="http://www.pragma-ade.com">www.pragma-ade.com</a> <a href="http://contextgarden.net">contextgarden.net</a> <a href="mailto:ntg-context@ntg.nl">ntg-context@ntg.nl</a>

## **3 Boxes**



low level

TEX

boxes

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### 3.1 Introduction

An average Con $\text{T}_{\text{E}}\text{X}$ t user will not use the low level box primitives but a basic understanding of how  $\text{T}_{\text{E}}\text{X}$  works doesn't hurt. In fact, occasionally using a box command might bring a solution not easily achieved otherwise, simply because a more high level interface can also be in the way.

The best reference is of course The  $\text{T}_{\text{E}}\text{X}$ book so if you're really interested in the details you should get a copy of that book. Below I will not go into details about all kind of glues, kerns and penalties, just boxes it is.

This explanation will be extended when I feel the need (or users have questions that can be answered here).

### 3.2 Boxes

This paragraph of text is made from lines that contain words that themselves contain symbolic representations of characters. Each line is wrapped in a so called horizontal box and eventually those lines themselves get wrapped in what we call a vertical box.

When we expose some details of a paragraph it looks like this:

H	This is a rather narrow	LH:0.000 LS:0.000	RS:13.888	RH:0.000
H	paragraph blown up a	LH:0.000 LS:0.000	RS:19.008	RH:0.000
H	bit. Here we use a flush	LH:0.000 LS:0.000	RS:7.007	RH:0.000
H	left, aka ragged right,	LH:0.000 LS:0.000	RS:19.706	RH:0.000
H	approach.	LH:0.000 LS:0.000	RS:0.000	RH:0.000

The left only shows the boxes, the variant at the right shows (font) kerns and glue too. Because we flush left, there is rather strong right skip glue at the right boundary of the box. If font kerns show up depends on the font, not all fonts have them (or have only a few). The glyphs themselves are also kind of boxed, as their dimensions determine the area that they occupy:

This is a rather ...

But, internally they are not really boxed, as they already are a single quantity. The same is true for rules: they are just blobs with dimensions. A box on the other hand wraps a linked list of so called nodes: glyphs, kerns, glue, penalties, rules, boxes, etc. It is a container with properties like width, height, depth and shift.

### 3.3 T<sub>E</sub>X primitives

The box model is reflected in T<sub>E</sub>X's user interface but not by that many commands, most noticeably `\hbox`, `\vbox` and `\vtop`. Here is an example of the first one:

```
\hbox width 10cm{text}
\hbox width 10cm height 1cm depth 5mm{text}
text \raise5mm\hbox{text} text
```

The `\raise` and `\lower` commands behave the same but in opposite directions. One could as well have been defined in terms of the other.

```
text \raise 5mm \hbox to 2cm {text}
text \lower -5mm \hbox to 2cm {text}
text \raise -5mm \hbox to 2cm {text}
text \lower 5mm \hbox to 2cm {text}
```

```

text.      text.      text.      text.
      |-----|      |-----|
      |text|          |text|
      |-----|          |-----|

```

A box can be moved to the left or right but, believe it or not, in ConT<sub>E</sub>Xt we never use that feature, probably because the consequences for the width are such that we can as well use kerns. Here are some examples:

```
text \vbox{\moveleft 5mm \hbox {left}}text !
text \vbox{\moveright 5mm \hbox{right}}text !
```

lefttext ! text      righttext !

```
text \vbox{\moveleft 25mm \hbox {left}}text !
text \vbox{\moveright 25mm \hbox{right}}text !
```

left      text text ! text                      righttext !

Code like this will produce a complaint about an underfull box but we can easily get around that:

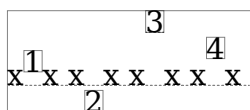
```
text \raise 5mm \hbox to 2cm {\hss text}
text \lower -5mm \hbox to 2cm {text\hss}
text \raise -5mm \hbox to 2cm {\hss text}
text \lower 5mm \hbox to 2cm {text\hss}
```

The `\hss` primitive injects a glue that when needed will fill up the available space. So, here we force the text to the right or left.

Instead of `\raise` you can also provide the shift (up or down) with a keyword:

```
\ruledhbox\bggroup
x\raise 5pt\ruledhbox                      {1}x
x\raise -10pt\ruledhbox                    {2}x
x\raise -5pt\ruledhbox shift -20pt{3}x
x\ruledhbox                                shift -10pt{4}x
```

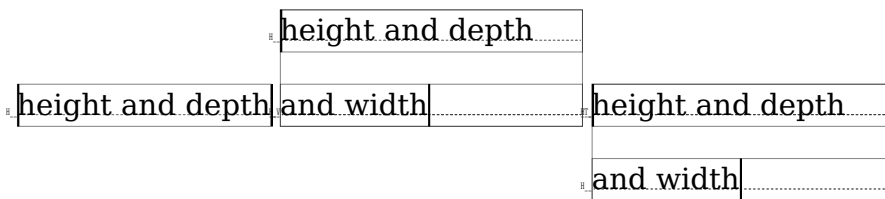
`\egroup`



We have three kind of boxes: `\hbox`, `\vbox` and `\vtop`. Actually we have a fourth type `\dbox` but that is a variant on `\vbox` to which we come back later.

```
\hbox{\strut height and depth\strut}
\vbox{\hsize 4cm \strut height and depth\par and width\strut}
\vtop{\hsize 4cm \strut height and depth\par and width\strut}
```

A `\vbox` aligns at the bottom and a `\vtop` at the top. I have added some so called struts to enforce a consistent height and depth. A strut is an invisible quantity (consider it a black box) that enforces consistent line dimensions: height and depth.



You can store a box in a register but you need to be careful not to use a predefined one. If you need a lot of boxes you can reserve some for your own:

```
\newbox\MySpecialBox
```

but normally you can do with one of the scratch registers, like 0, 2, 4, 6 or 8, for local boxes, and 1, 3, 5, 7 and 9 for global ones. Registers are used like:

```
\setbox0\hbox{here}
\global\setbox1\hbox{there}
```

In ConT<sub>E</sub>Xt you can also use

```
\setbox\scratchbox \hbox{here}
\setbox\scratchboxone\hbox{here}
\setbox\scratchboxtwo\hbox{here}
```

and some more. In fact, there are quite some predefined scratch registers (boxes, dimensions, counters, etc). Feel free to investigate further.

When a box is stored, you can consult its dimensions with `\wd`, `\ht` and `\dp`. You can of course store them for later use.

```
\scratchwidth \wd\scratchbox
\scratchheight\ht\scratchbox
\scratchdepth \dp\scratchbox
\scratchtotal \dimexpr\ht\scratchbox+\dp\scratchbox\relax
\scratchtotal \htdp\scratchbox
```

The last line is ConT<sub>E</sub>Xt specific. You can also set the dimensions

```
\wd\scratchbox 10cm
\ht\scratchbox 10mm
\dp\scratchbox 5mm
```

So you can cheat! A box is placed with `\copy`, which keeps the original intact or `\box` which just inserts the box and then wipes the register. In practice you seldom need a

copy, which is more expensive in runtime anyway. Here we use copy because it serves the examples.

```
\copy\scratchbox
\box \scratchbox
```

### 3.4 $\varepsilon$ - $\text{T}_{\text{E}}\text{X}$ primitives

The  $\varepsilon$ - $\text{T}_{\text{E}}\text{X}$  extensions don't add something relevant for boxes, apart from that you can use the expressions mechanism to mess around with their dimensions. There is a mechanism for typesetting r2l within a paragraph but that has limited capabilities and doesn't change much as it's mostly a way to trick the backend into outputting a stretch of text in the other direction. This feature is not available in Lua $\text{T}_{\text{E}}\text{X}$  because it has an alternative direction mechanism.

### 3.5 Lua $\text{T}_{\text{E}}\text{X}$ primitives

The concept of boxes is the same in Lua $\text{T}_{\text{E}}\text{X}$  as in its predecessors but there are some aspects to keep in mind. When a box is typeset this happens in Lua $\text{T}_{\text{E}}\text{X}$ :

1. A list of nodes is constructed. In Lua $\text{T}_{\text{E}}\text{X}$  this is a double linked list (so that it can easily be manipulated in Lua) but  $\text{T}_{\text{E}}\text{X}$  itself only uses the forward links.
2. That list is hyphenated, that is: so called discretionary nodes are injected. This depends on the language properties of the glyph (character) nodes.
3. Then ligatures are constructed, if the font has such combinations. When this built-in mechanism is used, in Con $\text{T}_{\text{E}}\text{X}$ t we speak of base mode.
4. After that inter-character kerns are applied, if the font provides them. Again this is a base mode action.
5. Finally the box gets packaged:
  - In the case of a horizontal box, the list is packaged in a hlist node, basically one liner, and its dimensions are calculated and set.
  - In the case of a vertical box, the paragraph is broken into one or more lines, without hyphenation, with optimal hyphenation or in the worst case with so called emergency stretch applied, and the result becomes a vlist node with its dimensions set.

In traditional  $\text{T}_{\text{E}}\text{X}$  the first four steps are interwoven but in  $\text{LuaT}_{\text{E}}\text{X}$  we need them split because the step 5 can be overloaded by a callback. In that case steps 3 and 4 (and maybe 2) are probably also overloaded, especially when you bring handling of fonts under Lua control.

New in  $\text{LuaT}_{\text{E}}\text{X}$  are three packers: `\hpack`, `\vpack` and `\tpack`, which are companions to `\hbox`, `\vbox` and `\vtop` but without the callbacks applied. Using them is a bit tricky as you never know if a callback should be applied, which, because users can often add their own Lua code, is not something predictable.

Another box related extension is `direction`. There are four possible directions but because in  $\text{LuaMetaT}_{\text{E}}\text{X}$  there are only two. Because this model has been upgraded, it will be discussed in the next section. A  $\text{ConT}_{\text{E}}\text{Xt}$  user is supposed to use the official  $\text{ConT}_{\text{E}}\text{Xt}$  interfaces in order to be downward compatible.

### 3.6 $\text{LuaMetaT}_{\text{E}}\text{X}$ primitives

There are two possible directions: left to right (the default) and right to left for Hebrew and Arabic. Here is an example that shows how it'd done with low level directives:

```
\hbox direction 0 {from left to right}
\hbox direction 1 {from right to left}
```

from left to right  
tfel ot thgir morf

A low level direction switch is done with:

```
\hbox direction 0
  {from left to right \textdirection 1 from right to left}
\hbox direction 1
  {from right to left \textdirection 1 from left to right}
```

from left to right tfel ot thgir morf  
thgir ot tfel morf tfel ot thgir morf

but actually this is kind of *not done* in  $\text{ConT}_{\text{E}}\text{Xt}$ , because there you are supposed to use the proper direction switches:

```
\naturalhbox {from left to right}
\reversehbox {from right to left}
\naturalhbox {from left to right \righttoleft from right to left}
```

`\reversebox {from right to left \lefttoright from left to right}`

from left to right

tfel ot thgir morf

from left to right tfel ot thgir morf

from left to right tfel ot thgir morf

Often more is needed to properly support right to left typesetting so using the ConT<sub>E</sub>Xt commands is more robust.

In LuaMetaT<sub>E</sub>X the box model has been extended a bit, this as a consequence of dropping the vertical directional typesetting, which never worked well. In previous sections we discussed the properties width, height and depth and the shift resulting from a `\raise`, `\lower`, `\moveleft` and `\moveright`. Actually, the shift is also used in for instance positioning math elements.

The way shifting influences dimensions can be somewhat puzzling. Internally, when T<sub>E</sub>X packages content in a box there are two cases:

- When a horizontal box is made, and height - shift is larger than the maximum height so far, that delta is taken. When depth + shift is larger than the current depth, then that depth is adapted. So, a shift up influences the height and a shift down influences the depth.
- In the case of vertical packaging, when width + shift is larger than the maximum box (line) width so far, that maximum gets bumped. So, a shift to the right can contribute, but a shift to the left cannot result in a negative width. This is also why vertical typesetting, where height and depth are swapped with width, goes wrong: we somehow need to map two properties onto one and conceptually T<sub>E</sub>X is really set up for horizontal typesetting. (And it's why I decided to just remove it from the engine.)

This is one of these cases where T<sub>E</sub>X behaves as expected but it also means that there is some limitation to what can be manipulated. Setting the shift using one of the four commands has a direct consequence when a box gets packaged which happens immediately because the box is an argument to the foursome.

There is in traditional T<sub>E</sub>X, probably for good reason, no way to set the shift of a box, if only because the effect would normally be none. But in LuaT<sub>E</sub>X we can cheat, and therefore, for educational purposed ConT<sub>E</sub>Xt has implements some cheats.

We use this sample box:



```

\setbox\scratchbox\hbox\bgroup
  \middlegray\vrule width 20mm depth -.5mm height 10mm
  \hskip-20mm
  \darkgray \vrule width 20mm height -.5mm depth 5mm
\egroup

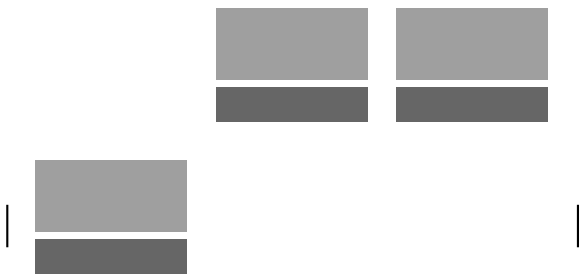
```

When we mess with the shift using the ConT<sub>E</sub>Xt `\shiftbox` helper, we see no immediate effect. We only get the shift applied when we use another helper, `\hpackbox`.

```

\hbox\bgroup
  \showstruts \strut
  \quad \copy\scratchbox
  \quad \shiftbox\scratchbox -20mm \copy\scratchbox
  \quad \hpackbox\scratchbox \box \scratchbox
  \quad \strut
\egroup

```



When instead we use `\vpackbox` we get a different result. This time we move left.

```

\hbox\bgroup
  \showstruts \strut
  \quad \copy\scratchbox
  \quad \shiftbox\scratchbox -10mm \copy\scratchbox
  \quad \vpackbox\scratchbox \copy\scratchbox
  \quad \strut
\egroup

```



The shift is set via Lua and the repackaging is also done in Lua, using the low level `hpack` and `vpack` helpers and these just happen to look at the shift when doing their job. At the T<sub>E</sub>X end this never happens.

This long exploration of shifting serves a purpose: it demonstrates that there is not that much direct control over boxes apart from their three dimensions. However this was never a real problem as one can just wrap a box in another one and use kerns to move the embedded box around. But nevertheless I decided to see if the engine can be a bit more helpful, if only because all that extra wrapping gives some overhead and complications when we want to manipulate boxes. And of course it is also a nice playground.

We start with changing the direction. Changing this property doesn't require repackaging because directions are not really dealt with in the frontend. When a box is converted to (for instance pdf) the reversion happens.

```
\setbox\scratchbox\hbox{whatever}
\the\boxdirection\scratchbox: \copy\scratchbox \crlf
\boxdirection\scratchbox 1
\the\boxdirection\scratchbox: \copy\scratchbox
```

0: whatever

1: revetahw

Another property that can be queried and set is an attribute. In order to get a private attribute we define one.

```
\newattribute\MyAt
\setbox\scratchbox\hbox attr \MyAt 123 {whatever}
[\the\boxattribute\scratchbox\MyAt]
\boxattribute\scratchbox\MyAt 456
[\the\boxattribute\scratchbox\MyAt]
[\ifnum\boxattribute\scratchbox\MyAt>400 okay\fi]
```

[123] [456] [okay]

The sum of the height and depth is available too. Because for practical reasons setting that property is also needed then, the choice was made to distribute the value equally over height and depth.

```
\setbox\scratchbox\hbox {height and depth}
[\the\ht\scratchbox]
[\the\dp\scratchbox]
[\the\boxtotal\scratchbox]
\boxtotal\scratchbox=20pt
[\the\ht\scratchbox]
```

```
[\the\dp\scratchbox]
[\the\boxtotal\scratchbox]
```

```
[8.35742pt] [2.44385pt] [10.80127pt] [10.0pt] [10.0pt] [20.0pt]
```

We've now arrived to a set of properties that relate to each other. They are a bit complex and given the number of possibilities one might need to revert to some trial and error: orientations and offsets. As with the dimensions, directions and attributes, they are passed as box specification. We start with the orientation.

```
\hbox \bgroup \showboxes
      \hbox orientation 0 {right}
\quad \hbox orientation 1 {up}
\quad \hbox orientation 2 {left}
\quad \hbox orientation 3 {down}
```

```
\egroup
```

The diagram shows four boxes arranged in a row. From left to right: a box containing the word 'right', a box containing 'dn' (representing up), a box containing 'left', and a box containing 'down'. Each box is a small rectangle with its content centered inside.

When the orientation is set, you can also set an offset. Where shifting around a box can have consequences for the dimensions, an offset is virtual. It gets effective in the backend, when the contents is converted to some output format.

```
\hbox \bgroup \showboxes
      \hbox orientation 0 yoffset 10pt {right}
\quad \hbox orientation 1 xoffset 10pt {up}
\quad \hbox orientation 2 yoffset -10pt {left}
\quad \hbox orientation 3 xoffset -10pt {down}
```

```
\egroup
```

The diagram shows four boxes arranged in a row. From left to right: a box containing the word 'right' shifted to the right, a box containing 'dn' shifted to the right, a box containing 'left' shifted to the left, and a box containing 'down' shifted to the left. Each box is a small rectangle with its content centered inside.

The reason that offsets are related to orientation is that we need to know in what direction the offsets have to be applied and this binding forces the user to think about it. You can also set the offsets using commands.

```
\setbox\scratchbox\hbox{whatever}%
1 \copy\scratchbox
2 \boxorientation\scratchbox 2 \copy\scratchbox
```

```

3 \boxxoffset \scratchbox -15pt \copy\scratchbox
4 \boxyoffset \scratchbox -15pt \copy\scratchbox
5

```

```

1 whatever2
3
4
5

```

```

\setbox\scratchboxone\hbox{whatever}%
\setbox\scratchboxtwo\hbox{whatever}%

```

```

1 \boxxoffset \scratchboxone -15pt \copy\scratchboxone
2 \boxyoffset \scratchboxone -15pt \copy\scratchboxone
3 \boxxoffset \scratchboxone -15pt \copy\scratchboxone
4 \boxyoffset \scratchboxone -15pt \copy\scratchboxone
5 \boxxmove \scratchboxtwo -15pt \copy\scratchboxtwo
6 \boxymove \scratchboxtwo -15pt \copy\scratchboxtwo
7 \boxxmove \scratchboxtwo -15pt \copy\scratchboxtwo
8 \boxymove \scratchboxtwo -15pt \copy\scratchboxtwo

```

```

whatever 2 3 4 whatever6 7 8
whatever whatever whatever whatever
whatever
whatever

```

The move commands are provided as convenience and contrary to the offsets they do adapt the dimensions. Internally, with the box, we register the orientation and the offsets and when you apply these commands multiple times the current values get overwritten. But ... because an orientation can be more complex you might not get the effects you expect when the options we discuss next are used. The reason is that we store the original dimensions too and these come into play when these other options are used: anchoring. So, normally you will apply an orientation and offsets once only.

The orientation specifier is actually a three byte number that best can be seen hexadecimal (although we stay within the decimal domain). There are three components: x-anchoring, y-anchoring and orientation:

```
0x<X><Y><O>
```

or in TeX speak:

```
"<X><Y><O>
```

The landscape and seascape variants both sit on top of the baseline while the flipped variant has its depth swapped with the height. Although this would be enough a bit more control is possible.

The vertical options of the horizontal variants anchor on the baseline, lower corner, upper corner or center.

```
\ruledhbox orientation "002 {\TEX} and
\ruledhbox orientation "012 {\TEX} and
\ruledhbox orientation "022 {\TEX} and
\ruledhbox orientation "032 {\TEX}
```

The horizontal options of the horizontal variants anchor in the center, left, right, halfway left and halfway right.

```
\ruledhbox orientation "002 {\TEX} and
\ruledhbox orientation "102 {\TEX} and
\ruledhbox orientation "202 {\TEX} and
\ruledhbox orientation "302 {\TEX} and
\ruledhbox orientation "402 {\TEX}
```

The orientation has consequences for the dimensions so they are dealt with in the expected way in constructing lines, paragraphs and pages, but the anchoring is virtual, like the offsets. There are two extra variants for orientation zero: on top of baseline or below, with dimensions taken into account.

```
\ruledhbox orientation "000 {\TEX} and
\ruledhbox orientation "004 {\TEX} and
\ruledhbox orientation "005 {\TEX}
```

The anchoring can look somewhat confusing but you need to keep in mind that it is normally only used in very controlled circumstances and not in running text. Wrapped in macros users don't see the details. We're talking boxes here, so for instance:

```
test\quad
\hbox orientation 3 \bgroup
  \strut test\hbox orientation "002 \bgroup\strut test\egroup test%
\egroup \quad
\hbox orientation 3 \bgroup
  \strut test\hbox orientation "002 \bgroup\strut test\egroup test%
```

```

\egroup \quad
\hbox orientation 3 \bgroup
  \strut test\hbox orientation "012 \bgroup\strut test\egroup test%
\egroup \quad
\hbox orientation 3 \bgroup
  \strut test\hbox orientation "022 \bgroup\strut test\egroup test%
\egroup \quad
\hbox orientation 3 \bgroup
  \strut test\hbox orientation "032 \bgroup\strut test\egroup test%
\egroup \quad
\hbox orientation 3 \bgroup
  \strut test\hbox orientation "042 \bgroup\strut test\egroup test%
\egroup
\quad test

```



Where a `\vtop` has the baseline at the top, a `\vbox` has it at the bottom. In LuaMeta-TeX we also have a `\dbox`, which is a `\vbox` with that behaves like a `\vtop` when it's appended to a vertical list: the height of the first box or rule determines the (base)line correction that gets applied. The following example demonstrates this:

<pre> XXXXXXXXXXXXXXXXXXXX We thrive in information-thick worlds be- cause of our marvelous and everyday ca- pacity to select, edit, single out, struc- ture, highlight, group, pair, merge, har- monize, synthesize, focus, organize, con- dense, reduce, boil down, choose, cate- gorize, catalog, classify, list, abstract, scan, look into, idealize, isolate, discrim- inate, distinguish, screen, pigeonhole, pick over, sort, integrate, blend, inspect, filter, lump, skip, smooth, chunk, aver- age, approximate, cluster, aggregate, outline, summarize, itemize, review, dip into, flip through, browse, glance into, leaf through, skim, refine, enumerate, glean, synopsisize, winnow the wheat from the chaff and separate the sheep from the goats. XXXXXXXXXXXXXXXXXXXX </pre>	<pre> XXXXXXXXXXXXXXXXXXXX We thrive in information-thick worlds be- cause of our marvelous and everyday ca- pacity to select, edit, single out, struc- ture, highlight, group, pair, merge, har- monize, synthesize, focus, organize, con- dense, reduce, boil down, choose, cate- gorize, catalog, classify, list, abstract, scan, look into, idealize, isolate, discrim- inate, distinguish, screen, pigeonhole, pick over, sort, integrate, blend, inspect, filter, lump, skip, smooth, chunk, aver- age, approximate, cluster, aggregate, outline, summarize, itemize, review, dip into, flip through, browse, glance into, leaf through, skim, refine, enumerate, glean, synopsisize, winnow the wheat from the chaff and separate the sheep from the goats. XXXXXXXXXXXXXXXXXXXX </pre>	<pre> XXXXXXXXXXXXXXXXXXXX We thrive in information-thick worlds be- cause of our marvelous and everyday ca- pacity to select, edit, single out, struc- ture, highlight, group, pair, merge, har- monize, synthesize, focus, organize, con- dense, reduce, boil down, choose, cate- gorize, catalog, classify, list, abstract, scan, look into, idealize, isolate, discrim- inate, distinguish, screen, pigeonhole, pick over, sort, integrate, blend, inspect, filter, lump, skip, smooth, chunk, aver- age, approximate, cluster, aggregate, outline, summarize, itemize, review, dip into, flip through, browse, glance into, leaf through, skim, refine, enumerate, glean, synopsisize, winnow the wheat from the chaff and separate the sheep from the goats. XXXXXXXXXXXXXXXXXXXX </pre>
<code>\vbox</code>	<code>\vtop</code>	<code>\dbox</code>

The `d` stands for ‘dual’ because we (sort of) have two baselines. The regular height and depth are those of a `\vbox`.

## 3.7 Splitting

When you feed T<sub>E</sub>X a paragraph of text it will accumulate the content in a list of nodes. When the paragraph is finished by `\par` or an empty line it will be fed into the par builder that will try to break the lines as good as possible. Normally that paragraph will be added to the page and at some point there can be breaks between lines in order not to overflow the page. When you collect the paragraph in a box you can use `\vsplit` to emulate this.

```
\setbox\scratchbox\vbox{\samplefile{tufte}}
```

```
\startlinecorrection
```

```
\ruledhbox{\vsplit\scratchbox to 2\lineheight}
```

```
\stoplinecorrection
```

We thrive in information-thick worlds because of our marvelous and everyday capacity to select, edit, single out, structure, highlight, group, pair, merge, harmonize, synthe-

The split off box is given the specified height, but in LuaMetaT<sub>E</sub>X you can also get the natural dimensions:

```
\setbox\scratchbox\vbox{\samplefile{tufte}}
```

```
\startlinecorrection
```

```
\ruledhbox{\vsplit\scratchbox upto 2\lineheight}
```

```
\stoplinecorrection
```

We thrive in information-thick worlds because of our marvelous and everyday capacity to select, edit, single out, structure, highlight, group, pair, merge, harmonize, synthe-

We can force a resulting box type by using `\vsplit`, `\tsplit` and `\dsplit` (here we use the visualized variants):

```
\setbox\scratchbox\vbox{\samplefile{tufte}}
```

```
\startlinecorrection
```

```
\ruledtsplit \scratchbox upto 2\lineheight
```

```
\stoplinecorrection
```

We thrive in information-thick worlds because of our marvelous and everyday capacity to select, edit, single out, structure, highlight, group, pair, merge, harmonize, synthe-

```
\setbox\scratchbox\vbox{\samplefile{tufte}}
```

```

\startlinecorrection
\ruledvsplit \scratchbox upto 2\lineheight
\stoplinecorrection

```

We thrive in information-thick worlds because of our marvelous and everyday capacity to select, edit, single out, structure, highlight, group, pair, merge, harmonize, synthe-

```

\setbox\scratchbox\vbox{\samplefile{tufte}}

```

```

\startlinecorrection
\ruleddsplit \scratchbox upto 2\lineheight
\stoplinecorrection

```

We thrive in information-thick worlds because of our marvelous and everyday capacity to select, edit, single out, structure, highlight, group, pair, merge, harmonize, synthe-

The engine provides vertical splitters but ConT<sub>E</sub>Xt itself also has a horizontal one.<sup>1</sup>

```

\starttexdefinition Test #1#2#3
  \par
  \dontleavehmode
  \strut
  \llap{{\infofont #2}\quad}
  \blackrule[width=#2,color=darkblue]
  \par
  \setbox\scratchbox\hbox{\samplefile{#1}}
  \hsplit\scratchbox
    to          #2
    depth       \strutdp
    height      \strutht
    shrinkcriterium #3 % badness
  \par
\stoptexdefinition

\dostepwiserecurse {100} {120} {2} {
  \Test{tufte}{#1mm}{1000}
  \Test{tufte}{#1mm}{-100}
}

```

100mm

We thrive in information-thick worlds because of

<sup>1</sup> At some point I might turn that one into a native engine primitive.



100mm We thrive in information-thick worlds because of

102mm We thrive in information-thick worlds because of our

102mm We thrive in information-thick worlds because of our

104mm We thrive in information-thick worlds because of our

104mm We thrive in information-thick worlds because of our

106mm We thrive in information-thick worlds because of our

106mm We thrive in information-thick worlds because of our

108mm We thrive in information-thick worlds because of our

108mm We thrive in information-thick worlds because of our

110mm We thrive in information-thick worlds because of our

110mm We thrive in information-thick worlds because of our

112mm We thrive in information-thick worlds because of our mar-

112mm We thrive in information-thick worlds because of our mar-

114mm We thrive in information-thick worlds because of our mar-

114mm We thrive in information-thick worlds because of our mar-

116mm We thrive in information-thick worlds because of our mar-

116mm We thrive in information-thick worlds because of our mar-

118mm We thrive in information-thick worlds because of our mar-

118mm We thrive in information-thick worlds because of our mar-

120mm We thrive in information-thick worlds because of our mar-

120mm

We thrive in information-thick worlds because of our mar-

A split off box gets packed at its natural size and a badness as well as overshoot amount is calculated. When the overshoot is positive and the the badness is larger than the stretch criterium, the box gets repacked to the natural size. The same happens when the overshoot is negative and the badness exceeds the shrink criterium. When the overshoot is zero (basically we have a fit) but the badness still exceeds the stretch or shrink we also repack. Indeed this is a bit fuzzy, but so is badness.

```
\starttexdefinition Test #1#2#3
  \par
  \dontleavehmode
  \strut
  \llap{{\infofont #2}\quad}
  \blackrule[width=#2,color=darkblue]
  \par
  \setbox\scratchbox\hbox{\samplefile{#1}}
  \doloop {
    \ifvoid\scratchbox
      \exitloop
    \else
      \hsplit\scratchbox
        to      #2
        depth  \strutdp
        height \strutht
        #3
      \par
      \allowbreak
    \fi
  }
\stoptexdefinition

\Test{tufte}{100mm}{shrinkcriterium 1000}
\Test{tufte}{100mm}{shrinkcriterium 0}
\Test{tufte}{100mm}{}

```

100mm

We thrive in information-thick worlds because of our marvelous and everyday capacity to select, edit, single out, structure, highlight, group, pair, merge,

harmonize, synthesize, focus, organize, condense, reduce, boil down, choose, categorize, catalog, classify, list, abstract, scan, look into, idealize, isolate, discriminate, distinguish, screen, pigeonhole, pick over, sort, integrate, blend, inspect, filter, lump, skip, smooth, chunk, average, approximate, cluster, aggregate, outline, summarize, itemize, review, dip into, flip through, browse, glance into, leaf through, skim, refine, enumerate, glean, synopsise, winnow the wheat from the chaff and separate the sheep from the goats.

100mm

We thrive in information-thick worlds because of our marvelous and everyday capacity to select, edit, single out, structure, highlight, group, pair, merge, harmonize, synthesize, focus, organize, condense, reduce, boil down, choose, categorize, catalog, classify, list, abstract, scan, look into, idealize, isolate, discriminate, distinguish, screen, pigeonhole, pick over, sort, integrate, blend, inspect, filter, lump, skip, smooth, chunk, average, approximate, cluster, aggregate, outline, summarize, itemize, review, dip into, flip through, browse, glance into, leaf through, skim, refine, enumerate, glean, synopsise, winnow the wheat from the chaff and separate the sheep from the goats.

100mm

We thrive in information-thick worlds because of our marvelous and everyday capacity to select, edit, single out, structure, highlight, group, pair, merge, harmonize, synthesize, focus, organize, condense, reduce, boil down, choose, categorize, catalog, classify, list, abstract, scan, look into, idealize, isolate, discriminate, distinguish, screen, pigeonhole, pick over, sort, integrate, blend, inspect, filter, lump, skip, smooth, chunk, average, approximate, cluster, aggregate, outline, summarize, itemize, review, dip into, flip through, browse, glance into, leaf through, skim, refine, enumerate, glean, synopsise, winnow the wheat from the chaff and separate the sheep from the goats.

BBB  
size, winnow the wheat from the chaff and separate  
the sheep from the goats.

Watch how the last line get stretched when we set the criterium to zero. I'm sure that users will find reasons to abuse this effect.

### 3.7 Colofon

Author Hans Hagen  
ConT<sub>E</sub>Xt 2025.07.04 21:26  
LuaMetaT<sub>E</sub>X 2.11.07 | 20250705  
Support [www.pragma-ade.com](http://www.pragma-ade.com)  
[contextgarden.net](http://contextgarden.net)  
[ntg-context@ntg.nl](mailto:ntg-context@ntg.nl)

111111

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AAA

## 4 Expansion

low level

TEX

expansion

## Contents

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### 4.1 Preamble

This short manual demonstrates a couple of properties of the macro language. It is not an in-depth philosophical expose about macro languages, tokens, expansion and such that some T<sub>E</sub>Xies like. I prefer to stick to the practical aspects. Occasionally it will be technical but you can just skip those paragraphs (or later return to them) when you can't follow the explanation. It's often not that relevant. I won't talk in terms of mouth, stomach and gut the way the T<sub>E</sub>Xbook does and although there is no way to avoid the word 'token' I will do my best to not complicate matters by too much token speak. Examples show best what we mean.

### 4.2 T<sub>E</sub>X primitives

The T<sub>E</sub>X language provides quite some commands and those built in are called primitives. User defined commands are called macros. A macro is a shortcut to a list of primitives and/or macro calls. All can be mixed with characters that are to be typeset somehow.

```
\def\MyMacro{b}
```

```
a\MyMacro c
```

When T<sub>E</sub>X reads this input the a gets turned into a glyph node with a reference to the current font set and the character a. Then the parser sees a macro call, and it will enter another input level where it expands this macro. In this case it sees just an b and it will give this the same treatment as the a. The macro ends, the input level decrements and the c gets its treatment.

Before we move on to more examples and differences between engines, it is good to stress that `\MyMacro` is not a primitive command: we made our command here. The b actually can be seen as a sort of primitive because in this macro it gets stored as so

called token with a primitive property. That primitive property can later on be used to determine what to do. More explicit examples of primitives are `\hbox`, `\advance` and `\relax`. It will be clear that ConT<sub>E</sub>Xt extends the repertoire of primitive commands with a lot of macro commands. When we typeset a source using module `m-scite` the primitives come out dark blue.

The amount of primitives differs per engine. It all starts with T<sub>E</sub>X as written by Don Knuth. Later  $\epsilon$ -T<sub>E</sub>X added some more primitives and these became official extensions adopted by other variants of T<sub>E</sub>X. The pdfT<sub>E</sub>X engine added quite some and as follow up on that LuaT<sub>E</sub>X added more but didn't add all of pdfT<sub>E</sub>X. A few new primitives came from Omega (Aleph). The LuaMetaT<sub>E</sub>X engine drops a set of primitives that comes with LuaT<sub>E</sub>X and adds plenty new ones. The nature of this engine (no backend and less frontend) makes that we need to implement some primitives as macros. But the basic set is what good old T<sub>E</sub>X comes with.

Internally these so called primitives are grouped in categories that relate to their nature. They can be directly expanded (a way of saying that they get immediately interpreted) or delayed (maybe stored for later usage). They can involve definitions, calculations, setting properties and values or they can result in some typesetting. This is what makes T<sub>E</sub>X confusing to new users: it is a macro programming language, an interpreter but at the same time an executor of typesetting instructions.

A group of primitives is internally identified as a command (they have a `cmd` code) and the sub commands are flagged by their `chr` code. This sounds confusing but just thing of the fact that most of what we input are characters and therefore they make up most sub commands. For instance the 'letter `cmd`' is used for characters that are seen as letters that can be used in the name of user commands, can be typeset, are valid for hyphenation etc. The letter related `cmd` can have many `chr` codes (all of Unicode). I'd like to remark that the grouping is to a large extend functional, so sometimes primitives that you expect to be similar in nature are in different groups. This has to do with the fact that T<sub>E</sub>X needs to be able to determine efficiently if a primitive is operating (or forbidden) in horizontal, vertical and/or math mode.

There are more than 150 internal `cmd` groups. if we forget about the mentioned character related ones, some, have only a few sub commands (`chr`) and others many more (just consider all the OpenType math spacing related parameters). A handful of these commands deal with what we call macros: user defined combinations of primitives and other macros, consider them little programs. The `\MyMacro` example above is an example. There are differences between engines. In standard T<sub>E</sub>X there are `\outer` and `\long` commands, and most engines have these. However, in LuaMetaT<sub>E</sub>X the later to be discussed `\protected` macros have their own specific 'call `cmd`'. Users don't need to bother about this.



So, when from now on we talk about primitives, we mean the built in, hard coded commands, and when we talk about macros we mean user commands. Although internally there are less cmd categories than primitives, from the perspective of the user they are all unique. Users won't consult the source anyway but when they do they are warned. Also, when in LuaMetaTeX you use the low level interfacing to TeX you have to figure out these subtle aspects because there this grouping does matter.

Before we continue I want to make clear that expansion (as discussed in this document) can refer to a macro being expanded (read: its meaning gets injected into the input, so the engine kind of sidetracks from what it was doing) but also to direct consequences of running into a primitive. However, users only need to consider expansion in the perspective of macros. If a user has `\advance` in the input it immediately gets done. But when it's part of a macro definition it only is executed when the macro expands. A good check in (traditional) TeX is to compare what happens in `\def` and `\edef` which is why we will use these two in the upcoming examples. You put something in a macro and then check what `\meaning` or `\show` reports.

Now back to user defined macros. A macro can contain references to macros so in practice the input can go several levels up and some applications push back a lot so this is why your TeX input stack can be configured to be huge.

```
\def\MyMacroA{ and }
\def\MyMacroB{1\MyMacroA 2}

a\MyMacroA b
```

When `\MyMacroB` is defined, its body gets three so called tokens: the character token `1` with property 'other', a token that is a reference to the macro `\MyMacroB`, and a character token `2`, also with property 'other'. The meaning of `\MyMacroA` is five tokens: a reference to a space token, then three character tokens with property 'letter', and finally a space token.

```
\def \MyMacroA{ and }
\edef\MyMacroB{1\MyMacroA 2}

a\MyMacroA b
```

In the second definition an `\edef` is used, where the `e` indicates expansion. This time the meaning gets expanded immediately. So we get effectively the same as in:

```
\def\MyMacroB{1 and 2}
```

Characters are easy: they just expand to themselves or trigger adding a glyph node, but not all primitives expand to their meaning or effect.

```
\def\MyMacroA{\scratchcounter = 1 }
\def\MyMacroB{\advance\scratchcounter by 1}
\def\MyMacroC{\the\scratchcounter}
```

```
\MyMacroA a
\MyMacroB b
\MyMacroB c
\MyMacroB d
\MyMacroC
```

a b c d 4

```
macro:\scratchcounter = 1
macro:\advance \scratchcounter by 1
macro:\the \scratchcounter
```

Let's assume that \scratchcounter is zero to start with and use \edef's:

```
\edef\MyMacroA{\scratchcounter = 1 }
\edef\MyMacroB{\advance\scratchcounter by 1}
\edef\MyMacroC{\the\scratchcounter}
```

```
\MyMacroA a
\MyMacroB b
\MyMacroB c
\MyMacroB d
\MyMacroC
```

a b c d 0

```
macro:\scratchcounter = 1
macro:\advance \scratchcounter by 1
macro:0
```

So, this time the third macro has its meaning frozen, but we can prevent this by applying a \noexpand when we do this:

```
\edef\MyMacroA{\scratchcounter = 1 }
\edef\MyMacroB{\advance\scratchcounter by 1}
\edef\MyMacroC{\noexpand\the\scratchcounter}
```

```
\MyMacroA a
\MyMacroB b
```

```
\MyMacroB c
\MyMacroB d
\MyMacroC
```

```
a b c d 4
```

```
macro:\scratchcounter = 1
macro:\advance \scratchcounter by 1
macro:\the \scratchcounter
```

Of course this is a rather useless example but it serves its purpose: you'd better be aware what gets expanded immediately in an `\edef`. In most cases you only need to worry about `\the` and embedded macros (and then of course their meanings).

You can also store tokens in a so-called token register. Here we use a predefined scratch register:

```
\def\MyMacroA{ and }
\def\MyMacroB{1\MyMacroA 2}
\scratchtoks {\MyMacroA}
```

The content of `\scratchtoks` is: “`\MyMacroA`”, so no expansion has happened here.

```
\def\MyMacroA{ and }
\def\MyMacroB{1\MyMacroA 2}
\scratchtoks \expandafter {\MyMacroA}
```

Now the content of `\scratchtoks` is: “`and`”, so this time expansion has happened.

```
\def\MyMacroA{ and }
\def\MyMacroB{1\MyMacroA 2}
\scratchtoks \expandafter {\MyMacroB}
```

Indeed the macro gets expanded but only one level: “`1\MyMacroA 2`”. Compare this with:

```
\def\MyMacroA{ and }
\edef\MyMacroB{1\MyMacroA 2}
\scratchtoks \expandafter {\MyMacroB}
```

The trick is to expand in two steps with an intermediate `\edef`: “`1 and 2`”. Later we will see that other engines provide some more expansion tricks. The only way to get some grip on expansion is to just play with it.

The `\expandafter` primitive expands the token (which can be a macro) standing after the next next one and then injects its meaning into the stream. So:

```
\expandafter \MyMacroA \MyMacroB
```

works okay. In a normal document you will never need this kind of hackery: it only happens in a bit more complex macros. Here is an example:

```
\scratchcounter 1
\bggroup
\advance\scratchcounter 1
\egroup
\the\scratchcounter
```

```
\scratchcounter 1
\bggroup
\advance\scratchcounter 1
\expandafter
\egroup
\the\scratchcounter
```

The first one gives 1, while the second gives 2.

### 4.3 $\varepsilon$ -T<sub>E</sub>X primitives

In this engine a couple of extensions were added and later on pdfT<sub>E</sub>X added some more. We only discuss a few that relate to expansion. There is however a pitfall here. Before  $\varepsilon$ -T<sub>E</sub>X showed up, ConT<sub>E</sub>Xt already had a few mechanism that also related to expansion and it used some names for macros that clash with those in  $\varepsilon$ -T<sub>E</sub>X. This is why we will use the `\normal` prefix here to indicate the primitive.<sup>2</sup>

```
\def\MyMacroA{a}
\def\MyMacroB{b}
\normalprotected\def\MyMacroC{c}
\edef\MyMacroABC{\MyMacroA\MyMacroB\MyMacroC}
```

These macros have the following meanings:

```
macro:a
macro:b
```

<sup>2</sup> In the meantime we no longer have a low level `\protected` macro so one can use the primitive

```
protected macro:c
macro:ab\MyMacroC
```

In ConT<sub>E</sub>Xt you will use the `\unexpanded` prefix instead, because that one did something similar in older versions of ConT<sub>E</sub>Xt. As we were early adopters of  $\varepsilon$ -T<sub>E</sub>X, this later became a synonym to the  $\varepsilon$ -T<sub>E</sub>X primitive.

```
\def\MyMacroA{a}
\def\MyMacroB{b}
\normalprotected\def\MyMacroC{c}
\normalexpanded{\scratchtoks{\MyMacroA\MyMacroB\MyMacroC}}
```

Here the wrapper around the token register assignment will expand the three macros, unless they are protected, so its content becomes “ab\MyMacroC”. This saves either a lot of more complex `\expandafter` usage or the need to use an intermediate `\edef`. In ConT<sub>E</sub>Xt the `\expanded` macro does something simpler but it doesn't expand the first token as this is meant as a wrapper around a command, like:

```
\expanded{\chapter{...}} % a ConTeXt command
```

where we do want to expand the title but not the `\chapter` command (not that this would happen actually because `\chapter` is a protected command.)

The counterpart of `\normalexpanded` is `\normalunexpanded`, as in:

```
\def\MyMacroA{a}
\def\MyMacroB{b}
\normalprotected\def\MyMacroC{c}
\normalexpanded {\scratchtoks
  {\MyMacroA\normalunexpanded {\MyMacroB}\MyMacroC}}
```

The register now holds “a\MyMacroB\MyMacroC”: three tokens, one character token and two macro references.

Tokens can represent characters, primitives, macros or be special entities like starting math mode, beginning a group, assigning a dimension to a register, etc. Although you can never really get back to the original input, you can come pretty close, with:

```
\detokenize{this can $ be anything \bgroup}
```

This (when typeset monospaced) is: this can \$ be anything \bgroup. The detokenizer is like `\string` applied to each token in its argument. Compare this to:

```
\normalexpanded {
```

```
\normaldetokenize{10pt}
}
```

We get four tokens: 10pt.

```
\normalexpanded {
  \string 1\string 0\string p\string t
}
```

So that was the same operation: 10pt, but in both cases there is a subtle thing going on: characters have a catcode which distinguishes them. The parser needs to know what makes up a command name and normally that's only letters. The next snippet shows these catcodes:

```
\normalexpanded {
  \noexpand\the\catcode` \string 1 \noexpand\enspace
  \noexpand\the\catcode` \string 0 \noexpand\enspace
  \noexpand\the\catcode` \string p \noexpand\enspace
  \noexpand\the\catcode` \string t \noexpand
}
```

The result is “12 12 11 11”: two characters are marked as ‘letter’ and two fall in the ‘other’ category.

## 4.4 LuaTeX primitives

This engine adds a little to the expansion repertoire. First of all it offers a way to extend token lists registers:

```
\def\MyMacroA{a}
\def\MyMacroB{b}
\normalprotected\def\MyMacroC{b}
\scratchtoks{\MyMacroA\MyMacroB}
```

The result is: “\MyMacroA\MyMacroB”.

```
\toksapp\scratchtoks{\MyMacroA\MyMacroB}
```

We're now at: “\MyMacroA\MyMacroB\MyMacroA\MyMacroB\MyMacroA\MyMacroB”.

```
\etoksapp\scratchtoks{\MyMacroA\space\MyMacroB\space\MyMacroC}
```

The register has this content: “\MyMacroA\MyMacroB\MyMacroA\MyMacroB a b \MyMacroC a b \MyMacroC”, so the additional context got expanded in the process, except of course the protected macro \MyMacroC.

There is a bunch of these combinators: \toksapp and \tokspre for local appending and prepending, with global companions: \gtoksapp and \gtokspre, as well as expanding variant: \etoksapp, \etokspre, \xtoksapp and \xtokspre.

These are not beforehand more efficient than using intermediate expanded macros or token lists, simply because in the process  $\text{T}_{\text{E}}\text{X}$  has to create token lists too, but sometimes they're just more convenient to use. In  $\text{ConT}_{\text{E}}\text{Xt}$  we actually do benefit from these.

## 4.5 LuaMeta $\text{T}_{\text{E}}\text{X}$ primitives

We already saw that macro's can be defined protected which means that

```
\def\TestA{A}
\protected \def\TestB{B}
\edef\TestC{\TestA\TestB}
```

gives this:

```
\TestC : A\TestB
```

One way to get \TestB expanded is to prefix it with \expand:

```
\def\TestA{A}
\protected \def\TestB{B}
\edef\TestC{\TestA\TestB}
\edef\TestD{\TestA\expand\TestB}
```

We now get:

```
\TestC : A\TestB
\TestD : AB
```

There are however cases where one wishes this to happen automatically, but that will also make protected macros expand which will create havoc, like switching fonts.

```
\def\TestA{A}
\protected \def\TestB{B}
\semiprotected \def\TestC{C}
\edef\TestD{\TestA\TestB\TestC}
```

```

\edef\TestE{\normalexpanded{\TestA\TestB\TestC}}
\edef\TestF{\semiexpanded {\TestA\TestB\TestC}}

```

This time `\TestC` loses its protection:

```

\TestA : A
\TestB : B
\TestC : C
\TestD : A\TestB \TestC
\TestE : A\TestB \TestC
\TestF : A\TestB C

```

Actually adding `\fullyexpanded` would be trivial but it makes not much sense to add the overhead (at least not now). This feature is experimental anyway so it might go away when I see no real advantage from it.

When you store something in a macro or token register you always need to keep an eye on category codes. A dollar in the input is normally treated as math shift, a hash indicates a macro parameter or preamble entry. Characters like ‘A’ are letters but ‘[’ and ‘]’ are tagged as ‘other’. The  $\TeX$  scanner acts according to these codes. If you ever find yourself in a situation that changing catcodes is no option or cumbersome, you can do this:

```

\edef\TestOA{\expandtoken\othercatcode `A}
\edef\TestLA{\expandtoken\lettercatcode `A}

```

In both cases the meaning is A but in the first case it's not a letter but a character flagged as ‘other’.

A whole new category of commands has to do with so called local control. When  $\TeX$  scans and interprets the input, a process takes place that is called tokenizing: (sequences of) characters get a symbolic representation and travel through the system as tokens. Often they immediately get interpreted and are then discarded. But when for instance you define a macro they end up as a linked list of tokens in the macro body. We already saw that expansion plays a role. In most cases, unless  $\TeX$  is collecting tokens, the main action is dealt with in the so-called main loop. Something gets picked up from the input but can also be pushed back, for instance because of some lookahead that didn't result in an action. Quite some time is spent in pushing and popping from the so-called input stack.

When we are in Lua, we can pipe back into the engine but all is collected till we're back in  $\TeX$  where the collected result is pushed into the input. Because  $\TeX$  is a mix of programming and action there basically is only that main loop. There is no real way



to start a sub run in Lua and do all kind of things independent of the current one. This makes sense when you consider the mix: it would get too confusing.

However, in LuaTeX and even better in LuaMetaTeX, we can enter a sort of local state and this is called ‘local control’. When we are in local control a new main loop is entered and the current state is temporarily forgotten: we can for instance expand where one level up expansion was not done. It sounds complicated and indeed it is complicated so examples have to clarify it.

```
1 \setbox0\hbox to 10pt{2} \count0=3 \the\count0 \multiply\count0 by 4
```

This snippet of code is not that useful but illustrates what we're dealing with:

- The 1 gets typeset. So, characters like that are seen as text.
- The \setbox primitive triggers picking up a register number, then goes on scanning for a box specification and that itself will typeset a sequence of whatever until the group ends.
- The count primitive triggers scanning for a register number (or reference) and then scans for a number; the equal sign is optional.
- The the primitive injects some value into the current input stream and it does so by entering a new input level.
- The multiply primitive picks up a register specification and multiplies that by the next scanned number. The by is optional.

We now look at this snippet again but with an expansion context:

```
\def \TestA{1 \setbox0\hbox{2} \count0=3 \the\count0}
```

```
\edef\TestB{1 \setbox0\hbox{2} \count0=3 \the\count0}
```

These two macros have a slightly different body. Make sure you see the difference before reading on.

---

#### control sequence: TestA

---

594369	12	49	other char	1	U+00031	
594087	10	32	spacer			
593618	129	0	set box			setbox
405874	12	48	other char	0	U+00030	
594825	31	14	make box			hbox
593763	1	123	left brace			

594694	12	50	other char	2	U+00032	
334376	2	125	right brace			
481024	10	32	spacer			
588700	122	1	register			count
594828	12	48	other char	0	U+00030	
588800	12	61	other char	=	U+0003D	
594426	12	51	other char	3	U+00033	
594642	10	32	spacer			
594792	140	0	the			the
334384	122	1	register			count
215622	12	48	other char	0	U+00030	

---

**control sequence: TestB**


---

595071	12	49	other char	1	U+00031	
594766	10	32	spacer			
594900	129	0	set box			setbox
595073	12	48	other char	0	U+00030	
334342	31	14	make box			hbox
594388	1	123	left brace			
594830	12	50	other char	2	U+00032	
594676	2	125	right brace			
593919	10	32	spacer			
481064	122	1	register			count
593795	12	48	other char	0	U+00030	
481077	12	61	other char	=	U+0003D	
300399	12	51	other char	3	U+00033	
594576	10	32	spacer			
595126	12	49	other char	1	U+00031	

---

We now introduce a new primitive `\localcontrolled`:

```
\edef\TestB{1 \setbox0\hbox{2} \count0=3 \the\count0}
```

```
\edef\TestC{1 \setbox0\hbox{2} \localcontrolled{\count0=3} \the\count0}
```

Again, watch the subtle differences:

---

**control sequence: TestB**


---

594220	12	49	other char	1	U+00031	
593625	10	32	spacer			
481029	129	0	set box			setbox

595096	12	48	other char	0	U+00030	
580355	31	14	make box			hbox
592742	1	123	left brace			
595045	12	50	other char	2	U+00032	
593489	2	125	right brace			
595064	10	32	spacer			
594402	122	1	register			count
594132	12	48	other char	0	U+00030	
594304	12	61	other char	=	U+0003D	
481086	12	51	other char	3	U+00033	
593578	10	32	spacer			
594194	12	49	other char	1	U+00031	

---

**control sequence: TestC**


---

594206	12	49	other char	1	U+00031	
595071	10	32	spacer			
594766	129	0	set box			setbox
594900	12	48	other char	0	U+00030	
595073	31	14	make box			hbox
334342	1	123	left brace			
594388	12	50	other char	2	U+00032	
594830	2	125	right brace			
594676	10	32	spacer			
594576	10	32	spacer			
592690	12	51	other char	3	U+00033	

Another example:

```
\edef\TestB{1 \setbox0\hbox{2} \count0=3 \the\count0}
```

```
\edef\TestD{\localcontrolled{1 \setbox0\hbox{2} \count0=3 \the\count0}}
```

1 3 ← Watch how the results end up here!

---

**control sequence: TestB**


---

594240	12	49	other char	1	U+00031	
594581	10	32	spacer			
594294	129	0	set box			setbox
595105	12	48	other char	0	U+00030	
593535	31	14	make box			hbox
594565	1	123	left brace			

595089	12	50	other char	2	U+00032	
594696	2	125	right brace			
22591	10	32	spacer			
593856	122	1	register			count
594586	12	48	other char	0	U+00030	
593758	12	61	other char	=	U+0003D	
592743	12	51	other char	3	U+00033	
481079	10	32	spacer			
594266	12	51	other char	3	U+00033	

---

control sequence: TestD

---

<no tokens>

---

We can use this mechanism to define so called fully expandable macros:

```
\def\WidthOf#1%
  {\beginlocalcontrol
   \setbox0\hbox{#1}%
   \endlocalcontrol
   \wd0 }

\scratchdimen\WidthOf{The Rite Of Spring}
```

```
\the\scratchdimen
```

104.72021pt

When you want to add some grouping, it quickly can become less pretty:

```
\def\WidthOf#1%
  {\dimexpr
   \beginlocalcontrol
     \begingroup
       \setbox0\hbox{#1}%
       \expandafter
     \endgroup
   \expandafter
   \endlocalcontrol
   \the\wd0
  \relax}

\scratchdimen\WidthOf{The Rite Of Spring}
```

`\the\scratchdimen`

104.72021pt

A single token alternative is available too and its usage is like this:

```
\def\TestA{\scratchcounter=100 }
\edef\TestB{\localcontrol\TestA \the\scratchcounter}
\edef\TestC{\localcontrolled{\TestA} \the\scratchcounter}
```

The content of `\TestB` is '100' and of course the `\TestC` macro gives ' 100'.

We now move to the Lua end. Right from the start the way to get something into  $\TeX$  from Lua has been the print functions. But we can also go local (immediate). There are several methods:

- via a set token register
- via a defined macro
- via a string

Among the things to keep in mind are catcodes, scope and expansion (especially in when the result itself ends up in macros). We start with an example where we go via a token register:

```
\toks0={\setbox0\hbox{The Rite Of Spring}}
\toks2={\setbox0\hbox{The Rite Of Spring!}}
```

`\startluacode`

```
tex.runlocal(0) context("[1: %p]",tex.box[0].width)
tex.runlocal(2) context("[2: %p]",tex.box[0].width)
```

`\stopluacode`

[1: 104.72021pt][2: 109.14062pt]

We can also use a macro:

```
\def\TestA{\setbox0\hbox{The Rite Of Spring}}
\def\TestB{\setbox0\hbox{The Rite Of Spring!}}
```

`\startluacode`

```
tex.runlocal("TestA") context("[3: %p]",tex.box[0].width)
tex.runlocal("TestB") context("[4: %p]",tex.box[0].width)
```

`\stopluacode`

[3: 104.72021pt][4: 109.14062pt]

A third variant is more direct and uses a (Lua) string:

```
\startluacode
tex.runstring([[ \setbox0\hbox{The Rite Of Spring} ]])

context("[5: %p]", tex.box[0].width)

tex.runstring([[ \setbox0\hbox{The Rite Of Spring!} ]])

context("[6: %p]", tex.box[0].width)
\stopluacode
```

[5: 104.72021pt][6: 109.14062pt]

A bit more high level:

```
context.runstring([[ \setbox0\hbox{(Here \bf 1.2345)} ]])
context.runstring([[ \setbox0\hbox{(Here \bf   %.3f)} ]], 1.2345)
```

Before we had runstring this was the way to do it when staying in Lua was needed:

```
\startluacode
token.setmacro("TestX", [[ \setbox0\hbox{The Rite Of Spring} ]])
tex.runlocal("TestX")
context("[7: %p]", tex.box[0].width)
\stopluacode
```

[7: 104.72021pt]

```
\startluacode
tex.scantoks(0, tex.ctxcatcodes, [[ \setbox0\hbox{The Rite Of Spring!} ]])
tex.runlocal(0)
context("[8: %p]", tex.box[0].width)
\stopluacode
```

[8: 109.14062pt]

The order of flushing matters because as soon as something is not stored in a token list or macro body,  $\TeX$  will typeset it. And as said, a lot of this relates to pushing stuff into the input which is stacked. Compare:

```
\startluacode
```

```
context("[HERE 1]")
context("[HERE 2]")
\stopluacode
```

[HERE 1][HERE 2]

with this:

```
\startluacode
tex.pushlocal() context("[HERE 1]") tex.poplocal()
tex.pushlocal() context("[HERE 2]") tex.poplocal()
\stopluacode
```

[HERE 1][HERE 2]

You can expand a macro at the Lua end with `token.expandmacro` which has a peculiar interface. The first argument has to be a string (the name of a macro) or a userdata (a valid macro token). This macro can be fed with parameters by passing more arguments:

string	serialized to tokens
true	wrap the next string in curly braces
table	each entry will become an argument wrapped in braces
token	inject the token directly
number	change control to the given catcode table

There are more scanner related primitives, like the  $\varepsilon$ -TeX primitive `\detokenize`:

```
[\detokenize {test \relax}]
```

This gives: `[test \relax ]`. In LuaMetaTeX we also have complementary primitive(s):

```
[\tokenized catcodetable \vrbcatcodes {test {\bf test} test}]
[\tokenized {test {\bf test} test}]
[\retokenized \vrbcatcodes {test {\bf test} test}]
```

The `\tokenized` takes an optional keyword and the examples above give: `[test {\bf test} test [test test test] [test {\bf test} test]`. The LuaTeX primitive `\scantextokens` which is a variant of  $\varepsilon$ -TeX's `\scantokens` operates under the current catcode regime (the last one honors `\everyeof`). The difference with `\tokenized` is that this one first serializes the given token list (just like `\detokenize`).<sup>3</sup>

<sup>3</sup> The `\scan *tokens` primitives now share the same helpers as Lua, but they should behave the same as in LuaTeX.

With `\retokenized` the catcode table index is mandatory (it saves a bit of scanning and is easier on intermixed `\expandafter` usage. There often are several ways to accomplish the same:

```
\def\MyTitle{test {\bf test} test}
\detokenize          \expandafter{\MyTitle}: 0.46\crlf
\meaningless        \MyTitle : 0.47\crlf
\retokenized         \notcatcodes{\MyTitle}: 0.87\crlf
\tokenized   catcodetable \notcatcodes{\MyTitle}: 0.93\crlf
```

```
test {\bf test} test: 0.46
test {\bf test} test: 0.47
test {\bf test} test: 0.87
test {\bf test} test: 0.93
```

Here the numbers show the relative performance of these methods. The `\detokenize` and `\meaningless` win because they already know that a verbose serialization is needed. The last two first serialize and then reinterpret the resulting token list using the given catcode regime. The last one is slowest because it has to scan the keyword.

There is however a pitfall here:

```
\def\MyText {test}
\def\MyTitle{test \MyText\space test}
\detokenize          \expandafter{\MyTitle}\crlf
\meaningless        \MyTitle \crlf
\retokenized         \notcatcodes{\MyTitle}\crlf
\tokenized   catcodetable \notcatcodes{\MyTitle}\crlf
```

The outcome is different now because we have an expandable embedded macro call. The fact that we expand in the last two primitives is also the reason why they are ‘slower’.

```
test \MyText \space test
test \MyText \space test
test test test
test test test
```

To complete this picture, we show a variant than combines much of what has been introduced in this section:



```

\semiprotected\def\MyTextA {test}
\def\MyTextB {test}
\def\MyTitle{test \MyTextA\space \MyTextB\space test}
\detokenize           \expandafter{\MyTitle}\crlf
\meaningless         \MyTitle \crlf
\retokenized          \notcatcodes{\MyTitle}\crlf
\retokenized          \notcatcodes{\semiexpanded{\MyTitle}}\crlf
\tokenized   catcodetable \notcatcodes{\MyTitle}\crlf
\tokenized   catcodetable \notcatcodes{\semiexpanded{\MyTitle}}

```

This time compare the last four lines:

```

test \MyTextA \space \MyTextB \space test
test \MyTextA \space \MyTextB \space test
test \MyTextA test test
test test test test
test \MyTextA test test
test test test test

```

Of course the question remains to what extent we need this and eventually will apply in ConT<sub>E</sub>Xt. The `\detokenize` is used already. History shows that eventually there is a use for everything and given the way LuaMetaT<sub>E</sub>X is structured it was not that hard to provide the alternatives without sacrificing performance or bloating the source.

## 4.6 Dirty tricks

When I was updating this manual Hans vd Meer and I had some discussions about expansion and tokenization related issues when combining of xml processing with T<sub>E</sub>X macros where he did some manipulations in Lua. In these mixed cases you can run into catcode related problems because in xml you want for instance a `#` to be a hash mark (other character) and not a parameter identifier. Normally this is handled well in ConT<sub>E</sub>Xt but of course there are complex cases where you need to adapt.

Say that you want to compare two strings (officially we should say token lists) with mixed catcodes. Let's also assume that you want to use the normal `\if` construct (which was part of the discussion). We start with defining a test set. The reason that we present this example here is that we use commands discussed in previous sections:

```

\def\abc{abc}
\semiprotected \def\xyz{xyz}
\edef\pqr{\expandtoken\notcatcodes`p%

```

```
\expandtoken\notcatcodes`q%
\expandtoken\notcatcodes`r}
```

```
1: \ifcondition\similartokens{abc} {def}YES\else NOP\fi (NOP) \quad
2: \ifcondition\similartokens{abc}{\abc}YES\else NOP\fi (YES)

3: \ifcondition\similartokens{xyz} {pqr}YES\else NOP\fi (NOP) \quad
4: \ifcondition\similartokens{xyz}{\xyz}YES\else NOP\fi (YES)

5: \ifcondition\similartokens{pqr} {pqr}YES\else NOP\fi (YES) \quad
6: \ifcondition\similartokens{pqr}{\pqr}YES\else NOP\fi (YES)
```

So, we have a mix of expandable and semi expandable macros, and also a mix of cat-codes. A naive approach would be:

```
\permanent\protected\def\similartokens#1#2%
  {\iftok{#1}{#2}}
```

but that will fail on some cases:

```
1: NOP(NOP)   2: YES(YES)
3: NOP(NOP)   4: NOP(YES)
5: YES(YES)   6: NOP(YES)
```

So how about:

```
\permanent\protected\def\similartokens#1#2%
  {\iftok{\detokenize{#1}}{\detokenize{#2}}}
```

That one is even worse:

```
1: NOP(NOP)   2: NOP(YES)
3: NOP(NOP)   4: NOP(YES)
5: YES(YES)   6: NOP(YES)
```

We need to expand so we end up with this:

```
\permanent\protected\def\similartokens#1#2%
  {\normalexpanded{\noexpand\iftok
    {\noexpand\detokenize{#1}}
    {\noexpand\detokenize{#2}}}}
```

Better:

```
1: NOP(NOP)   2: YES(YES)
```

3: NOP(NOP)    4: NOP(YES)  
 5: YES(YES)    6: YES(YES)

But that will still not deal with the mildly protected macro so in the end we have:

```
\permanent\protected\def\similartokens#1#2%
  {\semiexpanded{\noexpand\iftok
    {\noexpand\detokenize{#1}}
    {\noexpand\detokenize{#2}}}}
```

Now we're good:

1: NOP(NOP)    2: YES(YES)  
 3: NOP(NOP)    4: YES(YES)  
 5: YES(YES)    6: YES(YES)

Finally we wrap this one in the usual `\doifelse...` macro:

```
\permanent\protected\def\doifelsesimilartokens#1#2%
  {\ifcondition\similartokens{#1}{#2}%
    \expandafter\firstoftwoarguments
  \else
    \expandafter\secondoftwoarguments
  \fi}
```

so that we can do:

```
\doifelsesimilartokens{pqr}{\pqr}{YES}{NOP}
```

A companion macro of this is `\wipetoken` but for that one you need to look into the source.

## 4.6 Colofon

Author	Hans Hagen
ConT <sub>E</sub> Xt	2025.07.04 21:26
LuaMetaT <sub>E</sub> X	2.11.07   20250705
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## 5 Registers

low level

TEX

registers

## Contents

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### 5.1 Preamble

Registers are sets of variables that are accessed by index and as such resemble registers in a processing unit. You can store a quantity in a register, retrieve it, and also manipulate it.

There is hardly any need to use them in ConT<sub>E</sub>Xt so we keep it simple.

### 5.2 T<sub>E</sub>X primitives

There are several categories:

- Integers (int): in order to be portable (at the time it surfaced) there are only integers and no floats. The only place where T<sub>E</sub>X uses floats internally is when glue gets effective which happens in the backend.
- Dimensions (dimen): internally these are just integers but when they are entered they are sliced into two parts so that we have a fractional part. The internal representation is called a scaled point.
- Glue (skip): these are dimensions with a few additional properties: stretch and shrink. Being a compound entity they are stored differently and thereby a bit less efficient than numbers and dimensions.
- Math glue (muskip): this is the same as glue but with a unit that adapts to the current math style properties. It's best to think about them as being relative measures.
- Token lists (toks): these contain a list of tokens coming from the input or coming from a place where they already have been converted.

The original T<sub>E</sub>X engine had 256 entries per set. The first ten of each set are normally reserved for scratch purposes: the even ones for local use, and the odd ones for global

usage. On top of that macro packages can reserve some for its own use. It was quite easy to reach the maximum but there were tricks around that. This limitation is no longer present in the variants in use today.

Let's set a few dimension registers:

```
\dimen 0 = 10 pt
\dimen2=10pt
\dimen4 10pt
\scratchdimen 10pt
```

We can serialize them with:

```
\the \dimen0
\number \dimen2
\meaning\dimen4
\meaning\scratchdimen
```

The results of these operations are:

```
10.0pt
655360
\dimen4
constant dimension 10.0pt
```

The last two is not really useful but it is what you see when tracing options are set. Here `\scratchdimen` is a shortcut for a register. This is *not* a macro but a defined register. The low level `\dimendef` is used for this but in a macro package you should not use that one but the higher level `\newdimen` macro that uses it.

```
\newdimen\MyDimenA
\def \MyDimenB{\dimen999}
\dimendef\MyDimenC998

\meaning\MyDimenA
\meaning\MyDimenB
\meaning\MyDimenC
```

Watch the difference:

```
\dimen269
macro:\dimen 999
\dimen998
```

The first definition uses a yet free register so you won't get a clash. The second one is just a shortcut using a macro and the third one too but again direct shortcut. Try to imagine how the second line gets interpreted:

```
\MyDimenA10pt \MyDimenA10.5pt
\MyDimenB10pt \MyDimenB10.5pt
\MyDimenC10pt \MyDimenC10.5pt
```

Also try to imagine what messing around with `\MyDimenC` will do when we also have defined a few hundred extra dimensions with `\newdimen`.

In the case of dimensions the `\number` primitive will make the register serialize as scaled points without unit `sp`.

Next we see some of the other registers being assigned:

```
\count 0 = 100
\skip 0 = 10pt plus 3pt minus 2pt
\skip 0 = 10pt plus 1fill
\muskip 0 = 10mu plus 3mu minus 2mu
\muskip 0 = 10mu minus 1 fil
\toks 0 = {hundred}
```

When a number is expected, you can use for instance this:

```
\scratchcounter\scratchcounterone
```

Here we use a few predefined scratch registers. You can also do this:

```
\scratchcounter\numexpr\scratchcounterone+\scratchcountertwo\relax
```

There are some quantities that also qualify as number:

```
\chardef\MyChar=123 % refers to character 123 (if present)
\scratchcounter\MyChar
```

In the past using `\chardef` was a way to get around the limited number of registers, but it still had (in traditional  $\TeX$ ) a limitation: you could not go beyond 255. The `\mathchardef` could go higher as it also encodes a family number and class. This limitation has been lifted in  $\text{Lua}\TeX$ .

A character itself can also be interpreted as number, in which case it has to be prefixed with a reverse quote: ```, so:



```
\scratchcounter\numexpr`0+5\relax
\char\scratchcounter
```

produces “5” because the ``0` expands into the (ascii and utf8) slot 48 which represents the character zero. In this case the next makes more sense:

```
\char\numexpr`0+5\relax
```

If you want to know more about all these quantities, “`TEX` By Topic” provides a good summary of what `TEX` has to offer, and there is no need to repeat it here.

### 5.3 $\varepsilon$ -`TEX` primitives

Apart from the ability to use expressions, the contribution to registers that  $\varepsilon$ -`TEX` brought was that suddenly we could use upto 65K of them, which is more than enough. The extra registers were not as efficient as the first 256 because they were stored in the hash table, but that was not really a problem. In Omega and later Lua`TEX` regular arrays were used, at the cost of more memory which in the meantime has become cheap. As `ConTEXt` moved to  $\varepsilon$ -`TEX` rather early its users never had to worry about it.

### 5.4 Lua`TEX` primitives

The Lua`TEX` engine introduced attributes. These are numeric properties that are bound to the nodes that are the result of typesetting operations. They are basically like integer registers but when set their values get bound and when unset they are kind of invisible.

- Attribute (attribute): a numeric property that when set becomes part of the current attribute list that gets assigned to nodes.

Attributes can be used to communicate properties to Lua callbacks. There are several functions available for setting them and querying them.

```
\attribute999 = 123
```

Using attributes this way is dangerous (of course I can only speak for `ConTEXt`) because an attribute value might trigger some action in a callback that gives unwanted side effects. For convenience `ConTEXt` provides:

```
\newattribute\MyAttribute
```

Which currently defines `\MyAttribute` as constant integer 1026 and is meant to be used as:<sup>4</sup>

```
\attribute\MyAttribute = 123
```

Just be aware that defining attributes can have an impact on performance. As you cannot access them at the `TEX` end you seldom need them. If you do you can better use the proper more high level definers (not discussed here).

## 5.5 LuaMetaT<sub>E</sub>X primitives

The fact that scanning stops at a non-number or `\relax` can be sort of unpredictable which is why in LuaMetaT<sub>E</sub>X we also support the following variant:

```
\scratchdimen\dimexpr 10pt + 3pt \relax
\scratchdimen\dimexpr {10pt + 3pt}
```

At the cost of one more token braces can be used as boundaries instead of the single `\relax` boundary.

An important property of registers is that they can be accessed by a number. This has big consequences for the implementation: they are part of the big memory store and consume dedicated ranges. If we had only named access T<sub>E</sub>X's memory layout could be a bit leaner. In principle we could make the number of registers smaller because any limit on the amount at some point can be an obstacle. It is for that reason that we also have name-only variants:

```
\dimensiondef \MyDimenA 12pt
\integerdef \MyIntegerA 12
\gluespecdef \MyGlueA 12pt + 3pt minus 4pt
\mugluespecdef\MyMuA 12mu + 3mu minus 4mu
```

These are as efficient but not accessible by number but they behave like registers which means that you (can) use `\the`, `\advance`, `\multiply` and `\divide` with them.<sup>5</sup> In case you wonder why there is no alternative for `\toksdef`, there actually are multiple: they are called macros.

*todo: expressions*

<sup>4</sup> The low level `\attributedef` command is rather useless in the perspective of ConT<sub>E</sub>Xt.

<sup>5</sup> There are also the slightly more efficient `\advanceby`, `\multiplyby` and `\divideby` that don't check for the by keyword.

## 5.6 Units

The LuaMetaT<sub>E</sub>X engine supports the following units. The first batch is constant with hard coded fine tuned values. The second set is related to the current font. The last group is kind of special, the es is a replacement for the in and has a little sister in ts. The dk is dedicated to the master and makes a nice offset for so called T<sub>E</sub>X pages that we use for demos.

---

pt	1.0	point
bp	1.00374	big point (aka postscript point)
in	72.26999	inch
cm	28.45274	centimeter
mm	2.84526	milimeter
dd	1.07	didot
cc	12.8401	cicero
pc	12.0	pica
sp	0.00002	scaled points
px	0.00002	pixel

---

ex	5.70947	ex height
em	11.0	em width
mu	1.0	math unit

---

ts	7.11317	tove
es	71.13177	edith
eu	71.13177	european unit
dk	6.43985	knuth

---

The fi[lll] unit is not really a unit but a multiplier for infinite stretch and shrink; original T<sub>E</sub>X doesn't have the simple fi.

In addition to these we can have many more. In principle a user can define additional ones but there's always a danger of clashing. For users we reserve the units starting with an u. Here is how you define your own, we show three variants:

```

\newdimension \FooA   \FooA 1.23pt
\newdimen    \FooB   \FooB 12.3pt
\protected\def\FooC  {\the\dimexpr\FooA +\FooB\relax}

\pushoverloadmode % just in case
  \newuserunit\FooA ua
  \newuserunit\FooB ub
  \newuserunit\FooC uc

```

**\popoverloadmode**

And this is how they show up:

2.45999pt 24.6pt 27.06pt

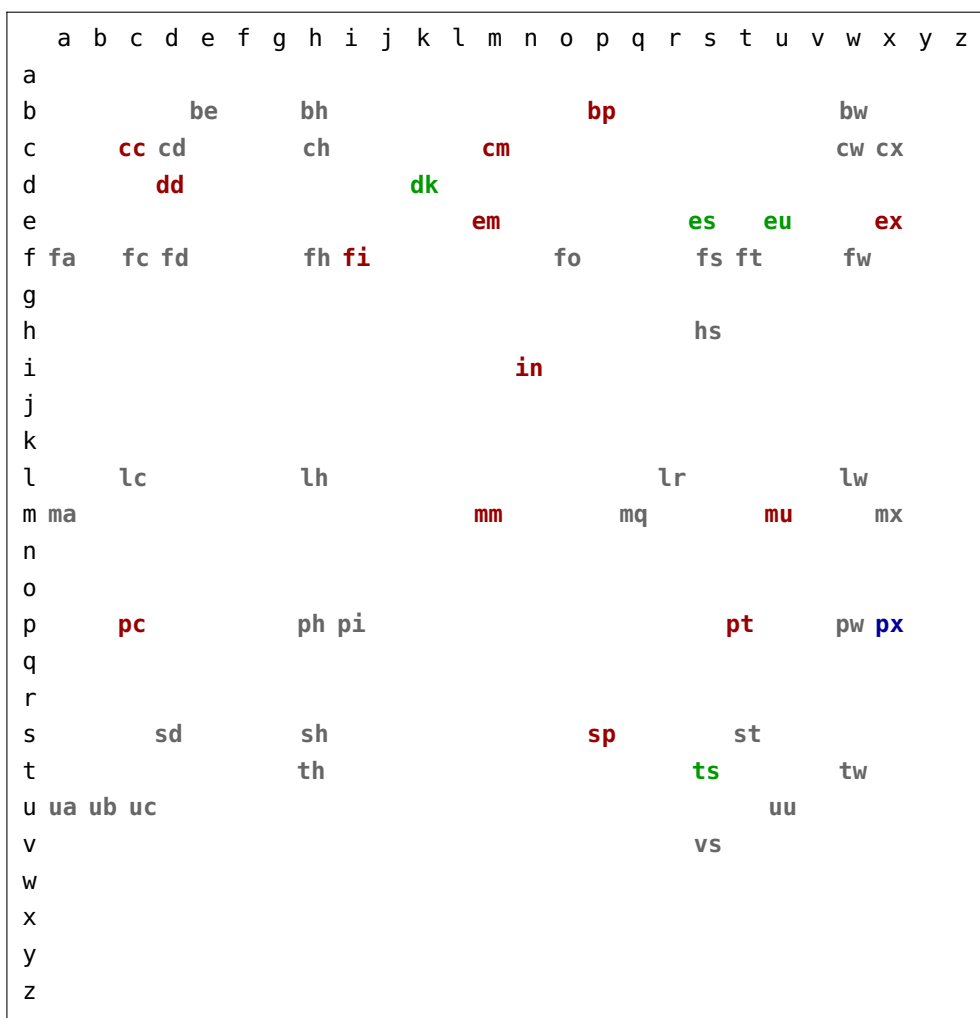
with

```
\the\dimexpr 2 ua \relax\quad
\the\dimexpr 2 ub \relax\quad
\the\dimexpr 2 uc \relax
```

The following additional units are predefined (reserved). The values are in points and some depend on the current layout and document font.

pi	3.14159	$\pi$ for Mikael
ft	867.23999	foot for Alan
fs	11.0	(global body) font size
tw	483.69687	(layout) text width
th	645.88272	(layout) text height
hs	483.69687	(current) hsize
vs	645.88272	(current) vsize
cd	0.0	(when set) column distance
cw	483.69687	(when set) column width
cx	236.34843	combination cell width
uu	28.45274	user unit (MetaFun)
fw	0.0	framed width
fh	0.0	framed height
fo	0.0	framed offset
lw	0.4	line width
sh	11.51031	strut height
sd	4.47621	strut depth
st	15.98653	strut total
ch	6.99854	width of zero (css)
fa	8.35742	font ascender
fd	1.71338	font descender
fc	8.01904	font cap height

Here is an example of usage:



**Figure 5.1** A map of available units

```

\startcombination[nx=4,ny=1]
  {\ruledhbox to 1cx{\strut one}} {1}
  {\ruledhbox to 1cx{\strut two}} {2}
  {\ruledhbox to 1cx{\strut three}} {3}
  {\ruledhbox to 1cx{\strut four}} {4}
\stopcombination

```

one.....	two.....	three.....	four.....
1	2	3	4

The uu can be set by users using the \uunit dimension variable. The default value is 1cm. Its current value is also known at the MetaPost end, as demonstrated in figure 5.2.

```

\startcombination[nx=2,ny=1]

```

```

\startcontent
  \uunit=1cm
  \framed[offset=1uu]
    \bgroup
      \startMPcode
        fill fullcircle scaled 3uu withcolor "darkred" ;
        fill fullcircle scaled 2cm withcolor "darkgreen" ;
      \stopMPcode
    \egroup
\stopcontent
\startcaption
  \type {\uunit = 1cm}
\stopcaption
\startcontent
  \uunit=1cx
  \framed[offset=.1uu]
    \bgroup
      \startMPcode
        fill fullcircle scaled .5uu withcolor "darkblue" ;
        fill fullcircle scaled 2cm withcolor "darkyellow" ;
      \stopMPcode
    \egroup
\stopcontent
\startcaption
  \type {\uunit = 1cx}
\stopcaption
\stopcombination

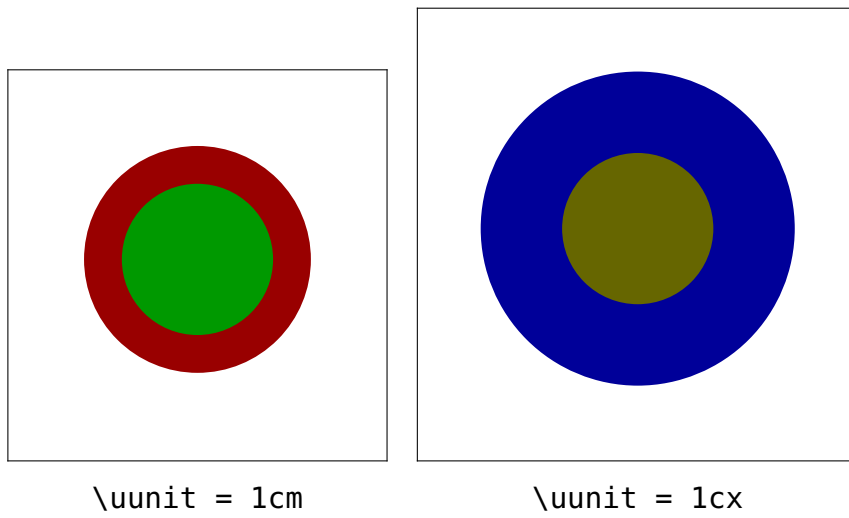
```

There is one catch here. If you use your own uu as numeric, you might need this:

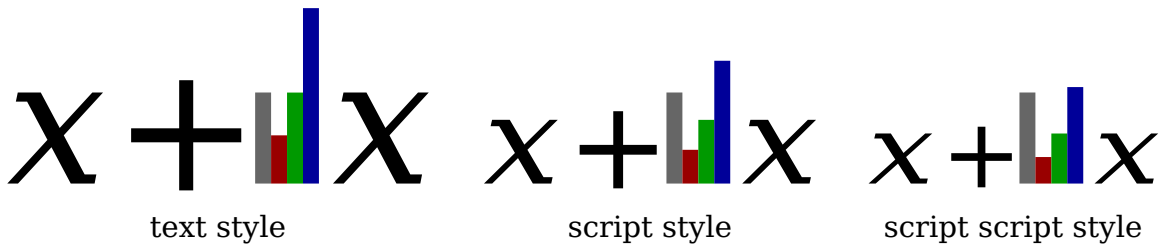
```
save uu ; numeric uu ; uu := 1cm ;
```

That is: make sure the meaning is restored afterwards and explicitly declare the variable. But this is good practice anyway when you generate multiple graphics using the same MetaPost instance.

There are a few units not mentioned yet and those concern math, where we need to adapt to the current style.



**Figure 5.2** Shared user units in  $\text{T}_{\text{E}}\text{X}$  and MetaFun.



The bars show  $1\text{ex}$ ,  $1\text{ma}$  (axis),  $1\text{mx}$  (ex-height) and  $1\text{mq}$  (em-width or quad). The last three adapt themselves to the style. Often the  $\text{mx}$  makes more sense than  $\text{ex}$ .

## 5.6 Colofon

Author	Hans Hagen
ConT <sub>E</sub> Xt	2025.07.04 21:26
LuaMetaT <sub>E</sub> X	2.11.07   20250705
Support	<a href="http://www.pragma-ade.com">www.pragma-ade.com</a> <a href="http://contextgarden.net">contextgarden.net</a> <a href="mailto:ntg-context@ntg.nl">ntg-context@ntg.nl</a>



## **6 Macros**

low level

TEX

macros

## Contents

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## 6.1 Preamble

This chapter overlaps with other chapters but brings together some extensions to the macro definition and expansion parts. As these mechanisms were stepwise extended, the other chapters describe intermediate steps in the development.

Now, in spite of the extensions discussed here the main idea is still that we have  $\text{T}_{\text{E}}\text{X}$  act like before. We keep the charm of the macro language but these additions make for easier definitions, but (at least initially) none that could not be done before using more code.

## 6.2 Definitions

A macro definition normally looks like like this:<sup>6</sup>

```
\def\macro#1#2%
  {\dontleavehmode\hbox to 6em{\vl\type{#1}\vl\type{#2}\vl\hss}}
```

Such a macro can be used as:

```
\macro {1}{2}
\macro {1} {2} middle space gobbled
\macro 1 {2} middle space gobbled
\macro {1} 2 middle space gobbled
```

<sup>6</sup> The `\dontleavehmode` command make the examples stay on one line.

```
\macro 1 2      middle space gobbled
```

We show the result with some comments about how spaces are handled:

```
12|
12|      middle space gobbled
12|      middle space gobbled
12|      middle space gobbled
12|      middle space gobbled
```

A definition with delimited parameters looks like this:

```
\def\macro[#1]%
  {\dontleavehmode\hbox to 6em{\v\type{#1}\v\hss}}
```

When we use this we get:

```
\macro [1]
\macro [ 1]      leading space kept
\macro [1 ]      trailing space kept
\macro [ 1 ]      both spaces kept
```

Again, watch the handling of spaces:

```
1|
1|      leading space kept
1|      trailing space kept
1|      both spaces kept
```

Just for the record we show a combination:

```
\def\macro[#1]#2%
  {\dontleavehmode\hbox to 6em{\v\type{#1}\v\type{#2}\v\hss}}
```

With this:

```
\macro [1]{2}
\macro [1] {2}
\macro [1] 2
```

we can again see the spaces go away:

```
12|
12|
```

```
|12|
```

A definition with two separately delimited parameters is given next:

```
\def\macro[#1#2]%
  {\dontleavehmode\hbox to 6em{\vl\type{#1}\vl\type{#2}\vl\hss}}
```

When used:

```
\macro [12]
\macro [ 12]    leading space gobbled
\macro [12 ]   trailing space kept
\macro [ 12 ]  leading space gobbled, trailing space kept
\macro [1 2]   middle space kept
\macro [ 1 2 ] leading space gobbled, middle and trailing space kept
```

We get ourselves:

```
|12|
|12|    leading space gobbled
|12|    trailing space kept
|12|    leading space gobbled, trailing space kept
|1 2|   middle space kept
|1 2|   leading space gobbled, middle and trailing space kept
```

These examples demonstrate that the engine does some magic with spaces before (and therefore also between multiple) parameters.

We will now go a bit beyond what traditional T<sub>E</sub>X engines do and enter the domain of LuaMetaT<sub>E</sub>X specific parameter specifiers. We start with one that deals with this hard coded space behavior:

```
\def\macro[#^#^]%
  {\dontleavehmode\hbox to 6em{\vl\type{#1}\vl\type{#2}\vl\hss}}
```

The #<sup>^</sup> specifier will count the parameter, so here we expect again two arguments but the space is kept when parsing for them.

```
\macro [12]
\macro [ 12]
\macro [12 ]
\macro [ 12 ]
\macro [1 2]
```

```
\macro [ 1 2 ]
```

Now keep in mind that we could deal well with all kind of parameter handling in Con-TeXt for decades, so this is not really something we missed, but it complements the to be discussed other ones and it makes sense to have that level of control. Also, availability triggers usage. Nevertheless, some day the #^ specifier will come in handy.

```
|12|
| 12|
|12 |
| 12 |
|1 2|
|1 2 |
```

We now come back to an earlier example:

```
\def\macro[#1]%
  {\dontleavehmode\hbox spread 1em{\vl\type{#1}\vl\hss}}
```

When we use this we see that the braces in the second call are removed:

```
\macro [1]
\macro [{1}]
```

```
|1| |1|
```

This can be prohibited by the #+ specifier, as in:

```
\def\macro[#+]%
  {\dontleavehmode\hbox spread 1em{\vl\type{#1}\vl\hss}}
```

As we see, the braces are kept:

```
\macro [1]
\macro [{1}]
```

Again, we could easily get around that (for sure intended) side effect but it just makes nicer code when we have a feature like this.

```
|1| |{1}|
```

Sometimes you want to grab an argument but are not interested in the results. For this we have two specifiers: one that just ignores the argument, and another one that keeps counting but discards it, i.e. the related parameter is empty.

```
\def\macro[#1][#0][#3][#-][#4]%
  {\dontleavehmode\hbox spread 1em
   {\v\type{#1}\v\type{#2}\v\type{#3}\v\type{#4}\v\hss}}
```

The second argument is empty and the fourth argument is simply ignored which is why we need #4 for the fifth entry.

```
\macro [1][2][3][4][5]
```

Here is proof that it works:

```
|1|3|5|
```

The reasoning behind dropping arguments is that for some cases we get around the nine argument limitation, but more important is that we don't construct token lists that are not used, which is more memory (and maybe even cpu cache) friendly.

Spaces are always kind of special in T<sub>E</sub>X, so it will be no surprise that we have another specifier that relates to spaces.

```
\def\macro[#1]#*[#2]%
  {\dontleavehmode\hbox spread 1em{\v\type{#1}\v\type{#2}\v\hss}}
```

This permits usage like the following:

```
\macro [1][2]
\macro [1] [2]
```

```
|1|2| |1|2|
```

Without the optional 'grab spaces' specifier the second line would possibly throw an error. This because T<sub>E</sub>X then tries to match ][ so the ] [ in the input is simply added to the first argument and the next occurrence of ][ will be used. That one can be someplace further in your source and if not T<sub>E</sub>X complains about a premature end of file. But, with the #\* option it works out okay (unless of course you don't have that second argument [2]).

Now, you might wonder if there is a way to deal with that second delimited argument being optional and of course that can be programmed quite well in traditional macro code. In fact, ConT<sub>E</sub>Xt does that a lot because it is set up as a parameter driven system with optional arguments. That subsystem has been optimized to the max over years and it works quite well and performance wise there is very little to gain. However, as soon as you enable tracing you end up in an avalanche of expansions and that is no fun.

This time the solution is not in some special specifier but in the way a macro gets defined.

```
\tolerant\def\macro[#1]#* [#2]%
  {\dontleavehmode\hbox spread 1em{\vl\type{#1}\vl\type{#2}\vl\hss}}
```

The magic `\tolerant` prefix with delimited arguments and just quits when there is no match. So, this is acceptable:

```
\macro [1][2]
\macro [1] [2]
\macro [1]
\macro
```

```
|12| |12| |1| |
```

We can check how many arguments have been processed with a dedicated conditional:

```
\tolerant\def\macro[#1]#* [#2]%
  {\ifarguments 0\or 1\or 2\or ?\fi: \vl\type{#1}\vl\type{#2}\vl}
```

We use this test:

```
\macro [1][2] \macro [1] [2] \macro [1] \macro
```

The result is: 2: |12| 2: |12| 1: |1| 0: | which is what we expect because we flush inline and there is no change of mode. When the following definition is used in display mode, the leading `n=` can for instance start a new paragraph and when code in `\everypar` you can loose the right number when macros get expanded before the `n` gets injected.

```
\tolerant\def\macro[#1]#* [#2]%
  {n=\ifarguments 0\or 1\or 2\or ?\fi: \vl\type{#1}\vl\type{#2}\vl}
```

In addition to the `\ifarguments` test primitive there is also a related internal counter `\lastarguments` set that you can consult, so the `\ifarguments` is actually just a shortcut for `\ifcase\lastarguments`.

We now continue with the argument specifiers and the next two relate to this optional grabbing. Consider the next definition:

```
\tolerant\def\macro#1#*#2%
  {\dontleavehmode\hbox spread 1em{\vl\type{#1}\vl\type{#2}\vl\hss}}
```

With this test:



```
\macro {1} {2}
\macro {1}
\macro
```

We get:

```
|12| |1\macro|
```

This is okay because the last `\macro` is a valid (single token) argument. But, we can make the braces mandate:

```
\tolerant\def\macro#=#*#=#%
  {\dontleavehmode\hbox spread 1em{\vl\type{#1}\vl\type{#2}\vl\hss}}
```

Here the `#=` forces a check for braces, so:

```
\macro {1} {2}
\macro {1}
\macro
```

gives this:

```
|12| |1| |
```

However, we do loose these braces and sometimes you don't want that. Of course when you pass the results downstream to another macro you can always add them, but it was cheap to add a related specifier:

```
\tolerant\def\macro#_#*#_#_%
  {\dontleavehmode\hbox spread 1em{\vl\type{#1}\vl\type{#2}\vl\hss}}
```

Again, the magic `\tolerant` prefix works will quit scanning when there is no match. So:

```
\macro {1} {2}
\macro {1}
\macro
```

leads to:

```
|{1}|{2}| |{1}| |
```

When you're tolerant it can be that you still want to pick up some argument later on. This is why we have a continuation option.

```

\tolerant\def\foo      [#1]*[#2]#:#3{!#1!#2!#3!}
\tolerant\def\oof[#1]*[#2]#:(#3)#:#4{!#1!#2!#3!#4!}
\tolerant\def\ofo      [#1]#:(#2)#:#3{!#1!#2!#3!}

```

Hopefully the next example demonstrates how it works:

```

\foo{3} \foo[1]{3} \foo[1][2]{3}
\oof{4} \oof[1]{4} \oof[1][2]{4}
\oof[1][2](3){4} \oof[1](3){4} \oof(3){4}
\ofo{3} \ofo[1]{3}
\ofo[1](2){3} \ofo(2){3}

```

As you can see we can have multiple continuations using the #: directive:

```

!!!3! !1!!!3! !1!2!3!
!!!!4! !1!!!!4! !1!2!!!4!
!1!2!3!4! !1!!!3!4! !!!3!4!
!!!3! !1!!!3!
!1!2!3! !2!3!

```

The last specifier doesn't work well with the \ifarguments state because we no longer know what arguments were skipped. This is why we have another test for arguments. A zero value means that the next token is not a parameter reference, a value of one means that a parameter has been set and a value of two signals an empty parameter. So, it reports the state of the given parameter as a kind of \ifcase.

```

\def\foo#1#2{ [\ifparameter#1\or(ONE)\fi\ifparameter#2\or(TWO)\fi] }

```

Of course the test has to be followed by a valid parameter specifier:

```

\foo{1}{2} \foo{1}{} \foo{}{2} \foo{}{}

```

The previous code gives this:

```

[(ONE)(TWO)] [(ONE)] [(TWO)] []

```

A combination check \ifparameters, again a case, matches the first parameter that has a value set.

We could add plenty of specifiers but we need to keep in mind that we're not talking of an expression scanner. We need to keep performance in mind, so nesting and backtracking are no option. We also have a limited set of useable single characters, but here's one that uses a symbol that we had left:

```
\def\startfoo[#/#]\stopfoo{ [#1](#2) }
```

The slash directive removes leading and trailing so called spacers as well as tokens that represent a paragraph end:

```
\startfoo [x ] x \stopfoo
\startfoo [ x ] x \stopfoo
\startfoo [ x] x \stopfoo
\startfoo [ x] \par x \par \par \stopfoo
```

So we get this:

```
[x](x) [x](x) [x](x) [x](x)
```

The next directive, the quitter #;, is demonstrated with an example. When no match has occurred, scanning picks up after this signal, otherwise we just quit.

```
\tolerant\def\foo[#1]#;( #2){/#1/#2/}
```

```
\foo[1]\quad\foo[2]\quad\foo[3]\par
\foo(1)\quad\foo(2)\quad\foo(3)\par
```

```
\tolerant\def\foo[#1]#;#={/#1/#2/}
```

```
\foo[1]\quad\foo[2]\quad\foo[3]\par
\foo{1}\quad\foo{2}\quad\foo{3}\par
```

```
\tolerant\def\foo[#1]#;#2{/#1/#2/}
```

```
\foo[1]\quad\foo[2]\quad\foo[3]\par
\foo{1}\quad\foo{2}\quad\foo{3}\par
```

```
\tolerant\def\foo[#1]#;( #2)#;#={/#1/#2/#3/}
```

```
\foo[1]\quad\foo[2]\quad\foo[3]\par
\foo(1)\quad\foo(2)\quad\foo(3)\par
\foo{1}\quad\foo{2}\quad\foo{3}\par
```

```
/1// /2// /3//
//1/ //2/ //3/
/1// /2// /3//
//1/ //2/ //3/
/1// /2// /3//
//1/ //2/ //3/
```

```

/1/// /2/// /3///
//1// //2// //3//
///1/ ///2/ ///3/

```

I have to admit that I don't really need it but it made some macros that I was redefining behave better, so there is some self-interest here. Anyway, I considered some other features, like picking up a detokenized argument but I don't expect that to be of much use. In the meantime we ran out of reasonable characters, but some day #? and #! might show up, or maybe I find a use for #< and #>. A summary of all this is given here:

---

+	keep the braces
-	discard and don't count the argument
/	remove leading and trailing spaces and pars
=	braces are mandate
_	braces are mandate and kept
^	keep leading spaces

---

1-9	an argument
0	discard but count the argument

---

*	ignore spaces
:	pick up scanning here
;	quit scanning

---

.	ignore pars and spaces
,	push back space when quit

---

The last two have not been discussed and were added later. The period directive gobbles space and par tokens and discards them in the process. The comma directive is like \* but it pushes back a space when the matching quits.

```
\tolerant\def\foo[#1]#;(#2){/#1/#2/}
```

```
\foo[1]\quad\foo[2]\quad\foo[3]\par
\foo(1)\quad\foo(2)\quad\foo(3)\par
```

```
\tolerant\def\foo[#1]#;#={/#1/#2/}
```

```
\foo[1]\quad\foo[2]\quad\foo[3]\par
\foo{1}\quad\foo{2}\quad\foo{3}\par
```

```
\tolerant\def\foo[#1]#;#2{/#1/#2/}
```

```
\foo[1]\quad\foo[2]\quad\foo[3]\par
\foo{1}\quad\foo{2}\quad\foo{3}\par
```

```
\tolerant\def\foo[#1]#;(#2)#;#={/#1/#2/#3/}
```

```
\foo[1]\quad\foo[2]\quad\foo[3]\par
```

```
\foo(1)\quad\foo(2)\quad\foo(3)\par
```

```
\foo{1}\quad\foo{2}\quad\foo{3}\par
```

```
/1// /2// /3//
```

```
//1/ //2/ //3/
```

```
/1// /2// /3//
```

```
//1/ //2/ //3/
```

```
/1// /2// /3//
```

```
//1/ //2/ //3/
```

```
/1/// /2/// /3///
```

```
//1// //2// //3//
```

```
///1/ ///2/ ///3/
```

Gobbling spaces versus pushing back is an interface design decision because it has to do with consistency.

## 6.3 Runaway arguments

There is a particular troublesome case left: a runaway argument. The solution is not pretty but it's the only way: we need to tell the parser that it can quit.

```
\tolerant\def\foo[#1=#2]%
  {\ifarguments 0\or 1\or 2\or 3\or 4\fi:\vl\type{#1}\vl\type{#2}\vl}
```

The outcome demonstrates that one still has to do some additional checking for sane results and there are alternative way to (ab)use this mechanism. It all boils down to a clever combination of delimiters and `\ignorearguments`.

```
\dontleavehmode \foo[a=1]
```

```
\dontleavehmode \foo[b=]
```

```
\dontleavehmode \foo[=]
```

```
\dontleavehmode \foo[x]\ignorearguments
```

All calls are accepted:

```
2:a|1|
```

```
2:b|
```

```
2:|
```

```
1:x|]
```

Just in case you wonder about performance: don't expect miracles here. On the one hand there is some extra overhead in the engine (when defining macros as well as when collecting arguments during a macro call) and maybe using these new features can sort of compensate that. As mentioned: the gain is mostly in cleaner macro code and less clutter in tracing. And I just want the ConT<sub>E</sub>Xt code to look nice: that way users can look in the source to see what happens and not drown in all these show-off tricks, special characters like underscores, at signs, question marks and exclamation marks.

For the record: I normally run tests to see if there are performance side effects and as long as processing the test suite that has thousands of files of all kind doesn't take more time it's okay. Actually, there is a little gain in ConT<sub>E</sub>Xt but that is to be expected, but I bet users won't notice it, because it's easily offset by some inefficient styling. Of course another gain of loosing some indirectness is that error messages point to the macro that the user called for and not to some follow up.

## 6.4 Introspection

A macro has a meaning. You can serialize that meaning as follows:

```
\tolerant\protected\def\foo#1[#2]*[#3]%
  {(1=#1) (2=#3) (3=#3)}
```

```
\meaning\foo
```

The meaning of \foo comes out as:

```
tolerant protected macro:#1[#2]*[#3]->(1=#1) (2=#3) (3=#3)
```

When you load the module system-tokens you can also say:

```
\luatokenable\foo
```

This produces a table of tokens specifications:

```
tolerant protected macro:#1[#2]*[#3]->(1=#1) (2=#3) (3=#3)
```

---

### **tolerant protected control sequence: foo**

---

597708	19	49	match		argument 1
597669	12	91	other char	[ U+0005B	
596436	19	50	match		argument 2
98043	12	93	other char	] U+0005D	
597827	19	42	match		argument *

593954	12	91	other char	[	U+0005B	
597519	19	51	match			argument 3
598166	12	93	other char	]	U+0005D	
594961	20	0	end match			
<hr/>						
594212	12	40	other char	(	U+00028	
592739	12	49	other char	1	U+00031	
595682	12	61	other char	=	U+0003D	
596674	21	1	parameter reference			
595528	12	41	other char	)	U+00029	
588801	10	32	spacer			
595228	12	40	other char	(	U+00028	
594408	12	50	other char	2	U+00032	
594619	12	61	other char	=	U+0003D	
597899	21	3	parameter reference			
597743	12	41	other char	)	U+00029	
598007	10	32	spacer			
594470	12	40	other char	(	U+00028	
596027	12	51	other char	3	U+00033	
594826	12	61	other char	=	U+0003D	
595904	21	3	parameter reference			
596661	12	41	other char	)	U+00029	

A token list is a linked list of tokens. The magic numbers in the first column are the token memory pointers. and because macros (and token lists) get recycled at some point the available tokens get scattered, which is reflected in the order of these numbers. Normally macros defined in the macro package are more sequential because they stay around from the start. The second and third row show the so called command code and the specifier. The command code groups primitives in categories, the specifier is an indicator of what specific action will follow, a register number a reference, etc. Users don't need to know these details. This macro is a special version of the online variant:

```
\showluatokens\foo
```

That one is always available and shows a similar list on the console. Again, users normally don't want to know such details.

## 6.5 nesting

You can nest macros, as in:

```
\def\foo#1#2{\def\oof##1{<#1>##1<#2>}}
```

At first sight the duplication of # looks strange but this is what happens. When T<sub>E</sub>X scans the definition of \foo it sees two arguments. Their specification ends up in the preamble that defines the matching. When the body is scanned, the #1 and #2 are turned into a parameter reference. In order to make nested macros with arguments possible a # followed by another # becomes just one #. Keep in mind that the definition of \oof is delayed till the macro \foo gets expanded. That definition is just stored and the only thing that get's replaced are the two references to a macro parameter

---

**control sequence: foo**

---

597848	19	49	match		argument 1
597779	19	50	match		argument 2
596648	20	0	end match		
597838	128	1	def		def
481042	146	0	tolerant call		oof
596535	6	35	parameter		
596429	12	49	other char	1	U+00031
594225	1	123	left brace		
595288	12	60	other char	<	U+0003C
597711	21	1	parameter reference		
594179	12	62	other char	>	U+0003E
597870	6	35	parameter		
597846	12	49	other char	1	U+00031
598046	12	60	other char	<	U+0003C
598005	21	2	parameter reference		
596674	12	62	other char	>	U+0003E
597839	2	125	right brace		

---

Now, when we look at these details, it might become clear why for instance we have ‘variable’ names like #4 and not #whatever (with or without hash). Macros are essentially token lists and token lists can be seen as a sequence of numbers. This is not that different from other programming environments. When you run into buzzwords like ‘bytecode’ and ‘virtual machines’ there is actually nothing special about it: some high level programming (using whatever concept, and in the case of T<sub>E</sub>X it's macros) eventually ends up as a sequence of instructions, say bytecodes. Then you need some machinery to run over that and act upon those numbers. It's something you arrive at naturally when you play with interpreting languages.<sup>7</sup>

---

<sup>7</sup> I actually did when I wrote an interpreter for some computer assisted learning system, think of a kind of interpreted Pascal, but later realized that it was a a bytecode plus virtual machine thing. I'd just applied what I learned when playing with eight bit processors that took bytes, and interpreted opcodes and such.



So, internally a #4 is just one token, a operator-operand combination where the operator is “grab a parameter” and the operand tells “where to store” it. Using names is of course an option but then one has to do more parsing and turn the name into a number<sup>8</sup>, add additional checking in the macro body, figure out some way to retain the name for the purpose of reporting (which then uses more token memory or strings). It is simply not worth the trouble, let alone the fact that we loose performance, and when T<sub>E</sub>X showed up those things really mattered.

It is also important to realize that a # becomes either a preamble token (grab an argument) or a reference token (inject the passed tokens into a new input level). Therefore the duplication of hash tokens ## that you see in macro nested bodies also makes sense: it makes it possible for the parser to distinguish between levels. Take:

```
\def\foo#1{\def\oof##1{#1##1#1}}
```

Of course one can think of this:

```
\def\foo#fence{\def\oof#text{#fence#text#fence}}
```

But such names really have to be unique then! Actually ConT<sub>E</sub>Xt does have an input method that supports such names, but discussing it here is a bit out of scope. Now, imagine that in the above case we use this:

```
\def\foo[#1][#2]{\def\oof##1{#1##1#2}}
```

If you're a bit familiar with the fact that T<sub>E</sub>X has a model of category codes you can imagine that a predictable “hash followed by a number” is way more robust than enforcing the user to ensure that catcodes of ‘names’ are in the right category (read: is a bracket part of the name or not). So, say that we go completely arbitrary names, we then suddenly needs some escaping, like:

```
\def\foo[#{left}][#{right}]{\def\oof#{text}#{left}#{text}#{right}}
```

And, if you ever looked into macro packages, you will notice that they differ in the way they assign category codes. Asking users to take that into account when defining macros makes not that much sense.

So, before one complains about T<sub>E</sub>X being obscure (the hash thing), think twice. Your demand for simplicity for your coding demand will make coding more cumbersome for

---

There's nothing spectacular about all this and I only realized decades later that the buzzwords describes old natural concepts.

<sup>8</sup> This is kind of what MetaPost does with parameters to macros. The side effect is that in reporting you get text0, expr2 and such reported which doesn't make things more clear.

the complex cases that macro packages have to deal with. It's comparable using  $\TeX$  for input or using (say) mark down. For simple documents the later is fine, but when things become complex, you end up with similar complexity (or even worse because you lost the enforced detailed structure). So, just accept the unavoidable: any language has its peculiar properties (and for sure I do know why I dislike some languages for it). The  $\TeX$  system is not the only one where dollars, percent signs, ampersands and hashes have special meaning.

## 6.6 Prefixes

Traditional  $\TeX$  has three prefixes that can be used with macros: `\global`, `\outer` and `\long`. The last two are no-op's in  $\text{LuaMeta}\TeX$  and if you want to know what they do (did) you can look it up in the  $\TeX$ book. The  $\varepsilon$ - $\TeX$  extension gave us `\protected`.

In  $\text{LuaMeta}\TeX$  we have `\global`, `\protected`, `\tolerant` and overload related prefixes like `\frozen`. A protected macro is one that doesn't expand in an expandable context, so for instance inside an `\edef`. You can force expansion by using the `\expand` primitive in front which is also something  $\text{LuaMeta}\TeX$ .

Frozen macros cannot be redefined without some effort. This feature can to some extent be used to prevent a user from overloading, but it also makes it harder for the macro package itself to redefine on the fly. You can remove the lock with `\unletfrozen` and add a lock with `\letfrozen` so in the end users still have all the freedoms that  $\TeX$  normally provides.

```

                \def\foo{foo} 1: \meaning\foo
        \frozen\def\foo{foo} 2: \meaning\foo
    \unletfrozen  \foo      3: \meaning\foo
\protected\frozen\def\foo{foo} 4: \meaning\foo
    \unletfrozen  \foo      5: \meaning\foo

```

```

1: macro:foo
2: macro:foo
3: macro:foo
4: protected macro:foo
5: protected macro:foo

```

This actually only works when you have set `\overloadmode` to a value that permits redefining a frozen macro, so for the purpose of this example we set it to zero.

A `\tolerant` macro is one that will quit scanning arguments when a delimiter cannot be matched. We saw examples of that in a previous section.

These prefixes can be chained (in arbitrary order):

```
\frozen\tolerant\protected\global\def\foo[#1]*[#2]{...}
```

There is actually an additional prefix, `\immediate` but that one is there as signal for a macro that is defined in and handled by Lua. This prefix can then perform the same function as the one in traditional  $\TeX$ , where it is used for backend related tasks like `\write`.

Now, the question is of course, to what extent will Con $\TeX$ t use these new features. One important argument in favor of using `\tolerant` is that it gives (hopefully) better error messages. It also needs less code due to lack of indirectness. Using `\frozen` adds some safeguards although in some places where Con $\TeX$ t itself overloads commands, we need to defrost. Adapting the code is a tedious process and it can introduce errors due to mistypings, although these can easily be fixed. So, it will be used but it will take a while to adapt the code base.

One problem with frozen macros is that they don't play nice with for instance `\futurelet`. Also, there are places in Con $\TeX$ t where we actually do redefine some core macro that we also want to protect from redefinition by a user. One can of course `\unletfrozen` such a command first but as a bonus we have a prefix `\overloaded` that can be used as prefix. So, one can easily redefine a frozen macro but it takes a little effort. After all, this feature is mainly meant to protect a user for side effects of definitions, and not as final blocker.<sup>9</sup>

A frozen macro can still be overloaded, so what if we want to prevent that? For this we have the `\permanent` prefix. Internally we also create primitives but we don't have a prefix for that. But we do have one for a very special case which we demonstrate with an example:

```
\def\F00 % trickery needed to pick up an optional argument
  {\noalign{\vskip10pt}}

\noaligned\protected\tolerant\def\00F[#1]%
  {\noalign{\vskip\iftok{#1}\emptytoks10pt\else#1\fi}}

\starttabulate[|l|l|]
  \NC test \NC test \NC \NR
  \NC test \NC test \NC \NR
```

<sup>9</sup> As usual adding features like this takes some experimenting and we're now at the third variant of the implementation, so we're getting there. The fact that we can apply such features in large macro package like Con $\TeX$ t helps figuring out the needs and best approaches.

```

\F00
\NC test \NC test \NC \NR
\00F[30pt]
\NC test \NC test \NC \NR
\00F
\NC test \NC test \NC \NR
\stoptabulate

```

When  $\text{T}_{\text{E}}\text{X}$  scans input (from a file or token list) and starts an alignment, it will pick up rows. When a row is finished it will look ahead for a `\noalign` and it expands the next token. However, when that token is protected, the scanner will not see a `\noalign` in that macro so it will likely start complaining when that next macro does get expanded and produces a `\noalign` when a cell is built. The `\noaligned` prefix flags a macro as being one that will do some `\noalign` as part of its expansion. This trick permits clean macros that pick up arguments. Of course it can be done with traditional means but this whole exercise is about making the code look nice.

The table comes out as:

```

test test
test test

test test

test test

test test

```

One can check the flags with `\ifflags` which takes a control sequence and a number, where valid numbers are:

1 frozen	2 permanent	4 immutable	8 primitive
16 mutable	32 noaligned	64 instance	

The level of checking is controlled with the `\overloadmode` but I'm still not sure about how many levels we need there. A zero value disables checking, the values 1 and 3 give warnings and the values 2 and 4 trigger an error.

## 6.7 Arguments

The number of arguments that a macro takes is traditionally limited to nine (or ten if one takes the trailing `#` into account). That this is enough for most cases is demonstrated

by the fact that ConTeXt has only a handful of macros that use #9. The reason for this limitation is in part a side effect of the way the macro preamble and arguments are parsed. However, because in LuaMetaTeX we use a different implementation, it was not that hard to permit a few more arguments, which is why we support upto 15 arguments, as in:

```
\def\foo#1#2#3#4#5#6#7#8#9#A#B#C#D#E#F{...}
```

We can support the whole alphabet without much trouble but somehow sticking to the hexadecimal numbers makes sense. It is unlikely that the core of ConTeXt will use this option but sometimes at the user level it can be handy. The penalty in terms of performance can be neglected.

```
\tolerant\def\foo#=#=#=#=#=#=#=#=#=#=#=#=#=#=#=#=#=#=%
  {(#1)(#2)(#3)(#4)(#5)(#6)(#7)(#8)(#9)(#A)(#B)(#C)(#D)(#E)(#F)}

\foo{1}{2}
```

In the previous example we have 15 optional arguments where braces are mandate (otherwise we the scanner happily scoops up what follows which for sure gives some error).

## 6.8 Constants

The LuaMetaTeX engine has lots of efficiency tricks in the macro parsing and expansion code that makes it not only fast but also let is use less memory. However, every time that the body of a macro is to be injected the expansion machinery kicks in. This often means that a copy is made (pushed in the input and used afterwards). There are however cases where the body is just a list of character tokens (with category letter or other) and no expansion run over the list is needed.

It is tempting to introduce a string data type that just stores strings and although that might happen at some point it has the disadvantage that one need to tokenize that string in order to be able to use it, which then defeats the gain. An alternative has been found in constant macros, that is: a macro without parameters and a body that is considered to be expanded and never freed by redefinition. There are two variants:

```
\cdef      \foo        {whatever}
\cdefcsname foo\endcsname{whatever}
```

These are actually just equivalents to

```
\edef      \foo        {whatever}
```

```
\edefcsname foo\endcsname{whatever}
```

just to make sure that the body gets expanded at definition time but they are also marked as being constant which in some cases might give some gain, for instance when used in csname construction. The gain is less than one expects although there are a few cases in ConT<sub>E</sub>Xt where extreme usage of parameters benefits from it. Users are unlikely to use these two primitives.

Another example of a constant usage is this:

```
\lettonothing\foo
```

which gives `\foo` an empty body. That one is used in the core, if only because it gives a bit smaller code. Performance is not that different from

```
\let\foo\empty
```

but it saves one token (8 bytes) when used in a macro. The assignment itself is not that different because `\foo` is made an alias to `\empty` which in turn only needs incrementing a reference counter.

## 6.9 Passing parameters

When you define a macro, the `#1` and more parameters are embedded as a reference to a parameter that is passed. When we have four parameters, the parameter stack has four entries and when an entry is eventually accessed a new input level is pushed and tokens are fetched from that list. This has some side effects when we check a parameter. This can happen multiple times, depending on how often we access a parameter. Take the following:

```
\def\oof#1{#1}
```

```
\tolerant\def\foo[#1]#*[#2]%
  {1:\ifparameter#1\or Y\else N\fi\quad
  2:\ifparameter#2\or Y\else N\fi\quad
  \oof{3:\ifparameter #1\or Y\else N\fi\quad
    4:\ifparameter #2\or Y\else N\fi\quad}%
  \par}
```

```
\foo \foo[] \foo[][] \foo[A] \foo[A][B]
```

This gives:

```

1:N 2:N 3:N 4:N
1:N 2:N 3:N 4:N
1:N 2:N 3:N 4:N
1:Y 2:N 3:Y 4:N
1:Y 2:Y 3:Y 4:Y

```

as you probably expect. However the first two checks are different from the embedded checks because they can check against the parameter reference. When we expand `\oof` its argument gets passed to the macro as a list and when the scanner collects the next token it will then push the parameter content on the input stack. So, then, instead of a reference we get the referenced parameter list. Internally that means that in 3 and 4 we check for a token and not for the length of the list (as in case 1 & 2). This means that

```

\iftok{#1}\emptytoks Y\else N\fi
\ifparameter#1\or Y\else N\fi

```

are different. In the first case we have a proper token list and nested conditionals in that list are okay. In the second case we just look ahead to see if there is an `\or`, `\else` or other condition related command and if so we decide that there is no parameter. So, if `\ifparameter` is a suitable check for empty depends on the need for expansion.

When you define macros that themselves call macros that should operate on the arguments of its parent you can easily pass these:

```

\def\foo#1#2%
  {\oof{#1}{#2}{P}%
  \oof{#1}{#2}{Q}%
  \oof{#1}{#2}{R}}

```

```

\def\oof#1#2#3%
  {[#1][#1]%
  #3%
  [#2][#2]}

```

Here the nested call to `\oof` involved three passed parameters. You can avoid that as follows:

```

\def\foo#1#2%
  {\def\MyIndexOne{#1}%
  \def\MyIndexTwo{#2}%
  \oof{P}\oof{Q}\oof{R}}

```

```
\def\oof#1%
  {(\MyIndexOne)(\MyIndexOne)%
   #1%
   (\MyIndexTwo)(\MyIndexTwo)}
```

You can also do this:

```
\def\foo#1#2%
  {\def\oof##1%
   {/#1/#2/%
   ##1%
   /#1//#2/}%
   \oof{P}\oof{Q}\oof{R}}
```

These parameters indicated by # in the macro body are in fact references. When we call for instance `\foo{1}{2}` the two parameters get pushed on a parameter stack and the embodied references point to these stack entries. By the time that body gets expanded  $\TeX$  bumps the input level and pushes the parameter list onto the input stack. It then continues expansion. The parameter is not copied, because it can't be changed anyway. The only penalty in terms of performance and memory usage is the pushing and popping of the input. So how does that work out for these three cases?

When in the first case the `\oof{#1}{#2}{P}` is seen,  $\TeX$  starts expanding the `\oof` macro. That one expects three arguments. The #1 reference is seen and in this case a copy of that parameter is passed. The same is true for the other two. Then, inside `\oof` expansion happens on the parameters on the stack and no copies have to be made there.

The second case defines two macros so again two copies are made that make the bodies of these macros. This comes at the cost of some runtime and memory. However, this time with `\oof{P}` only one argument gets passed and instead expansion of the macros happen in there.

Normally macro arguments are not that large but there can be situations where we really want to avoid useless copying. This not only saves memory but also can give a bit better performance. In the examples above the second variant is some 10% faster than the first one. We can gain another 10% with the following trick:

```
\def\foo#1#2%
  {\parameterdef\MyIndexOne\plusone % 1
   \parameterdef\MyIndexTwo\plustwo % 2
   \oof{P}\oof{Q}\oof{R}\norelax}
```



```
\def\oof#1%
  {<\MyIndexOne><\MyIndexOne>%
  #1%
  <\MyIndexTwo><\MyIndexTwo>}
```

Here we define an explicit parameter reference that we access later on. There is the overhead of a definition but it can be neglected. We use that reference (abstraction) in `\oof`. Actually you can use that reference in any call down the chain.

When applied to `\foo{1}{2}` the four variants above give us:

```
[1][1]P[2][2][1][1]Q[2][2][1][1]R[2][2]
(1)(1)P(2)(2)(1)(1)Q(2)(2)(1)(1)R(2)(2)
/1/2/P/1//2//1/2/Q/1//2//1/2/R/1//2/
<1><1>P<2><2><1><1>Q<2><2><1><1>R<2><2>
```

Before we had `parameterdef` we had this:

```
\def\foo#1#2%
  {\integerdef\MyIndexOne\parameterindex\plusone % 1
  \integerdef\MyIndexTwo\parameterindex\plustwo % 2
  \oof{P}\oof{Q}\oof{R}\norelax}
```

```
\def\oof#1%
  {<\expandparameter\MyIndexOne><\expandparameter\MyIndexOne>%
  #1%
  <\expandparameter\MyIndexTwo><\expandparameter\MyIndexTwo>}
```

It involves more tokens, is a bit less abstract, but as it is a cheap extension we kept it. It actually demonstrates that one can access parameters in the stack by index, but it one then needs to keep track of where access takes place. In principle one can debug the call chain this way.

To come back to performance and memory usage, when the arguments become larger the fourth variant with the `\parameterdef` quickly gains over the others. But it only shows in exceptional usage. This mechanism is more about abstraction: it permits us to efficiently turn arguments into local variables without the overhead involved in creating macros. You can test if a parameter is set

```
\tolerant\protected\def\MyMacro[#1]#:#2%
  {\parameterdef\MyArgumentOne\plusone
  \parameterdef\MyArgumentTwo\plustwo
  \ifparameter\MyArgumentOne\or
```

```

(\MyArgumentOne)
\fi
/\MyArgumentTwo/}

```

```

\MyMacro[one]{two}
\MyMacro{two}

```

Indeed we get:

```
(one) /two/ /two/
```

Of course `\ifparameter#1\or...` is more efficient but once you use named parameters like this it's probably not something you're worry too much about,

## 6.10 Nesting

We also have a few preamble features that relate to nesting. Although we can do without (as shown for years in LMTX) they do have some benefits. They are discussed as group here and because they are only useful for low level programming we stick to simple examples. The `#L` and `#R` use the following token as delimiters. Here we use `[` and `]` but they can be a `\cs` as well. Nested delimiters are handled well.

The `#S` grabs the argument till the next final square bracket `]` but in the process will grab nested with it sees a `[`. The `#P` does the same for parentheses and `#X` for angle brackets. In the next examples the `#*` just gobbles optional spaces but we've seen that one already.

The `#G` argument just registers the next token as delimiter but it will grab multiple of them. The `#M` gobbles more: in addition to the delimiter spaces are gobbled.

```

\tolerant\def\fooA      [#1]{(#1)}
\tolerant\def\fooB      [#L[#R]#1]{(#1)}
\tolerant\def\fooC      #S#1{(#1)}
\tolerant\def\fooE      #S#1,{(#1)}
\tolerant\def\fooF      #S#1#*#S#2{(#1/#2)}
\tolerant\def\fooG      [#1]#S[#2]#*#S[#3]{(#1/#2/#3)}
\tolerant\def\fooH      [#1][#S#2]#*[#S#3]{(#1/#2/#3)}
\tolerant\def\fooI      #1=#2#G,{(#1=#2)}
\tolerant\def\fooJ      #1=#2#M,{(#1=#2)}

\fooA[x]                (x)                (x)
\fooB[x]                (x)                (x)

```

<code>\fooC[1[2]3[4]5]</code>	<code>([1[2]3[4]5])</code>	<code>(1[2]3[4]5)</code>
<code>\fooE X[, ]X,</code>	<code>(X[, ]X)</code>	<code>(X[, ]X)</code>
<code>\fooF[A] [B]</code>	<code>([A]/[B])</code>	<code>(A/B)</code>
<code>\fooF[] []</code>	<code>([]/[])</code>	<code>(/)</code>
<code>\fooG[a][b][c]</code>	<code>(a/b/c)</code>	<code>(a/b/c)</code>
<code>\fooG[a][b]</code>	<code>(a/b/)</code>	<code>(a/b/)</code>
<code>\fooG[a]</code>	<code>(a//)</code>	<code>(a//)</code>
<code>\fooG[a][x[x]x][c]</code>	<code>(a/x[x]x/c)</code>	<code>(a/x[x]x/c)</code>
<code>\fooH[a][x[x]x][c]</code>	<code>(a/x[x]x/c)</code>	<code>(a/x[x]x/c)</code>
<code>\fooI X=X, , ,</code>	<code>(X=X)</code>	<code>(X=X)</code>
<code>\fooJ X=X, , ,</code>	<code>(X=X)</code>	<code>(X=X)</code>

These features make it possible to support nested setups more efficiently and also makes it possible to accept values that contain balanced brackets in setup commands without additional overhead. Although it has never been an issue to let users specify:

```
\defineoverlay[whatever][{some \command[withparameters] here}]
```

```
\setupfoo[before={\blank[big]}]
```

it might be less confusing to permit:

```
\defineoverlay[whatever][some \command[withparameters] here]
```

```
\setupfoo[before=\blank[big]]
```

as well, if only because occasionally users get hit by this.

## 6.11 Duplicate hashes

In  $\TeX$  every character has a so called category code. Most characters are classified as ‘letter’ (they make up words) or as ‘other’. In Unicode we distinguish symbols, punctuation, and more, but in  $\TeX$  these are all of category ‘other’. In math however we can classify them differently but in this perspective we ignore that. The backslash has category ‘escape’ and it starts a control sequence. The curly braces are (internally) of category ‘left brace’ and ‘right brace’ aka ‘begin group’ and ‘end group’ but, no matter what they are called, they begin and end something: a group, argument, token list, box, etc. Any character can have those categories. Although it would look strange to a  $\TeX$  user, this can be made valid:

```
!protected !gdef !weird¶1
B
```

```

    something: ¶1
E
!weird BhereE

```

In such a setup spaces can be of category ‘invisible’. The paragraph symbol takes the place of the hash as parameter identifier. The next code shows how this is done. Here we wrap all in a macro so that we don't get catcode interference in the document source.

```

\def\NotSoTeX
  {\begingroup
   \catcode `B \begingroupcatcode
   \catcode `E \endgroupcatcode
   \catcode `¶ \parametercatcode
   \catcode `! \escapecatcode
   \catcode 32 \ignorecatcode
   \catcode 13 \ignorecatcode
   % this buffer has a definition:
   \getbuffer
   % which is now known globally
   \endgroup}
\NotSoTeX
\weird{there}

```

This results in:

```

something:here
something:there

```

In the first line the !, B and E are used as escape and argument delimiters, in the second one we use the normal characters. When we show the `\meaningasis` we get:

```
\global \protected \def \weird #1{something:#1}
```

or in more detail:

---

**protected control sequence: weird**

---

594662	19	49	match		argument 1
588777	20	0	end match		
596531	11	115	letter	s	U+00073
597729	11	111	letter	o	U+0006F
598566	11	109	letter	m	U+0006D
596568	11	101	letter	e	U+00065

---

597567	11	116	letter	t	U+00074
598001	11	104	letter	h	U+00068
596431	11	105	letter	i	U+00069
595482	11	110	letter	n	U+0006E
596708	11	103	letter	g	U+00067
597830	12	58	other char	:	U+0003A
597708	21	1	parameter reference		

---

So, no matter how we set up the system, in the end we get some generic representation. When we see #1 in 'print' it can be either two tokens, # (catcode parameter) followed by 1 with catcode other, or one token referring to parameter 1 where the character 1 is the opcode of an internal 'reference command'. In order to distinguish a reference from the two token case, parameter hash tokens get shown as doubles.

```
\def\test #1{x#1x##1x####1x}
\def\tset ¶1{x¶1x¶¶1x¶¶¶¶1x}
```

And with \meaning we get, consistent with the input:

```
macro:#1->x#1x##1x####1x
macro:#1->x#1x¶¶1x¶¶¶¶1x
```

These are equivalent, apart from the parameter character in the body of the definition:

---

**control sequence: test**

---

593944	19	49	match		argument 1
598652	20	0	end match		

---

598574	11	120	letter	x	U+00078
597492	21	1	parameter reference		
595827	11	120	letter	x	U+00078
597247	6	35	parameter		
599026	12	49	other char	1	U+00031
597829	11	120	letter	x	U+00078
594992	6	35	parameter		
594242	6	35	parameter		
598647	12	49	other char	1	U+00031
597031	11	120	letter	x	U+00078

---

**control sequence: tset**

---

598777	19	49	match		argument 1
--------	----	----	-------	--	------------

597858	20	0	end match		
<hr/>					
598891	11	120	letter	x	U+00078
594766	21	1	parameter reference		
596993	11	120	letter	x	U+00078
594676	6	182	parameter		
598778	12	49	other char	1	U+00031
598047	11	120	letter	x	U+00078
598612	6	182	parameter		
597826	6	182	parameter		
593503	12	49	other char	1	U+00031
598908	11	120	letter	x	U+00078
<hr/>					

Watch how every ‘parameter’ is just a character with the Unicode index of the used input character as property. Let us summarize the process. When a single parameter character is seen in the input, the next character determines how it will be interpreted. If there is a digit then it becomes a reference to a parameter in the preamble, and when followed by another parameter character it will be appended to the body of the macro and that second one is dropped. So, two parameter characters become one, and four become two. One parameter character becomes a reference and from that you can guess what three in a row become. However, when  $\text{T}_{\text{E}}\text{X}$  is showing the macro definition (using meaning) the hashes get duplicated in order to distinguish parameter references from parameter characters that were kept (e.g. for nested definitions). One can make an argument for  $\backslash\text{parameterchar}$  as we also have  $\backslash\text{escapechar}$  but by now this convention is settled and it doesn't look that bad anyway.

We now come to the more tricky part with respect to the doubling of hashes. When  $\text{T}_{\text{E}}\text{X}$  was written its application landscape looked a bit different. For instance, fonts were limited and therefore it was natural to access special characters by name. Using  $\backslash\#$  to get a hash in the text was not that problematic, if one needed that character at all. The same can be said for the braces, backslash and even the dollar (after all  $\text{T}_{\text{E}}\text{X}$  is free software).

But what if we have more visualization and/or serialization than meanings and tracing? When we opened up the internals in  $\text{LuaT}_{\text{E}}\text{X}$  and even more in  $\text{LuaMetaT}_{\text{E}}\text{X}$  the duplicating of hashes became a bit of a problem. There we don't need to distinguish between a parameter reference and a parameter character because by that time these references are resolved. All hashes that we encounter are just that: hashes. And this is why in  $\text{LuaMetaT}_{\text{E}}\text{X}$  we disable the duplication for those cases where it serves no purpose.

When the engine scans a macro definition it starts with picking up the name of the macro. Then it starts scanning the preamble up to the left brace. In the preamble of a

### Duplicate hashes

macro the scanner converts hashes followed by another token into single match token. Then when the macro body is scanned single hashes followed by a number become a reference, while double hashes become one hash and get interpreted at expansion time (possibly triggering an error when not followed by a valid specifier like a number). In traditional T<sub>E</sub>X we basically had this:

```
\def\test#1{#1}
\def\test#1{##}
\def\test#1{#X}
\def\test#1{##1}
```

There can be a trailing # in the preamble for special purposes but we forget about that now. The first definition is valid, the second definition is invalid when the macro is expanded and the third definition triggers an error at definition time. The last definition will again trigger an error at expansion time.

However, in LuaMetaT<sub>E</sub>X we have an extended preamble where the following preamble parameters are handled (some only in tolerant mode):

#n	parameter	index 1 upto E
#0	throw away parameter	increment index
#-	ignore parameter	keep index
#*	gobble white space	
#+	keep (honor) the braces	
#.	ignore pars and spaces	
#,	push back space when no match	
#/	remove leading and trailing spaces and pars	
#=	braces are mandate	
#^	keep leading spaces	
#_	braces are mandate and kept (obey)	
#@	par delimiter	only for internal usage
#:	pick up scanning here	
#;	quit scanning	
#L	left delimiter token	followed by token
#R	right delimiter token	followed by token
#G	gobble token	followed by token
#M	gobble token and spaces	followed by token

### Duplicate hashes

#S	nest square brackets	only inner pairs
#X	nest angle brackets	only inner pairs
#P	nest parentheses	only inner pairs

As mentioned these will become so called match tokens and only when we show the meaning the hash will show up again.

```
\def\test[#1]#*[*S#2]{.#1.#2.}
```

---

**control sequence: test**

---

594271	12	91	other char	[	U+0005B	
481151	19	49	match			argument 1
596027	12	93	other char	]	U+0005D	
592748	19	42	match			argument *
598347	12	91	other char	[	U+0005B	
597696	12	42	other char	*	U+0002A	
596988	11	83	letter	S	U+00053	
596640	19	50	match			argument 2
594390	12	93	other char	]	U+0005D	
598135	20	0	end match			
<hr/>						
597978	12	46	other char	.	U+0002E	
597023	21	1	parameter reference			
598660	12	46	other char	.	U+0002E	
597934	21	2	parameter reference			
599056	12	46	other char	.	U+0002E	

---

This means that in the body of a macro you will not see `#*` show up. It is just a directive that tells the macro parser that spaces are to be skipped. The `#S` directive makes the parser for the second parameter handle nested square bracket. The only hash that we can see end up in the body is the one that we entered as double hash (then turned single) followed by (in traditional terms) a number that when all gets parsed with then become a reference: the sequence `##1` internally is `#1` and becomes 'reference to parameter 1' assuming that we define a macro in that body. If no number is there, an error is issued. This opens up the possibility to add more variants because it will only break compatibility with respect to what is seen as error. As with the preamble extensions, old documents that have them would have crashed before they became available.

So, this means that in the body, and actually anywhere in the document apart from preambles, we now support the following general parameter specifiers. Keep in mind that they expand in an expansion context which can be tricky when they overlap with

### Duplicate hashes



preamble entries, like for instance `#R` in such an expansion. Future extensions can add more so *any* hashed shortcut is sensitive for that.

<code>#I</code>	current iterator	<code>\currentloopiterator</code>
<code>#P</code>	parent iterator	<code>\previousloopiterator 1</code>
<code>#G</code>	grandparent iterator	<code>\previousloopiterator 2</code>
<code>#H</code>	hash escape	<code>#</code>
<code>#S</code>	space escape	<code>␣</code>
<code>#T</code>	tab escape	<code>\t</code>
<code>#L</code>	newline escape	<code>\n</code>
<code>#R</code>	return escape	<code>\r</code>
<code>#X</code>	backslash escape	<code>\</code>
<code>#N</code>	nbsp	U+00A0 (under consideration)
<code>#Z</code>	zws	U+200B (under consideration)

Some will now argue that we already have `^^` escapes in  $\TeX$  and `^^^^` and `^^^^^^` in  $\text{Lua}\TeX$  and that is true. However, these can be disabled, and in  $\text{Con}\TeX$ t they are, where we instead enable the prescript, postscript, and index features in mathmode and there type `^` and `_` are used. Even more: in  $\text{Con}\TeX$ t we just let `^`, `_` and `&` be what they are. Occasionally I consider `$` to be just that but as I don't have dollars I will happily leave that for inline math. When users are not defining macros or are using the alternative definitions we can consider making the `#` a hash. An excellent discussion of how  $\TeX$  reads it's input and changes state accordingly can be found in Victor Eijkhout's " $\TeX$  By Topic", section 2.6: when `^^` is followed by a character with  $v < 128$  the interpreter will inject a character with code  $v - 64$ . When followed by two (!) lowercase hexadecimal characters, the corresponding character will be injected. Anyway, it not only looks kind of ugly, it also is somewhat weird because what follows is interpreted mixed way. The substitution happens early on (which is okay). But, how about the output? Traditional  $\TeX$  serializes special characters with a similar syntax but that has become optional when eight bit mode was added to the engines, it is configurable in  $\text{Lua}\TeX$  and has been dropped in  $\text{LuaMeta}\TeX$ : we operate in a utf universum.

## 6.11 Colofon

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$\text{LuaMeta}\TeX$	2.11.07   20250705
Support	<a href="http://www.pragma-ade.com">www.pragma-ade.com</a> <a href="http://contextgarden.net">contextgarden.net</a> <a href="mailto:ntg-context@ntg.nl">ntg-context@ntg.nl</a>

## 7 Grouping

low level

TEX

grouping

## Contents

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## 7.1 Introduction

This is a rather short explanation. I decided to write it after presenting the other topics at the 2019 ConT<sub>E</sub>Xt meeting where there was a question about grouping.

### 7.1.1 Pascal

In a language like Pascal, the language that T<sub>E</sub>X has been written in, or Modula, its successor, there is no concept of grouping like in T<sub>E</sub>X. But we can find keywords that suggests this:

```
for i := 1 to 10 do begin ... end
```

This language probably inspired some of the syntax of T<sub>E</sub>X and MetaPost. For instance an assignment in MetaPost uses := too. However, the begin and end don't really group but define a block of statements. You can have local variables in a procedure or function but the block is just a way to pack a sequence of statements.

### 7.1.2 T<sub>E</sub>X

In T<sub>E</sub>X macros (or source code) the following can occur:

```
\begingroup
...
\endgroup
```

as well as:

```
\bgroup
...
\egroup
```

Here we really group in the sense that assignments to variables inside a group are forgotten afterwards. All assignments are local to the group unless they are explicitly done global:

```
\scratchcounter=1
\def\foo{foo}
\begingroup
  \scratchcounter=2
  \global\globalscratchcounter=2
  \gdef\foo{F00}
\endgroup
```

Here `\scratchcounter` is still one after the group is left but its global counterpart is now two. The `\foo` macro is also changed globally.

Although you can use both sets of commands to group, you cannot mix them, so this will trigger an error:

```
\bgroup
\endgroup
```

The bottomline is: if you want a value to persist after the group, you need to explicitly change its value globally. This makes a lot of sense in the perspective of  $\text{T}_{\text{E}}\text{X}$ .

### 7.1.3 MetaPost

The MetaPost language also has a concept of grouping but in this case it's more like a programming language.

```
begingroup ;
  n := 123 ;
endgroup ;
```

In this case the value of `n` is 123 after the group is left, unless you do this (for numerics there is actually no need to declare them):

```
begingroup ;
  save n ; numeric n ; n := 123 ;
endgroup ;
```

Given the use of MetaPost (read: MetaFont) this makes a lot of sense: often you use macros to simplify code and you do want variables to change. Grouping in this language

serves other purposes, like hiding what is between these commands and let the last expression become the result. In a `vardef` grouping is implicit.

So, in MetaPost all assignments are global, unless a variable is explicitly saved inside a group.

### 7.1.4 Lua

In Lua all assignments are global unless a variable is defines local:

```
local x = 1
local y = 1
for i = 1, 10 do
    local x = 2
    y = 2
end
```

Here the value of `x` after the loop is still one but `y` is now two. As in Lua<sub>TEX</sub> we mix <sub>TEX</sub>, MetaPost and Lua you can mix up these concepts. Another mixup is using `:=`, `endfor`, `fi` in Lua after done some MetaPost coding or using `end` instead of `endfor` in MetaPost which can make the library wait for more without triggering an error. Proper syntax highlighting in an editor clearly helps.

### 7.1.5 C

The Lua language is a mix between Pascal (which is one reason why I like it) and C.

```
int x = 1 ;
int y = 1 ;
for (i=1; i<=10;i++) {
    int x = 2 ;
    y = 2 ;
}
```

The semicolon is also used in Pascal but there it is a separator and not a statement end, while in MetaPost it does end a statement (expression).

## 7.2 Kinds of grouping

Explicit grouping is accomplished by the two grouping primitives:

```
\begingroup
  \sl render slanted here
\endgroup
```

However, often you will find this being used:

```
{\sl render slanted here}
```

This is not only more compact but also avoids the `\endgroup` gobbling following spaces when used inline. The next code is equivalent but also suffers from the gobbling:

```
\bgroup
  \sl render slanted here
\egroup
```

The `\bgroup` and `\egroup` commands are not primitives but aliases (made by `\let`) to the left and right curly brace. These two characters have so called category codes that signal that they can be used for grouping. The *can be* here suggest that there are other purposes and indeed there are, for instance in:

```
\toks 0 = {abs}
\hbox {def}
```

In the case of a token list assignment the curly braces fence the assignment, so scanning stops when a matching right brace is found. The following are all valid:

```
\toks 0 = {a{b}s}
\toks 0 = \bgroup a{b}s}
\toks 0 = {a{\bgroup b}s}
\toks 0 = {a{\egroup b}s}
\toks 0 = \bgroup a{\bgroup b}s}
\toks 0 = \bgroup a{\egroup b}s}
```

They have in common that the final fence has to be a right brace. That the first one can be an alias is due to the fact that the scanner searches for a brace equivalent when it looks for the value. Because the equal is optional, there is some lookahead involved which involves expansion and possibly push back while once scanning for the content starts just tokens are collected, with a fast check for nested and final braces.

In the case of the box, all these specifications are valid:

```
\hbox {def}
\hbox \bgroup def\egroup
```

## Kinds of grouping

```

\hbox \bgroup def}
\hbox \bgroup d{e\egroup f}
\hbox {def\egroup

```

This is because now the braces and equivalent act as grouping symbols so as long as they match we're fine. There is a pitfall here: you cannot mix and match different grouping, so the next issues an error:

```

\bgroup xxx\endgroup % error
\begingroup xxx\egroup % error

```

This can make it somewhat hard to write generic grouping macros without trickery that is not always obvious to the user. Fortunately it can be hidden in macros like the helper `\groupedcommand`. In LuaMetaTeX we have a clean way out of this dilemma:

```

\beginsimplegroup xxx\endsimplegroup
\beginsimplegroup xxx\endgroup
\beginsimplegroup xxx\egroup

```

When you start a group with `\beginsimplegroup` you can end it in the three ways shows above. This means that the user (or calling macro) doesn't take into account what kind of grouping was used to start with.

When we are in math mode things are different. First of all, grouping with `\begin-group` and `\endgroup` in some cases works as expected, but because the math input is converted in a list that gets processed later some settings can become persistent, like changes in style or family. You can bet better use `\beginmathgroup` and `\endmathgroup` as they restore some properties. We also just mention the `\frozen` prefix that can be used to freeze assignments to some math specific parameters inside a group.

## 7.3 Hooks

In addition to the original `\aftergroup` primitive we have some more hooks. They can best be demonstrated with an example:

```

\begingroup \bf
%
\aftergroup A \aftergroup 1
\atendofgroup B \atendofgroup 1
%
\aftergrouped {A2}
\atendofgrouped {B2}

```



```

%
test
\endgroup

```

These collectors are accumulative. Watch how the bold is applied to what we inject before the group ends.

**test B1B2A1A2**

## 7.4 Local versus global

When  $\text{T}_{\text{E}}\text{X}$  enters a group and an assignment is made the current value is stored on the save stack, and at the end of the group the original value is restored. In  $\text{LuaMetaT}_{\text{E}}\text{X}$  this mechanism is made a bit more efficient by avoiding redundant stack entries. This is also why the next feature can give unexpected results when not used wisely.

Now consider the following example:

```

\newdimension\MyDimension

\starttabulate[||||]
  \NC      \MyDimension10pt \the\MyDimension
  \NC \advance\MyDimension10pt \the\MyDimension
  \NC \advance\MyDimension10pt \the\MyDimension \NC \NR
  \NC      \MyDimension10pt \the\MyDimension
  \NC \advance\MyDimension10pt \the\MyDimension
  \NC \advance\MyDimension10pt \the\MyDimension \NC \NR
\stoptabulate

10.0pt 10.0pt 10.0pt
10.0pt 10.0pt 10.0pt

```

The reason why we get the same values is that cells are a group and therefore the value gets restored as we move on. We can use the `\global` prefix to get around this

```

\starttabulate[||||]
  \NC \global      \MyDimension10pt \the\MyDimension
  \NC \global\advance\MyDimension10pt \the\MyDimension
  \NC \global\advance\MyDimension10pt \the\MyDimension \NC \NR
  \NC \global      \MyDimension10pt \the\MyDimension
  \NC \global\advance\MyDimension10pt \the\MyDimension
  \NC \global\advance\MyDimension10pt \the\MyDimension \NC \NR

```

**\stoptabulate**

```
10.0pt 20.0pt 30.0pt
10.0pt 20.0pt 30.0pt
```

Instead of using a global assignment and increment we can also use the following

**\constrained\MyDimension\zeropoint****\starttabulate[||||]**

```
\NC \retained \MyDimension10pt \the\MyDimension
\NC \retained\advance\MyDimension10pt \the\MyDimension
\NC \retained\advance\MyDimension10pt \the\MyDimension \NC \NR
\NC \retained \MyDimension10pt \the\MyDimension
\NC \retained\advance\MyDimension10pt \the\MyDimension
\NC \retained\advance\MyDimension10pt \the\MyDimension \NC \NR
```

**\stoptabulate**

```
10.0pt 20.0pt 30.0pt
10.0pt 20.0pt 30.0pt
```

So what is the difference with the global approach? Say we have these two buffers:

**\startbuffer[one]**

```
\global\MyDimension\zeropoint
\framed {
  \framed {\global\advance\MyDimension10pt \the\MyDimension}
  \framed {\global\advance\MyDimension10pt \the\MyDimension}
}
\framed {
  \framed {\global\advance\MyDimension10pt \the\MyDimension}
  \framed {\global\advance\MyDimension10pt \the\MyDimension}
}
```

**\stopbuffer****\startbuffer[two]**

```
\global\MyDimension\zeropoint
\framed {
  \framed {\global\advance\MyDimension10pt \the\MyDimension}
  \framed {\global\advance\MyDimension10pt \the\MyDimension}
  \getbuffer[one]
}
\framed {
```

```

    \framed {\global\advance\MyDimension10pt \the\MyDimension}
    \framed {\global\advance\MyDimension10pt \the\MyDimension}
    \getbuffer[one]
}
\stopbuffer

```

Typesetting the second buffer gives us:

10.0pt	20.0pt	10.0pt	20.0pt	30.0pt	40.0pt
50.0pt	60.0pt	10.0pt	20.0pt	30.0pt	40.0pt

When we want to have these entities independent and not use different variables, the global settings bleeding from one into the other entity is messy. Therefore we can use this:

```

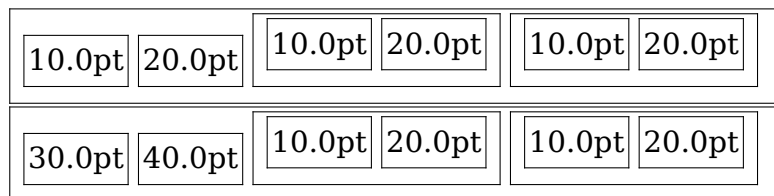
\startbuffer[one]
  \constrained\MyDimension\zeropoint
  \framed {
    \framed {\retained \MyDimension10pt \the\MyDimension}
    \framed {\retained\advance\MyDimension10pt \the\MyDimension}
  }
  \framed {
    \framed {\retained \MyDimension10pt \the\MyDimension}
    \framed {\retained\advance\MyDimension10pt \the\MyDimension}
  }
\stopbuffer

\startbuffer[two]
  \constrained\MyDimension\zeropoint
  \framed {
    \framed {\retained\advance\MyDimension10pt \the\MyDimension}
    \framed {\retained\advance\MyDimension10pt \the\MyDimension}
    \getbuffer[one]
  }
  \framed {
    \framed {\retained\advance\MyDimension10pt \the\MyDimension}
    \framed {\retained\advance\MyDimension10pt \the\MyDimension}
    \getbuffer[one]
  }
}

```

## `\stopbuffer`

Now we get this:



The `\constrained` prefix makes sure that we have a stack entry, without being clever with respect to the current value. Then the `\retained` prefix can do its work reliably and avoid pushing the old value on the stack. Without the constrain it gets a bit unpredictable because then it all depends on where further up the chain the value was put on the stack. Of course one can argue that we should not have the “save stack redundant entries optimization” but that's not going to be removed.

## 7.5 Files

Although it doesn't really fit in this chapter, here are some hooks into processing files:

```

Hello World!\atendoffiled           {\writestatus{FILE}{ATEOF B #1}}\par
Hello World!\atendoffiled           {\writestatus{FILE}{ATEOF A #1}}\par
Hello World!\atendoffiled reverse  {\writestatus{FILE}{ATEOF C #1}}\par
Hello World!\begingroup\sl \atendoffiled {\endgroup}\par

```

Inside a file you can register tokens that will be expanded when the file ends. You can also do that beforehand using a variant of the `\input` primitive:

```
\eofinput {\writestatus{FILE}{DONE}} {thatfile.tex}
```

This feature is mostly there for consistency with the hooks into groups and paragraphs but also because `\everyeof` is kind of useless given that one never knows beforehand if a file loads another file. The hooks mentioned above are bound to the current file.

## 7.5 Colofon

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## **8 Security**

low level

TEX

security

## Contents

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### 8.1 Preamble

Here I will discuss a moderate security subsystem of LuaMetaT<sub>E</sub>X and therefore ConT<sub>E</sub>Xt LMTX. This is not about security in the sense of the typesetting machinery doing harm to your environment, but more about making sure that a user doesn't change the behavior of the macro package in ways that introduce interference and thereby unwanted side effect. It's all about protecting macros.

This is all very experimental and we need to adapt the ConT<sub>E</sub>Xt source code to this. Actually that will happen a few times because experiments trigger that. It might take a few years before the security model is finalized and all files are updated accordingly. There are lots of files and macros involved. In the process the underlying features in the engine might evolve.

### 8.2 Flags

Before we go into the security levels we see what flags can be set. The T<sub>E</sub>X language has a couple of so called prefixes that can be used when setting values and defining macros. Any engine that has traditional T<sub>E</sub>X with  $\varepsilon$ -T<sub>E</sub>X extensions can do this:

```

                \def\foo{foo}
\global        \def\foo{foo}
\global\protected\def\foo{foo}

```

And LuaMetaT<sub>E</sub>X adds another one:

```

                \tolerant      \def\foo{foo}
\global\tolerant      \def\foo{foo}
\global\tolerant\protected\def\foo{foo}

```

What these prefixes do is discussed elsewhere. For now it is enough to know that the two optional prefixes `\protected` and `\tolerant` make for four distinctive cases of macro calls.

But there are more prefixes:

---

frozen	a macro that has to be redefined in a managed way
permanent	a macro that had better not be redefined
primitive	a primitive that normally will not be adapted
immutable	a macro or quantity that cannot be changed, it is a constant
mutable	a macro that can be changed no matter how well protected it is

---

instance	a macro marked as (for instance) be generated by an interface
----------	---

---

noaligned	the macro becomes acceptable as <code>\noalign</code> alias
-----------	---

---

overloaded	when permitted the flags will be adapted
enforced	all is permitted (but only in zero mode or ini mode)
aliased	the macro gets the same flags as the original

---

These prefixed set flags to the command at hand which can be a macro but basically any control sequence.

To what extent the engine will complain when a property is changed in a way that violates the above depends on the parameter `\overloadmode`. When this parameter is set to zero no checking takes place. More interesting are values larger than zero. If that is the case, when a control sequence is flagged as mutable, it is always permitted to change. When it is set to immutable one can never change it. The other flags determine the kind of checking done. Currently the following overload values are used:

		<b>immutable</b>	<b>permanent</b>	<b>primitive</b>	<b>frozen</b>	<b>instance</b>
1	warning	*	*	*		
2	error	*	*	*		
3	warning	*	*	*	*	
4	error	*	*	*	*	
5	warning	*	*	*	*	*
6	error	*	*	*	*	*

The even values (except zero) will abort the run. In `ConTEXt` we plug in a callback that deals with the messages. A value of 255 will freeze this parameter. At level five and above the instance flag is also checked but no drastic action takes place. We use this to signal to the user that a specific instance is redefined (of course the definition macros can check for that too).

So, how does it work. The following is okay:

```
\def\MacroA{A}
\def\MacroB{B}
```



```
\let\MyMacro\MacroA
\let\MyMacro\MacroB
```

The first two macros are ordinary ones, and the last two lines just create an alias. Such an alias shares the definition, but when for instance `\MacroA` is redefined, its new meaning will not be reflected in the alias.

```
\permanent\protected\def\MacroA{A}
\permanent\protected\def\MacroB{B}
\let\MyMacro\MacroA
\let\MyMacro\MacroB
```

This also works, because the `\let` will create an alias with the protected property but it will not take the permanent property along. For that we need to say:

```
\permanent\protected\def\MacroA{A}
\permanent\protected\def\MacroB{B}
\permanent\let\MyMacro\MacroA
\permanent\let\MyMacro\MacroB
```

or, when we want to copy all properties:

```
\permanent\protected\def\MacroA{A}
\permanent\protected\def\MacroB{B}
\aliased\let\MyMacro\MacroA
\aliased\let\MyMacro\MacroB
```

However, in `ConTEXt` we have commands that we like to protect against overloading but at the same time have a different meaning depending on the use case. An example is the `\NC` (next column) command that has a different implementation in each of the table mechanisms.

```
\permanent\protected\def\NC_in_table {...}
\permanent\protected\def\NC_in_tabulate{...}
\aliased\let\NC\NC_in_table
\aliased\let\NC\NC_in_tabulate
```

Here the second aliasing of `\NC` fails (assuming of course that we enabled overload checking). One can argue that grouping can be used but often no grouping takes place when we redefine on the fly. Because `frozen` is less restrictive than `primitive` or `permanent`, and of course `immutable`, the next variant works:

```
\frozen\protected\def\NC_in_table {...}
```

```
\frozen\protected\def\NC_in_tabulate{...}
\overloaded\let\NC\NC_in_table
\overloaded\let\NC\NC_in_tabulate
```

However, in practice, as we want to keep the overload checking, we have to do:

```
\frozen\protected\def\NC_in_table  {...}
\frozen\protected\def\NC_in_tabulate{...}
\overloaded\frozen\let\NC\NC_in_table
\overloaded\frozen\let\NC\NC_in_tabulate
```

or use `\aliased`, but there might be conflicting permissions. This is not that nice, so there is a kind of dirty trick possible. Consider this:

```
\frozen\protected\def\NC_in_table  {...}
\frozen\protected\def\NC_in_tabulate{...}
\def\setNCintable  {\enforced\let\frozen\let\NC\NC_in_table}
\def\setNCintabulate{\enforced\let\frozen\let\NC\NC_in_tabulate}
```

When we're in so called `initex` mode or when the overload mode is zero, the `\enforced` prefix is internalized in a way that signals that the follow up is not limited by the overload mode and permissions. This definition time binding mechanism makes it possible to use permanent macros that users cannot redefine, but existing macros can, unless of course they tweak the mode parameter.

Now keep in mind that users can always cheat but that is intentional. If you really want to avoid that you can set the overload mode to 255 after which it cannot be set any more. However, it can be useful to set the mode to zero (or some warning level) when foreign macro packages are used.

## 8.3 Complications

One side effect of all this is that all those prefixes can lead to more code. On the other hand we save some due to the extended macro argument handling features. When you take the size of the format file as reference, in the end we get a somewhat smaller file. Every token that you add or remove gives a 8 bytes difference. The extra overhead that got added to the engine is compensated by the fact that some macro implementations can be more efficient. In the end, in spite of these new features and the more extensive testing of flags performance is about the same.<sup>10</sup>

<sup>10</sup> And if you wonder about memory, by compacting the used (often scattered) token memory before dumping I manages to save some 512K on the format file, so often the loss and gain are somewhere else.

## 8.4 Introspection

In case you want to get some details about the properties of a macro, you can check its meaning. The full variant shows all of them.

`% a macro with two optional arguments with optional spacing in between:`

```
\permanent\tolerant\protected\def\MyFoo[#1]#*[#2]{(#1)(#2)}
```

```
\meaningless\MyFoo\par
```

```
\meaning \MyFoo\par
```

```
\meaningfull\MyFoo\par
```

```
[#1]#*[#2]->(#1)(#2)
```

```
tolerant protected macro:[#1]#*[#2]->(#1)(#2)
```

```
permanent tolerant protected macro:[#1]#*[#2]->(#1)(#2)
```

## 8.4 Colofon

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## **9 Characters**

low level

TEX

characters

## Contents

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## 9.1 Introduction

This explanation is part of the low level manuals because in practice users will not have to deal with these matters in MkIV and even less in LMTX. You can skip to the last section for commands.

## 9.2 History

If we travel back in time to when T<sub>E</sub>X was written we end up in eight bit character universe. In fact, the first versions assumed seven bits, but for comfortable use with languages other than English that was not sufficient. Support for eight bits permits the usage of so called code pages as supported by operating systems. Although ascii input became kind of the standard soon afterwards, the engine can be set up for different encodings. This is not only true for T<sub>E</sub>X, but for many of its companions, like MetaFont and therefore MetaPost.<sup>11</sup>

Core components of a T<sub>E</sub>X engine are hyphenation of words, applying inter-character kerns and build ligatures. In traditional T<sub>E</sub>X engines those processes are interwoven into the par builder but in LuaT<sub>E</sub>X these are separate stages. The original approach is the reason that there is a relation between the input encoding and the font encoding: the character in the input is the slot used in a reference to a glyph. When producing the final result (e.g. pdf) there can also be a mapping to an index in a font resource.

```
input A [tex ->] font slot A [backend ->] glyph index A
```

The mapping that T<sub>E</sub>X does is normally one-to-one but an input character can undergo some transformation. For instance a character beyond ascii 126 can be made active and expand to some character number that then becomes the font slot. So, it is the

---

<sup>11</sup> This remapping to an internal representation (e.g. ebcdic) is not present in LuaT<sub>E</sub>X where we assume utf8 to be the input encoding. The MetaPost library that comes with LuaT<sub>E</sub>X still has that code but in LuaMetaT<sub>E</sub>X it's gone. There one can set up the machinery to be utf8 aware too.

expansion (or meaning) of a character that end up as numeric reference in the glyph node. Virtual fonts can introduce yet another remapping but that's only visible in the backend.

Actually, in Lua $\TeX$  the same happens but in practice there is no need to go active because (at least in Con $\TeX$ t) we assume a Unicode path so there the font slot is the Unicode got from the utf8 input.

In the eight bit universe macro packages (have to) provide all kind of means to deal with (in the perspective of English) special characters. For instance, `\"a` would put a diaeresis on top of the a or even better, refer to a character in the encoding that the chosen font provides. Because there are some limitations of what can go in an eight bit font, and because in different countries the used  $\TeX$  fonts evolved kind of independent, we ended up with quite some different variants of fonts. It was only with the Latin Modern project that this became better. Interesting is that when we consider the fact that such a font has often also hardly used symbols (like registered or copyright) coming up with an encoding vector that covers most (latin based) European languages (scripts) is not impossible<sup>12</sup> Special symbols could simply go into a dedicated font, also because these are always accessed via a macro so who cares about the input. It never happened.

Keep in mind that when utf8 is used with eight bit engines, Con $\TeX$ t will convert sequences of characters into a slot in a font (depending on the font encoding used which itself depends on the coverage needed). For this every first (possible) byte of a multi-byte utf sequence is an active character, which is no big deal because these are outside the ascii range. Normal ascii characters are single byte utf sequences and fall through without treatment.

Anyway, in Con $\TeX$ t MkII we dealt with this by supporting mixed encodings, depending on the (local) language, referencing the relevant font. It permits users to enter the text in their preferred input encoding and also get the words properly hyphenated. But we can leave these MkII details behind.

### 9.3 The heritage

In MkIV we got rid of input and font encodings, although one can still load files in a specific code page.<sup>13</sup> We also kept the means to enter special characters, if only because

---

<sup>12</sup> And indeed in the Latin Modern project we came up with one but it was already to late for it to become popular.

<sup>13</sup> I'm not sure if users ever depend on an input encoding different from utf8.

text editors seldom support(ed) a wide range of visual editing of those. This is why we still have

```
\"u \^a \v{s} \AE \ij \eacute \oslash
```

and many more. The ones with one character names are rather common in the T<sub>E</sub>X community but it is definitely a weird mix of symbols. The next two are kind of outdated: in these days you delegate that to the font handler, where turning them into ‘single’ character references depends on what the font offers, how it is set up with respect to (for instance) ligatures, and even might depend on language or script.

The ones with the long names partly are tradition, but as we have a lot of them, in MkII they actually serve a purpose. These verbose names are used in the so called encoding vectors and are part of the utf expansion vectors. They are also used in labels so that we have a good indication if what goes in there: remember that in those times editors often didn't show characters, unless the font for display had them, or the operating system somehow provided them from another font. These verbose names are used for latin, greek and cyrillic and for some other scripts and symbols. They take up quite a bit of hash space and the format file.<sup>14</sup>

## 9.4 The LMTX approach

In the process of tagging all (public) macros in LMTX (which happened in 2020-2021) I wondered if we should keep these one character macros, the references to special characters and the verbose ones. When asked on the mailing list it became clear that users still expect the short ones to be present, often just because old bibT<sub>E</sub>X files are used that might need them. However, in MkIV and LMTX we load bibT<sub>E</sub>X files in a way that turn these special character references into proper utf8 input so it makes a weak argument. Anyway, although they could go, for now we keep them because users expect them. However, in LMTX the implementation is somewhat different now, a bit more efficient in terms of hash and memory, potentially a bit less efficient in runtime, but no one will notice that.

A new command has been introduced, the very short `\chr`.

```
\chr {a} \chr {a} \chr {a}
\chr {`a} \chr {'a} \chr {"a}
\chr {a acute} \chr {a grave} \chr {a umlaut}
```

<sup>14</sup> In MkII we have an abstract front-end with respect to encodings and also an abstract backend with respect to supported drivers but both approaches no longer make sense today.



```
\chr {aacute} \chr {agrave} \chr {aumlaut}
```

In the first line the composed character using two characters, a base and a so called mark. Actually, one doesn't have to use `\chr` in that case because `ConTEXt` does already collapse characters for you. The second line looks like the shortcuts `\``, `\'` and `\"`. The third and fourth lines could eventually replace the more symbolic long names, if we feel the need. Watch out: in Unicode input the marks come *after*.

```
à á ä
à á ä
á à a~mla~t
á à a~mla~t
```

Currently the repertoire is somewhat limited but it can be easily be extended. It all depends on user needs (doing Greek and Cyrillic for instance). The reason why we actually save code deep down is that the helpers for this have always been there.<sup>15</sup>

The `\"` commands are now just aliases to more verbose and less hackery looking macros:

<code>\withgrave</code>	à	<code>\`</code>	à
<code>\withacute</code>	á	<code>\'</code>	á
<code>\withcircumflex</code>	â	<code>\^</code>	â
<code>\withtilde</code>	ã	<code>\~</code>	ã
<code>\withmacron</code>	ā	<code>\=</code>	ā
<code>\withbreve</code>	ě	<code>\u</code>	ě
<code>\withdotaccent</code>	ċ	<code>\.</code>	.c
<code>\withdiaeresis</code>	ë	<code>\"</code>	ë
<code>\withring</code>	ŭ	<code>\r</code>	ŭ
<code>\withhungarumlaut</code>	ű	<code>\H</code>	ű
<code>\withcaron</code>	ě	<code>\v</code>	ě
<code>\withcedilla</code>	ç	<code>\c</code>	ç
<code>\withogonek</code>	ę	<code>\k</code>	ę

Not all fonts have these special characters. Most natural is to have them available as precomposed single glyphs, but it can be that they are just two shapes with the marks anchored to the base. It can even be that the font somehow overlays them, assuming (roughly) equal widths. The compose font feature in `ConTEXt` normally can handle most well.

<sup>15</sup> So if needed I can port this approach back to MkIV, but for now we keep it as is because we then have a reference.

An occasional ugly rendering doesn't matter that much: better have something than nothing. But when it's the main language (script) that needs them you'd better look for a font that handles them. When in doubt, in ConTEXt you can enable checking:

<b>command</b>	<b>equivalent to</b>
<code>\checkmissingcharacters</code>	<code>\enabletrackers[fonts.missing]</code>
<code>\removemissingcharacters</code>	<code>\enabletrackers[fonts.missing=remove]</code>
<code>\replacemissingcharacters</code>	<code>\enabletrackers[fonts.missing=replace]</code>
<code>\handlemissingcharacters</code>	<code>\enabletrackers[fonts.missing={decompose,replace}]</code>

The decompose variant will try to turn a composed character into its components so that at least you get something. If that fails it will inject a replacement symbol that stands out so that you can check it. The console also mentions missing glyphs. You don't need to enable this by default<sup>16</sup> but you might occasionally do it when you use a font for the first time.

In LMTX this mechanism has been upgraded so that replacements follow the shape and are actually real characters. The decomposition has not yet been ported back to MkIV.

The full list of commands can be queried when a tracing module is loaded:

```
\usemodule[s][characters-combinations]
```

```
\showcharactercombinations
```

We get this list:

acute	U+00301	´	<code>\withacute</code>
breve	U+00306	˘	<code>\withbreve</code>
caron	U+0030C	ˇ	<code>\withcaron</code>
caron below	U+0032C	ˇ̣	<code>\withcaronbelow</code>
cedilla	U+00327	¸	<code>\withcedilla</code>
circumflex	U+00302	ˆ	<code>\withcircumflex</code>
circumflex below	U+0032D	ˆ̣	<code>\withcircumflexbelow</code>
comma below	U+00326	¸	<code>\withcommabelow</code>
diaeresis	U+00308	¨	<code>\withdiaeresis</code>
dieresis	U+00308	¨	<code>\withdieresis</code>
dot	U+00307	·	<code>\withdot</code>
dot below	U+00323	·̣	<code>\withdotbelow</code>
double acute	U+0030B	˝	<code>\withdoubleacute</code>

<sup>16</sup> There is some overhead involved here.

double grave	U+0030F	“	\withdoublegrave
double vertical line	U+0030E	”	\withdoubleverticalline
grave	U+00300	`	\withgrave
hook	U+00309	´	\withhook
hook below	U+1FA9D		\withhookbelow
hungarumlaut	U+0030B	˝	\withhungarumlaut
inverted breve	U+00311	˘	\withinvertedbreve
line	U+00304	-	\withline
line below	U+00331	_	\withlinebelow
macron	U+00304	-	\withmacron
macron below	U+00331	_	\withmacronbelow
middle dot	U+000B7	·	\withmiddledot
ogonek	U+00328	˛	\withogonek
overline	U+00305	—	
ring	U+0030A	°	\withring
ring below	U+00325	◌◌	\withringbelow
slash	U+0002F	/	\withslash
stroke	U+0002F	/	\withstroke
tilde	U+00303	~	\withtilde
tilde below	U+00330	˜	\withtildebelow
vertical line	U+0030D	¦	\withverticalline

Some combinations are special for ConTEXt because Unicode doesn't specify decomposition for all composed characters.

## 9.5 spaces

The engine has no real concept of a space. When the input has one it is turned into a glue, likely with some stretch and shrink. When `\nospaces` is set to one, no glue will be inserted. A value of two will inject a zero width glue. When set to three a glyph will be inserted with the character code set by `\spacechar`.

```
\nospaces\plusthree
\spacechar\underscoreasciicode
\hccode\underscoreasciicode\underscoreasciicode
Where are the spaces?
```

The `hccode` tells the machinery that the underscore is a valid word separator (think compound words).

```
Where_are_the_spaces?
```

## 9.5 Colofon

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## **10 Scope**

low level

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scope

## Contents

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### 10.1 Introduction

When I visited the file where register allocations are implemented I wondered to what extent it made sense to limit allocation to global instances only. This chapter deals with this phenomena.

### 10.2 Registers

In  $\TeX$  definitions can be local or global. Most assignments are local within a group. Registers and definitions can be assigned global by using the `\global` prefix. There are also some properties that are global by design, like `\prevdepth`. A mixed breed are boxes. When you tweak its dimensions you actually tweak the current box, which can be an outer level. Compare:

```
\scratchcounter = 1
here the counter has value 1
\begingroup
  \scratchcounter = 2
  here the counter has value 2
\endgroup
here the counter has value 1
```

with:

```
\setbox\scratchbox=\hbox{}
here the box has zero width
\begingroup
  \wd\scratchbox=10pt
  here the box is 10pt wide
\endgroup
here the box is 10pt wide
```

It all makes sense so a remark like “Assignments to box dimensions are always global” are sort of confusing. Just look at this:

```

\setbox\scratchbox=\hbox to 20pt{}
here the box is \the\wd\scratchbox\ wide\par
\begingroup
  \setbox\scratchbox=\hbox{}
  here the box is \the\wd\scratchbox\ wide\par
  \begingroup
    \wd\scratchbox=15pt
    here the box is \the\wd\scratchbox\ wide\par
  \endgroup
  here the box is \the\wd\scratchbox\ wide\par
\endgroup
here the box is \the\wd\scratchbox\ wide\par

here the box is 20.0pt wide
here the box is 0.0pt wide
here the box is 15.0pt wide
here the box is 15.0pt wide
here the box is 20.0pt wide

```

If you don't think about it, what happens is what you expect. Now watch the next variant:

The `\global` is only effective for the current box. It is good to realize that when we talk registers, the box register behaves just like any other register but the manipulations happen to the current one.

```

\setbox\scratchbox=\hbox to 20pt{}
here the box is \the\wd\scratchbox\ wide\par
\begingroup
  \setbox\scratchbox=\hbox{}
  here the box is \the\wd\scratchbox\ wide\par
  \begingroup
    \global\wd\scratchbox=15pt
    here the box is \the\wd\scratchbox\ wide\par
  \endgroup
  here the box is \the\wd\scratchbox\ wide\par
\endgroup
here the box is \the\wd\scratchbox\ wide\par

here the box is 20.0pt wide
here the box is 0.0pt wide
here the box is 15.0pt wide

```



here the box is 15.0pt wide

here the box is 20.0pt wide

```
\scratchdimen=20pt
```

here the dimension is `\the\scratchdimen\par`

```
\begingroup
```

```
\scratchdimen=0pt
```

here the dimension is `\the\scratchdimen\par`

```
\begingroup
```

```
\global\scratchdimen=15pt
```

here the dimension is `\the\scratchdimen\par`

```
\endgroup
```

here the dimension is `\the\scratchdimen\par`

```
\endgroup
```

here the dimension is `\the\scratchdimen\par`

here the dimension is 20.0pt

here the dimension is 0.0pt

here the dimension is 15.0pt

here the dimension is 15.0pt

here the dimension is 15.0pt

## 10.3 Allocation

The plain  $\text{\TeX}$  format has set some standards and one of them is that registers are allocated with `\new...` commands. So we can say:

```
\newcount\mycounta
```

```
\newdimen\mydimena
```

These commands take a register from the pool and relate the given name to that entry. In  $\text{\ConTeXt}$  we have a bunch of predefined scratch registers for general use, like:

```
scratchcounter      : \meaningfull\scratchcounter
scratchcounterone  : \meaningfull\scratchcounterone
scratchcountertwo  : \meaningfull\scratchcountertwo
scratchdimen       : \meaningfull\scratchdimen
scratchdimenone    : \meaningfull\scratchdimenone
scratchdimentwo    : \meaningfull\scratchdimentwo
```

The meaning reveals what these are:

### Allocation

```

scratchcounter : global constant integer 1026
scratchcounterone : global constant integer 0
scratchcountertwo : global constant integer 0
scratchdimen : global constant dimension 15.0pt
scratchdimenone : global constant dimension 0.0pt
scratchdimentwo : global constant dimension 0.0pt

```

You can use the numbers directly but that is a bad idea because they can clash! In the original  $\TeX$  engine there are only 256 registers and some are used by the engine and the core of a macro package itself, so that leaves a little amount for users. The  $\varepsilon\text{-}\TeX$  extension lifted that limitation and bumped to 32K and Lua $\TeX$  upped that to 64K. One could go higher but what makes sense? These registers are taking part of the fixed memory slots because that makes nested (grouped) usage efficient and access fast. The number you see above is deduced from the so called command code (here indicated by `\count`) and an index encoded in the same token. So, `\scratchcounter` takes a single token contrary to the verbose `\count257` that takes four tokens where the number gets parsed every time it is needed. But those are details that a user can forget.

As mentioned, commands like `\newcount\foo` create a global control sequence `\foo` referencing a counter. You can locally redefine that control sequence unless in LuaMeta- $\TeX$  you have so called overload mode enabled. You can do local or global assignments to these registers.

```

\scratchcounter = 123
\begingroup
  \scratchcounter = 456
  \begingroup
    \global\scratchcounter = 789
  \endgroup
\endgroup

```

And in both cases count register 257 is set. When an assignment is global, all current values to that register get the same value. Normally this is all quite transparent: you get what you ask for. However the drawback is that as a user you cannot know what variables are already defined, which means that this will fail (that is: it will issue a message):

```
\newcount\scratchcounter
```

as will the second line in:

```
\newcount\myscratchcounter
```

```
\newcount\myscratchcounter
```

In ConT<sub>E</sub>Xt the scratch registers are visible but there are lots of internally used ones are protected from the user by more obscure names. So what if you want to use your own register names without ConT<sub>E</sub>Xt barking to you about not being able to define it. This is why in LMTX (and maybe some day in MkIV) we now have local definitions:

```
\begingroup
  \newlocaldimen\mydimena      \mydimena1\onpoint
  \newlocaldimen\mydimenb      \mydimenb2\onpoint
  (\the\mydimena,\the\mydimenb)
  \begingroup
    \newlocaldimen\mydimena      \mydimena3\onpoint
    \newlocaldimen\mydimenb      \mydimenb4\onpoint
    \newlocaldimen\mydimenc      \mydimenc5\onpoint
    (\the\mydimena,\the\mydimenb,\the\mydimenc)
    \begingroup
      \newlocaldimen\mydimena      \mydimena6\onpoint
      \newlocaldimen\mydimenb      \mydimenb7\onpoint
      (\the\mydimena,\the\mydimenb)
    \endgroup
    \newlocaldimen\mydimend      \mydimend8\onpoint
    (\the\mydimena,\the\mydimenb,\the\mydimenc,\the\mydimend)
  \endgroup
  (\the\mydimena,\the\mydimenb)
\endgroup
```

The allocated registers get zero values but you can of course set them to any value that fits their nature:

```
(1.0pt,2.0pt)
(3.0pt,4.0pt,5.0pt)
(6.0pt,7.0pt)
(3.0pt,4.0pt,5.0pt,8.0pt)
(1.0pt,2.0pt)
```

You can also use the next variant where you also pass the initial value:

```
\begingroup
  \setnewlocaldimen\mydimena      1\onpoint
  \setnewlocaldimen\mydimenb      2\onpoint
  (\the\mydimena,\the\mydimenb)
```

```

\begingroup
  \setnewlocaldimen\mydimena 3\onepoint
  \setnewlocaldimen\mydimenb 4\onepoint
  \setnewlocaldimen\mydimenc 5\onepoint
  (\the\mydimena,\the\mydimenb,\the\mydimenc)
\begingroup
  \setnewlocaldimen\mydimena 6\onepoint
  \setnewlocaldimen\mydimenb 7\onepoint
  (\the\mydimena,\the\mydimenb)
\endgroup
  \setnewlocaldimen\mydimend 8\onepoint
  (\the\mydimena,\the\mydimenb,\the\mydimenc,\the\mydimend)
\endgroup
(\the\mydimena,\the\mydimenb)
\endgroup

```

So, again we get:

```

(1.0pt,2.0pt)
(3.0pt,4.0pt,5.0pt)
(6.0pt,7.0pt)
(3.0pt,4.0pt,5.0pt,8.0pt)
(1.0pt,2.0pt)

```

When used in the body of the macro there is of course a little overhead involved in the repetitive allocation but normally that can be neglected.

## 10.4 Files

When adding these new allocators I also wondered about the read and write allocators. We don't use them in ConT<sub>E</sub>Xt but maybe users like them, so let's give an example and see what more demands they have:

```

\integerdef\StartHere\numexpr\inputlineno+2\relax
\starthiding
SOME LINE 1
SOME LINE 2
SOME LINE 3
SOME LINE 4
\stophiding
\integerdef\StopHere\numexpr\inputlineno-2\relax

```

```

\begingroup
  \newlocalread\myreada
  \immediate\openin\myreada {lowlevel-scope.tex}
  \dostepwiserecurse{\StopHere}{\StartHere}{-1}{
    \readline\myreada line #1 to \scratchstring #1 : \scratchstring \par
  }
  \blank
  \dostepwiserecurse{\StartHere}{\StopHere}{1}{
    \read    \myreada line #1 to \scratchstring #1 : \scratchstring \par
  }
  \immediate\closein\myreada
\endgroup

```

Here, instead of hard coded line numbers we used the stored values. The optional line keyword is a LMTX speciality.

```

281 : SOME LINE 4
280 : SOME LINE 3
279 : SOME LINE 2
278 : SOME LINE 1

```

```

278 : SOME LINE 1
279 : SOME LINE 2
280 : SOME LINE 3
281 : SOME LINE 4

```

Actually an application can be found in a small (demonstration) module:

```

\usemodule[system-readers]

```

This provides the code for doing this:

```

\startmarkedlines[test]
SOME LINE 1
SOME LINE 2
SOME LINE 3
\stopmarkedlines

```

```

\begingroup
  \newlocalread\myreada
  \immediate\openin\myreada {\markedfilename{test}}
  \dostepwiserecurse{\lastmarkedline{test}}{\firstmarkedline{test}}{-1}{
    \readline\myreada line #1 to \scratchstring #1 : \scratchstring \par
  }

```

```
}  
\immediate\closein\myreada  
\endgroup
```

As you see in these examples, we can locally define a read channel without getting a message about it already being defined.

## 10.4 Colofon

Author	Hans Hagen
ConT <sub>E</sub> Xt	2025.07.04 21:26
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Support	<a href="http://www.pragma-ade.com">www.pragma-ade.com</a> <a href="http://contextgarden.net">contextgarden.net</a> <a href="mailto:ntg-context@ntg.nl">ntg-context@ntg.nl</a>

## **11 Paragraphs**

low level

TEX

paragraphs



## Contents

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### 11.1 Introduction

This manual is mostly discussing a few low level wrappers around low level  $\text{T}_{\text{E}}\text{X}$  features. Its writing is triggered by an update to the MetaFun and LuaMetaFun manuals where we mess a bit with shapes. It gave a good reason to also cover some more paragraph related topics but it might take a while to complete. Remind me if you feel that takes too much time.

Because paragraphs and their construction are rather central to  $\text{T}_{\text{E}}\text{X}$ , you can imagine that the engine exposes dealing with them. This happens via commands (primitives) but only when it's robust. Then there are callbacks, and some provide detailed information about what we're dealing with. However, intercepting node lists can already be hairy and we do that a lot in  $\text{ConT}_{\text{E}}\text{Xt}$ . Intercepting and tweaking paragraph properties is even more tricky, which is why we try to avoid that in the core. But . . . in the following sections you will see that there are actually a couple of mechanism that do so. Often new features like this are built in stepwise and enabled locally for a while and when they seem okay they get enabled by default.<sup>17</sup>

---

<sup>17</sup> For this we have `\enableexperiments` which one can use in `cont-loc.mkxl` or `cont-exp.mkxl`, files that are loaded runtime when on the system. When you use them, make sure they don't interfere; they are not part of the updates, contrary to `cont-new.mkxl`.

## 11.2 Paragraphs

Before we demonstrate some trickery, let's see what a paragraph is. Normally a document source is formatted like this:

```
some text (line 1)
some text (line 2)

some more test (line 1)
some more test (line 2)
```

There are two blocks of text here separated by an empty line and they become two paragraphs. Unless configured otherwise an empty line is an indication that we end a paragraph. You can also explicitly do that:

```
some text (line 1)
some text (line 2)
\par
some more test (line 1)
some more test (line 2)
```

When  $\text{\TeX}$  starts a paragraph, it actually also does something think of:

```
[\the\everypar]some text      (line 1) some text      (line 2) \par
[\the\everypar]some more test (line 1) some more test (line 2) \par
```

or more accurate:

```
[\the\everypar]some text      some text      \par
[\the\everypar]some more test some more test \par
```

because the end-of-line character has become a space. As mentioned, an empty line is actually the end of a paragraph. But in LuaMeta $\text{\TeX}$  we can cheat a bit. If we have this:

```
line 1

line 2
```

We can do this (watch how we need to permit overloading a primitive when we have enabled `\overloadmode`):

```
\pushoverloadmode
\def\linepar{\removeunwantedspaces !\ignorespaces}
\popoverloadmode
```

line 1

line 2

This comes out as:

line 1

line 2

I admit that since it got added (as part of some cleanup halfway the overhaul of the engine) I never saw a reason to use it, but it is a cheap feature. The `\linepar` primitive is undefined (`\undefined`) by default so no user sees it anyway. Just don't use it unless maybe for some pseudo database trickery (I considered using it for the database module but it is not needed). In a similar fashion, just don't redefine `\par`: it's asking for troubles and 'not done' in ConT<sub>E</sub>Xt anyway.

Back to reality. In LuaT<sub>E</sub>X we get a node list that starts with a so called `localpar` node and ends with a `\parfillskip`. The first node is prepended automatically. That list travels through the system: hyphenation, applying font properties, break the effectively one line into lines, wrap them and add them to a vertical list, etc. Each stage can be intercepted via callbacks.

When the paragraph is broken into lines hanging indentation or a so called par shape can be applied, and we will see more of that later, here we talk `\par` and show another LuaMetaT<sub>E</sub>X trick:

```
\def\foo{{\bf test:} \ignorepars}
```

```
\foo
```

line

The macro typesets some text and then skips to the next paragraph:

**test:** line

Think of this primitive as being a more powerful variant of `\ignorespaces`. This leaves one aspect: how do we start a paragraph. Technically we need to force T<sub>E</sub>X into so called horizontal mode. When you look at plain T<sub>E</sub>X documents you will notice commands like `\noindent` and `\indent`. In ConT<sub>E</sub>Xt we have more high level variants, for instance we have `\noindentation`.

A robust way to make sure that you get in horizontal mode is using `\dontleavehmode` which is a wink to `\leavevmode`, a command that you should never use in ConT<sub>E</sub>Xt, so

when you come from plain or L<sup>A</sup>T<sub>E</sub>X, it's one of the commands you should wipe from your memory.

When T<sub>E</sub>X starts with a paragraph the `\everypar` token list is expanded and again this is a primitive you should not mess with yourself unless in very controlled situations. If you change its content, you're on your own with respect to interferences and side effects.

One of the things that T<sub>E</sub>X does in injecting the indentation. Even when there is none, it gets added, not as skip but as an empty horizontal box of a certain width. This is easier on the engine when it constructs the paragraph from the one liner: starting with a skip demands a bit more testing in the process (a nice trick so to say). However, in ConT<sub>E</sub>Xt we enable the LuaMetaT<sub>E</sub>X feature that does use a skip instead of a box. It's part of the normalization that is discussed later. Instead of checking for a box with property `indent`, we check for a skip with such property. This is often easier and cleaner.

A bit off topic is the fact that in traditional T<sub>E</sub>X empty lines or `\par` primitives can trigger an error. This has to do with the fact that the program evolved in a time where paper terminals were used and runtime could be excessive. So, in order to catch a possible missing brace, a concept of `\long` macros, permitting `\par` or equivalents in arguments, was introduced as well as not permitting them in for instance `display math`. In ConT<sub>E</sub>Xt MkII most macros that could be sensitive for this were defined as `\long` so that users never had to bother about it and probably were not even aware of it. Right from the start in LuaT<sub>E</sub>X these error-triggers could be disabled which of course we enable in ConT<sub>E</sub>Xt and in LuaMetaT<sub>E</sub>X these features have been removed altogether. I don't think users will complain about this.

If you want to enforce a newline but not a new paragraph you can use the `\crlf` command. When used on its own it will produce an empty line. Don't use this to create whitespace between lines.

If you want to do something after so called par tokens are seen you can do this:

```
\def\foo{{\bf >>>> }}
\expandafterpars\foo
```

this is a new paragraph ...

```
\expandafterpars\foo
\par\par\par\par
```

this is a new paragraph ...

This not to be confused with `\everypar` which is a token list that  $\TeX$  itself injects before each paragraph (also nested ones).

```
>>>> this is a new paragraph ...
```

```
>>>> this is a new paragraph ...
```

This is typically a primitive that will only be used in macros. You can actually program it using macros: pickup a token, check and push it back when it's not a par equivalent token. The primitive is is just nicer (and easier on the log when tracing is enabled).

## 11.3 Properties

A paragraph is just a collection of lines that result from one input line that got broken. This process of breaking into lines is influenced by quite some parameters. In traditional  $\TeX$  and also in LuaMeta $\TeX$  by default the values that are in effect when the end of the paragraph is met are used. So, when you change them in a group and then ends the paragraph after the group, the values you've set in the group are not used.

However, in LuaMeta $\TeX$  we can optionally store them with the paragraph. When that happens the values current at the start are frozen. You can still overload them but that has to be done explicitly then. The advantage is that grouping no longer interferes with the line break algorithm. The magic primitive is `\snapshotpar` which takes a number made from categories mentioned below:

<b>variable</b>	<b>category</b>	<b>code</b>
<code>\hsize</code>	hsize	0x00000001
<code>\leftskip</code>	skip	0x00000002
<code>\rightskip</code>	skip	0x00000002
<code>\hangindent</code>	hang	0x00000004
<code>\hangafter</code>	hang	0x00000004
<code>\parindent</code>	indent	0x00000008
<code>\parfillleftskip</code>	parfill	0x00000010
<code>\parfillskip</code>	parfill	0x00000010
<code>\parinitleftskip</code>	parfill	0x00000010
<code>\parinitrightskip</code>	parfill	0x00000010
<code>\emergencyleftskip</code>	emergency	0x00800000
<code>\emergencyrightskip</code>	emergency	0x00800000
<code>\adjustspacing</code>	adjust	0x00000020
<code>\protrudechars</code>	protrude	0x00000040
<code>\pretolerance</code>	tolerance	0x00000080

<code>\tolerance</code>	tolerance	0x00000080
<code>\emergencystretch</code>	stretch	0x00000100
<code>\looseness</code>	looseness	0x00000200
<code>\lastlinefit</code>	lastline	0x00000400
<code>\linepenalty</code>	linepenalty	0x00000800
<code>\interlinepenalty</code>	linepenalty	0x00000800
<code>\clubpenalty</code>	clubpenalty	0x00001000
<code>\widowpenalty</code>	widowpenalty	0x00002000
<code>\displaywidowpenalty</code>	displaypenalty	0x00004000
<code>\lefttwindemerits</code>	twindemerits	0x20000000
<code>\righttwindemerits</code>	twindemerits	0x20000000
<code>\brokenpenalty</code>	brokenpenalty	0x00008000
<code>\adjdemerits</code>	demerits	0x00010000
<code>\doublehyphdemerits</code>	demerits	0x00010000
<code>\finalhyphdemerits</code>	demerits	0x00010000
<code>\parshape</code>	shape	0x00020000
<code>\interlinepenalties</code>	linepenalty	0x00000800
<code>\clubpenalties</code>	clubpenalty	0x00001000
<code>\widowpenalties</code>	widowpenalty	0x00002000
<code>\displaywidowpenalties</code>	displaypenalty	0x00004000
<code>\brokenpenalties</code>	brokenpenalty	0x00008000
<code>\orphanpenalties</code>	orphanpenalty	0x00200000
<code>\toddlrpenalties</code>	toddlerpenalty	0x00400000
<code>\fitnessclasses</code>	fitnessclasses	0x40000000
<code>\adjacentdemerits</code>	demerits	0x00010000
<code>\mathleftclass</code>	orphanpenalty	0x00200000
<code>\baselineskip</code>	line	0x00040000
<code>\lineskip</code>	line	0x00040000
<code>\lineskiplimit</code>	line	0x00040000
<code>\adjustspacingstep</code>	adjust	0x00000020
<code>\adjustspacingshrink</code>	adjust	0x00000020
<code>\adjustspacingstretch</code>	adjust	0x00000020
<code>\hyphenationmode</code>	hyphenation	0x00080000
<code>\shapingpenaltiesmode</code>	shapingpenalty	0x00100000
<code>\shapingpenalty</code>	shapingpenalty	0x00100000
<code>\emergencyextrastretch</code>	emergency	0x00800000
<code>\parpasses</code>	parpasses	0x01000000
<code>\linebreakchecks</code>	linebreakchecks	0x10000000
<code>\singlelinepenalty</code>	singlelinepenalty	0x02000000

### Properties

```
\hyphenpenalty      hyphenpenalty    0x04000000
\exhyphenpenalty    exhyphenpenalty  0x08000000
```

As you can see here, there are more paragraph related parameters than in for instance pdf $\TeX$  and Lua $\TeX$  and these are (to be) explained in the LuaMeta $\TeX$  manual. You can imagine that keeping this around with the paragraph adds some extra overhead to the machinery but most users won't notice that because it is compensated by gains elsewhere.

This is pretty low level and there are a bunch of helpers that support this but these are not really user level macros. As with everything  $\TeX$  you can mess around as much as you like, and the code gives plenty of examples but when you do this, you're on your own because it can interfere with Con $\TeX$ t core functionality.

In L $\mathcal{M}\mathcal{T}\mathcal{X}$  taking these snapshots is turned on by default and because it thereby fundamentally influences the par builder, users can run into compatibility issues but in practice there has been no complaints (and this feature has been in use quite a while before this document was written). One reason for users not noticing is that one of the big benefits is probably handled by tricks mentioned on the mailing list. Imagine that you have this:

```
{\bf watch out:} here is some text
```

In this small example the result will be as expected. But what if something magic with the start of a paragraph is done? Like this:

```
\placefigure[left]{A cow!}{\externalfigure[cow.pdf]}
```

```
{\bf watch out:} here is some text ... of course much more is needed to
  get a flow around the figure!
```

The figure will hang at the left side of the paragraph but it is put there when the text starts and that happens inside the bold group. It means that the properties we set in order to get the shape around the figure are lost as soon as we're at 'here is some text' and definitely is wrong when the paragraph ends and the par builder has to use them to get the shape right. We get text overlapping the figure. A trick to overcome this is:

```
\dontleavehmode {\bf watch out:} here is some text ... of course much
  more is needed to get a flow around the figure!
```

where the first macro makes sure we already start a paragraph before the group is entered (using a  $\backslash$ strut also works). It's not nice and I bet users have been bitten by

## Properties

this and by now know the tricks. But, with snapshots such fuzzy hacks are not needed any more! The same is true with this:

```
{\leftskip 1em some text \par}
```

where we had to explicitly end the paragraph inside the group in order to retain the skip. I suppose that users normally use the high level environments so they never had to worry about this. It's also why users probably won't notice that this new mechanism has been active for a while. Actually, when you now change a parameter inside the paragraph its new value will not be applied (unless you prefix it with `\frozen` or `snapshot` it) but no one did that anyway.

## 11.4 Wrapping up

In ConT<sub>E</sub>Xt LMTX we have a mechanism to exercise macros (or content) before a paragraph ends. This is implemented using the `\wrapuppar` primitive. The to be wrapped up material is bound to the current paragraph which in order to get this done has to be started when this primitive is used.

Although the high level interface has been around for a while it still needs a bit more testing (read: use cases are needed). In the few cases where we already use it application can be different because again it relates to snapshots. This because in the past we had to use tricks that also influenced the user interface of some macros (which made them less natural as one would expect). So the question is: where do we apply it in old mechanisms and where not.

*todo: accumulation, interference, where applied, limitations*

## 11.5 Hanging

There are two mechanisms for getting a specific paragraph shape: rectangular hanging and arbitrary shapes. Both mechanisms work top-down. The first mechanism uses a combination of `\hangafter` and `\hangindent`, and the second one depends on `\parshape`. In this section we discuss the rectangular one.

```
\hangafter 4 \hangindent 4cm \samplefile{tufte} \page
\hangafter -4 \hangindent 4cm \samplefile{tufte} \page
\hangafter 4 \hangindent -4cm \samplefile{tufte} \page
\hangafter -4 \hangindent -4cm \samplefile{tufte} \page
```

As you can see in figure 11.1, the four cases are driven by the sign of the values. If you want to hang into the margin you need to use different tricks, like messing with the



`\leftskip`, `\rightskip` or `\parindent` parameters (which then of course can interfere with other mechanisms uses at the same time).

## 11.6 Shapes

In `ConTeXt` we don't use `\parshape` a lot. It is used in for instance side floats but even there not in all cases. It's more meant for special applications. This means that in `MkII` and `MkIV` we don't have some high level interface. However, when `MetaFun` got upgraded to `LuaMetaFun`, and the manual also needed an update, one of the examples in that manual that used shapes also got done differently (read: nicer). And that triggered the arrival of a new low level shape mechanism.

One important property of the `\parshape` mechanism is that it works per paragraph. You define a shape in terms of a left margin and width of a line. The shape has a fixed number of such pairs and when there is more content, the last one is used for the rest of the lines. When the paragraph is finished, the shape is forgotten.<sup>18</sup>

The high level interface is a follow up on the example in the `MetaFun` manual and uses shapes that carry over to the next paragraph. In addition we can cycle over a shape. In this interface shapes are defined using keyword. Here are some examples:

```
\startparagraphshape[test]
  left 1mm right 1mm
  left 5mm right 5mm
\stopparagraphshape
```

This shape has only two entries so the first line will have a 1mm margin while later lines will get 5mm margins. This translates into a `\parshape` like:

```
\parshape 2
  1mm \dimexpr\hsize-1mm\relax
  5mm \dimexpr\hsize-5mm\relax
```

Watch the number 2: it tells how many specification lines follow. As you see, we need to calculate the width.

```
\startparagraphshape[test]
  left 1mm right 1mm
```

---

<sup>18</sup> Not discussed here is a variant that might end up in `LuaMetaTeX` that works with the progression, i.e. takes the height of the content so far into account. This is somewhat tricky because for that to work vertical skips need to be frozen, which is no real big deal but has to be done careful in the code.



**Figure 11.1** Hanging indentation

```
left 5mm right 5mm
repeat
```

**\stopparagraphshape**

This variant will alternate between 1mm and 5mm margins. The repeating feature is translated as follows. Maybe at some point I will introduce a few more options.

```
\parshape 2 options 1
 1mm \dimexpr\hsize-1mm\relax
 5mm \dimexpr\hsize-5mm\relax
```

A shape can have some repetition, and we can save keystrokes by copying the last entry. The resulting `\parshape` becomes rather long.

```
\startparagraphshape[test]
left 1mm right 1mm
left 2mm right 2mm
left 3mm right 3mm
copy 8
left 4mm right 4mm
left 5mm right 5mm
left 5mm hsize 10cm
```

**\stopparagraphshape**

Also watch the `hsize` keyword: we don't calculate the `hsize` from the left and right values but explicitly set it.

```
\startparagraphshape[test]
left 1mm right 1mm
right 3mm
left 5mm right 5mm
repeat
```

**\stopparagraphshape**

When a right keywords comes first the left is assumed to be zero. In the examples that follow we will use a couple of definitions:

```
\startparagraphshape[test]
both 1mm both 2mm both 3mm both 4mm both 5mm both 6mm
both 7mm both 6mm both 5mm both 4mm both 3mm both 2mm
```

**\stopparagraphshape**

```
\startparagraphshape[test-repeat]
```

```
both 1mm both 2mm both 3mm both 4mm both 5mm both 6mm
both 7mm both 6mm both 5mm both 4mm both 3mm both 2mm
repeat
```

```
\stopparagraphshape
```

The last one could also be defines as:

```
\startparagraphshape[test-repeat]
  \rawparagraphshape{test} repeat
\stopparagraphshape
```

In the previous code we already introduced the repeat option. This will make the shape repeat at the engine level when the shape runs out of specified lines. In the application of a shape definition we can specify a method to be used and that determine if the next paragraph will start where we left off and discard afterwards (shift) or that we move the discarded lines up front so that we never run out of lines (cycle). It sounds complicated but just keep in mind that repeat is part of the `\parshape` and act within a paragraph while shift and cycle are applied when a new paragraph is started.

In figure 11.2 you see the following applied:

```
\startshapedparagraph[list=test]
  \dorecurse{8}{\showparagraphshape\samplefile{tufte} \par}
\stopshapedparagraph

\startshapedparagraph[list=test-repeat]
  \dorecurse{8}{\showparagraphshape\samplefile{tufte} \par}
\stopshapedparagraph
```

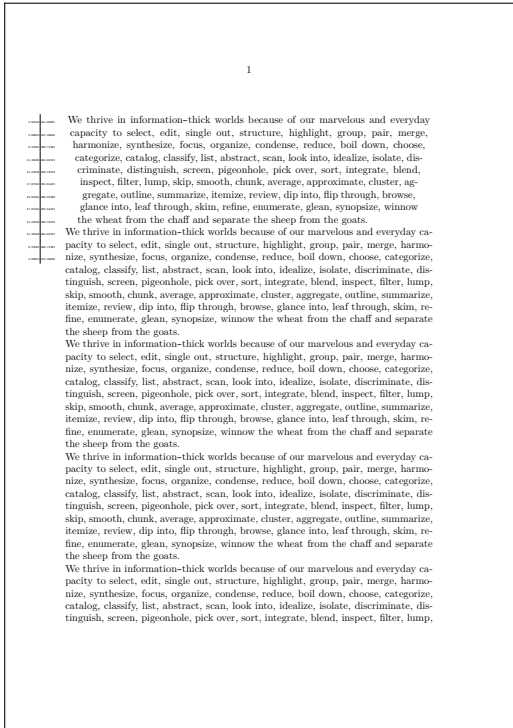
In figure 11.3 we use this instead:

```
\startshapedparagraph[list=test,method=shift]
  \dorecurse{8}{\showparagraphshape\samplefile{tufte} \par}
\stopshapedparagraph
```

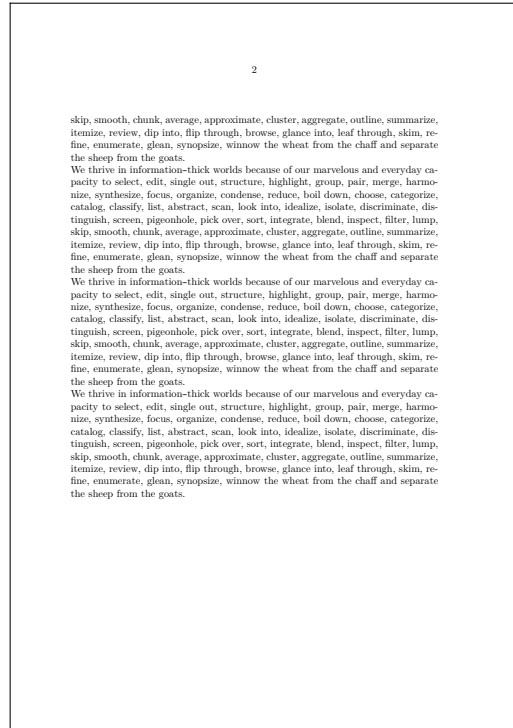
Finally, in figure 11.4 we use:

```
\startshapedparagraph[list=test,method=cycle]
  \dorecurse{8}{\showparagraphshape\samplefile{tufte} \par}
\stopshapedparagraph
```

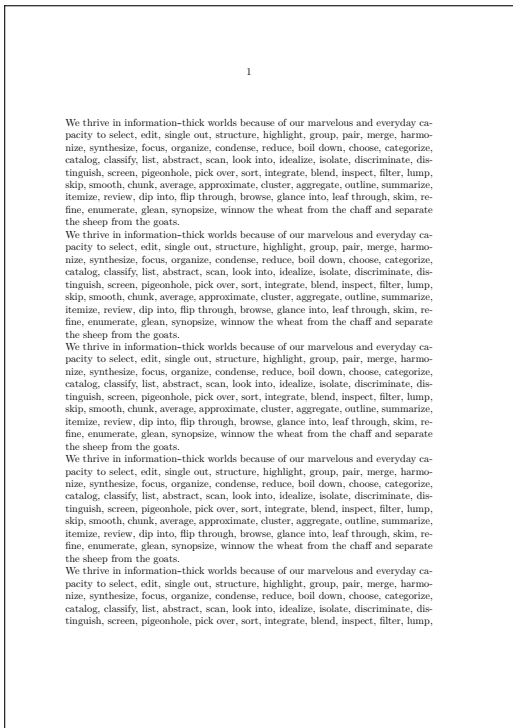
These examples are probably too small to see the details but you can run them yourself or zoom in on the details. In the margin we show the values used. Here is a simple



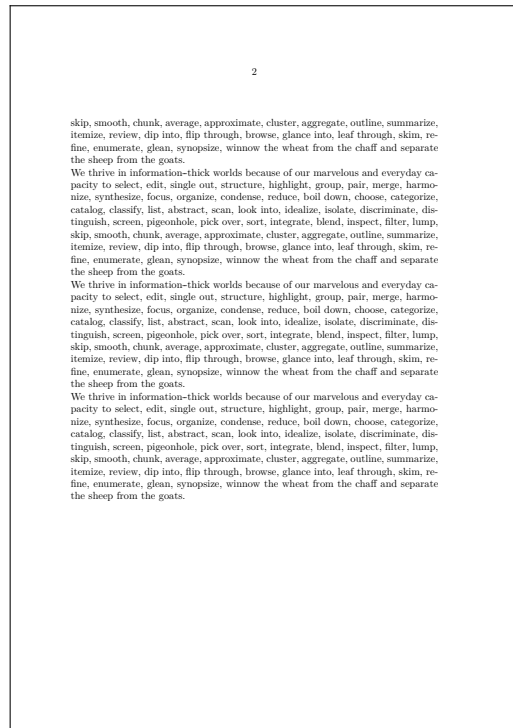
discard, finite shape, page 1



discard, finite shape, page 2



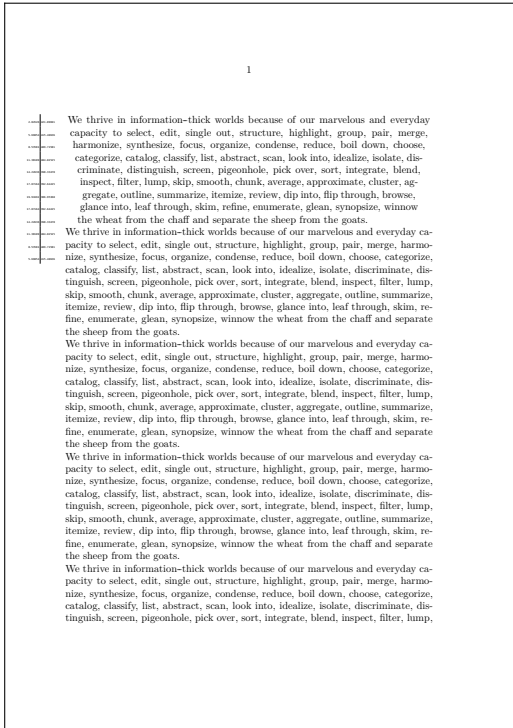
discard, repeat in shape, page 1



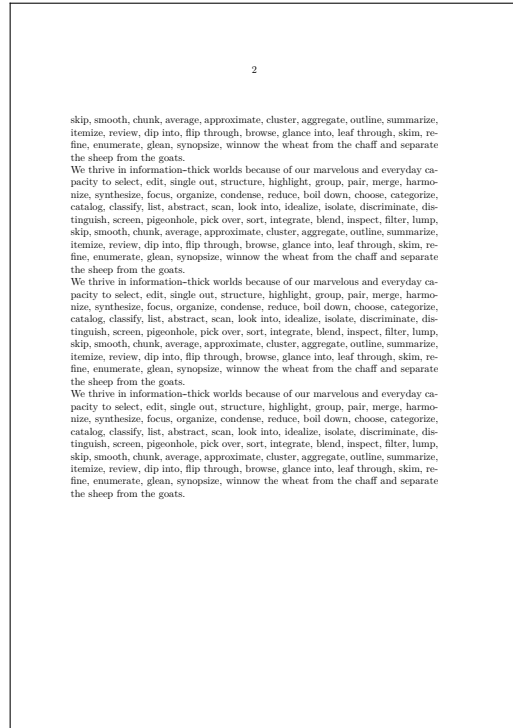
discard, repeat in shape, page 2

Figure 11.2 Discarded shaping

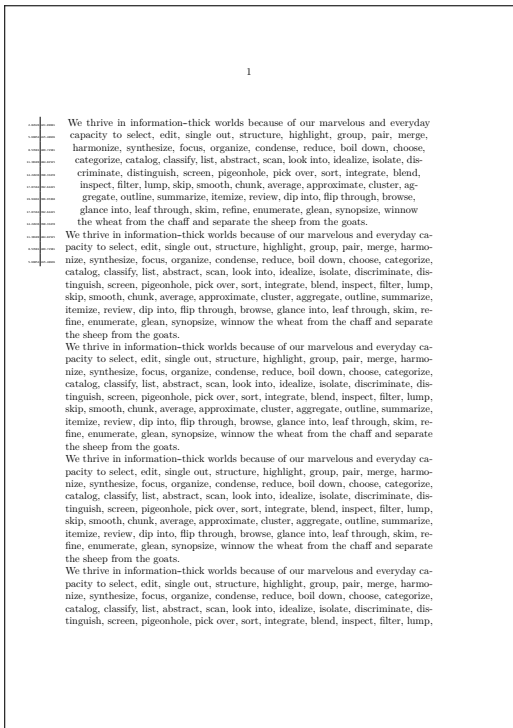




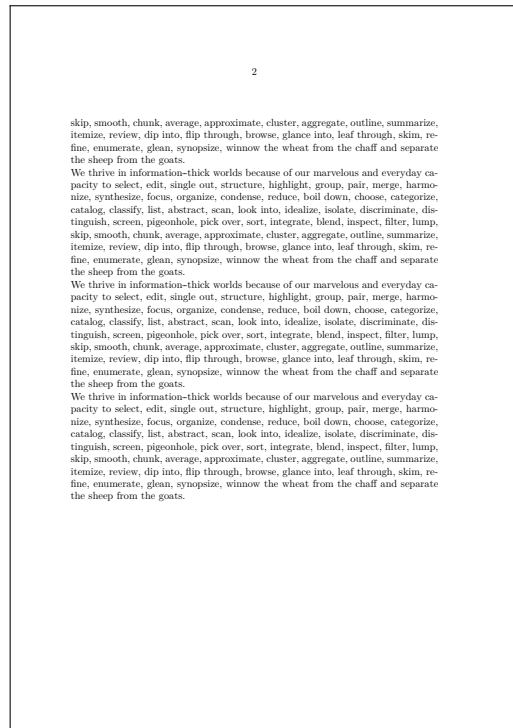
cycle, finite shape, page 1



cycle, finite shape, page 2



cycle, repeat in shape, page 1



cycle, repeat in shape, page 2

Figure 11.4 Cycled shaping

example of (non) poetry. There are other environments that can be used instead but this makes a good example anyway.

```
\startparagraphshape[test]
  left 0em right 0em
  left 1em right 0em
  repeat
\stopparagraphshape

\startshapedparagraph[list=test,method=cycle]
  verse line 1.1\crlf verse line 2.1\crlf
  verse line 3.1\crlf verse line 4.1\par
  verse line 1.2\crlf verse line 2.2\crlf
  verse line 3.2\crlf verse line 4.2\crlf
  verse line 5.2\crlf verse line 6.2\par
\stopshapedparagraph
```

```
verse line 1.1
  verse line 2.1
verse line 3.1
  verse line 4.1
```

```
verse line 1.2
verse line 2.2
verse line 3.2
verse line 4.2
verse line 5.2
verse line 6.2
```

Because the idea for this feature originates in MetaFun, we will now kick in some MetaPost. The following code creates a shape for a circle. We use a 2mm offset here:

```
\startuseMPgraphic{circle}
  path p ; p := fullcircle scaled TextWidth ;
  build_parshape(p,
    2mm, 0, 0,
    LineHeight, StrutHeight, StrutDepth, StrutHeight
  ) ;
\stopuseMPgraphic
```

We plug this into the already described macros:

```
\startshapedparagraph[mp=circle]%
```



```

\setupalign[verytolerant,stretch,last]%
\samplefile{tufte}
\samplefile{tufte}
\stopshapedparagraph

```

And get ourself a circular shape. Watch out, at this moment the shape environment does not add grouping so when as in this case you change the alignment it can influence the document.

We thrive in information-thick  
worlds because of our marvelous and every-  
day capacity to select, edit, single out, structure,  
highlight, group, pair, merge, harmonize, synthesize, focus,  
organize, condense, reduce, boil down, choose, categorize, cat-  
alog, classify, list, abstract, scan, look into, idealize, isolate, discrimi-  
nate, distinguish, screen, pigeonhole, pick over, sort, integrate, blend, in-  
spect, filter, lump, skip, smooth, chunk, average, approximate, cluster, aggre-  
gate, outline, summarize, itemize, review, dip into, flip through, browse, glance  
into, leaf through, skim, refine, enumerate, glean, synopsisize, winnow the wheat  
from the chaff and separate the sheep from the goats. We thrive in information-  
thick worlds because of our marvelous and everyday capacity to select, edit, single  
out, structure, highlight, group, pair, merge, harmonize, synthesize, focus, organize,  
condense, reduce, boil down, choose, categorize, catalog, classify, list, abstract, scan,  
look into, idealize, isolate, discriminate, distinguish, screen, pigeonhole, pick over, sort,  
integrate, blend, inspect, filter, lump, skip, smooth, chunk, average, approximate, clus-  
ter, aggregate, outline, summarize, itemize, review, dip into, flip through, browse,  
glance into, leaf through, skim, refine, enumerate, glean, synopsisize, winnow the  
wheat from the chaff and separate the sheep from the goats.

Assuming that the shape definition above is in a buffer we can do this:

```

\startshapedparagraph[mp=circle]%
\setupalign[verytolerant,stretch,last]%
\samplefile{tufte}
\samplefile{tufte}
\stopshapedparagraph

```

The result is shown in figure 11.5. Because all action happens in the framed environ-  
ment, we can also use this definition:

```

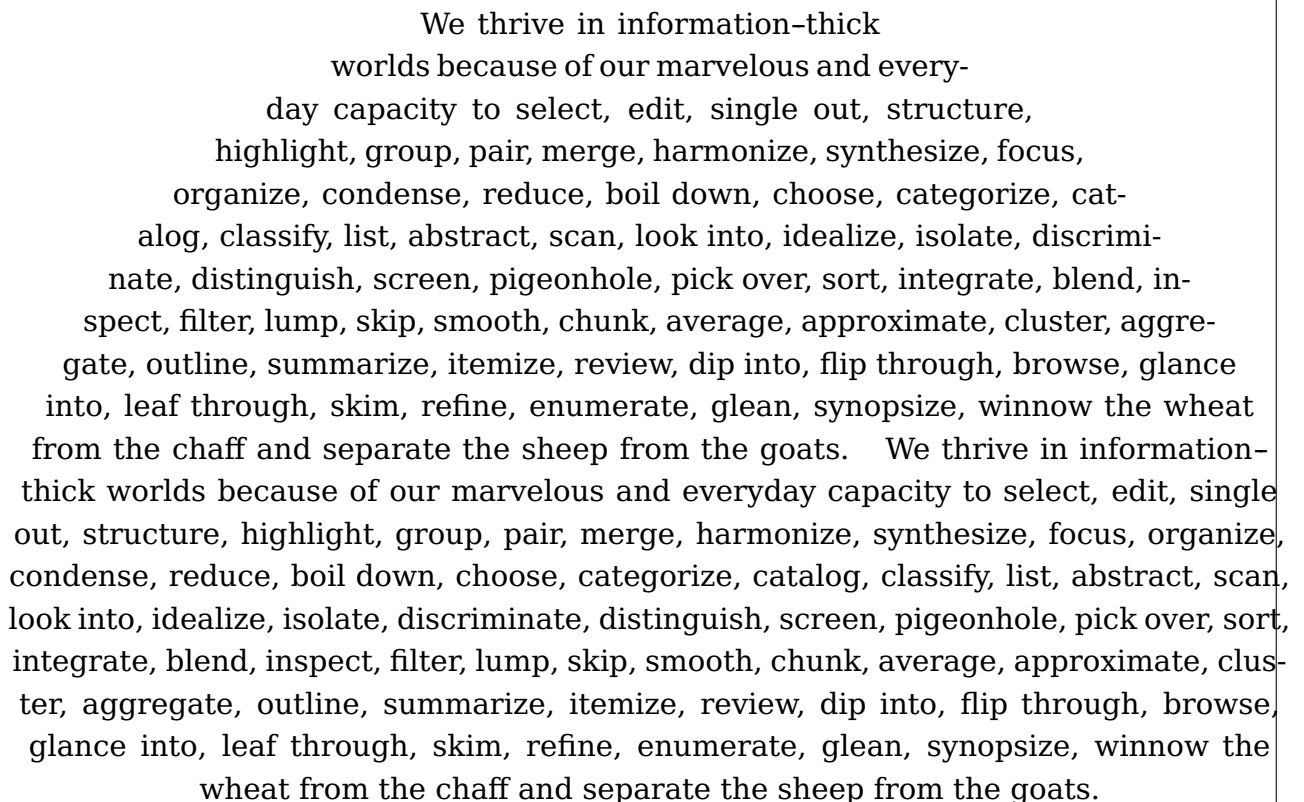
\startuseMPgraphic{circle}

```

```

path p ; p := fullcircle scaled \the\dimexpr\framedwidth+\framedoffset
*2\relax ;
build_parshape(p,
  \framedoffset, 0, 0,
  LineHeight, StrutHeight, StrutDepth, StrutHeight
) ;
draw p ;
\stopuseMPgraphic

```



We thrive in information-thick  
 worlds because of our marvelous and every-  
 day capacity to select, edit, single out, structure,  
 highlight, group, pair, merge, harmonize, synthesize, focus,  
 organize, condense, reduce, boil down, choose, categorize, cat-  
 alog, classify, list, abstract, scan, look into, idealize, isolate, discrimi-  
 nate, distinguish, screen, pigeonhole, pick over, sort, integrate, blend, in-  
 spect, filter, lump, skip, smooth, chunk, average, approximate, cluster, aggre-  
 gate, outline, summarize, itemize, review, dip into, flip through, browse, glance  
 into, leaf through, skim, refine, enumerate, glean, synopsisize, winnow the wheat  
 from the chaff and separate the sheep from the goats. We thrive in information-  
 thick worlds because of our marvelous and everyday capacity to select, edit, single  
 out, structure, highlight, group, pair, merge, harmonize, synthesize, focus, organize,  
 condense, reduce, boil down, choose, categorize, catalog, classify, list, abstract, scan,  
 look into, idealize, isolate, discriminate, distinguish, screen, pigeonhole, pick over, sort,  
 integrate, blend, inspect, filter, lump, skip, smooth, chunk, average, approximate, clus-  
 ter, aggregate, outline, summarize, itemize, review, dip into, flip through, browse,  
 glance into, leaf through, skim, refine, enumerate, glean, synopsisize, winnow the  
 wheat from the chaff and separate the sheep from the goats.

**Figure 11.5** A framed circular shape

A mechanism like this is often never completely automatic in the sense that you need to keep an eye on the results. Depending on user demands more features can be added. With weird shapes you might want to set up the alignment to be tolerant and have some stretch.

The interface described in the MetaFun manual is pretty old, the time stamp of the original code is mid 2000, but the principles didn't change. The examples in `meta-imp-txt.mkx1` can now be written as:

```

\startshapetext[test 1,test 2,test 3,test 4]
  \setupalign[verytolerant,stretch,normal]%

```

```

\samplefile{douglas} % Douglas R. Hofstadter
\stopshapetext
\startcombination[2*2]
  {\framed[offset=overlay,frame=off,background=test 1]{\getshapetext}}
  {test 1}
  {\framed[offset=overlay,frame=off,background=test 2]{\getshapetext}}
  {test 2}
  {\framed[offset=overlay,frame=off,background=test 3]{\getshapetext}}
  {test 3}
  {\framed[offset=overlay,frame=off,background=test 4]{\getshapetext}}
  {test 4}
\stopcombination

```

In figure 11.6 we see the result. Watch how for two shapes we have enabled tracing. Of course you need to tweak till all fits well but we're talking of special situations anyway.

Here is a bit more extreme example. Again we use a circle:

```

\startuseMPgraphic{circle}
  lmt_parshape [
    path      = fullcircle scaled 136mm,
    offset    = 2mm,
    bottomskip = - 1.5LineHeight,
  ] ;
\stopuseMPgraphic

```

But we output a longer text:

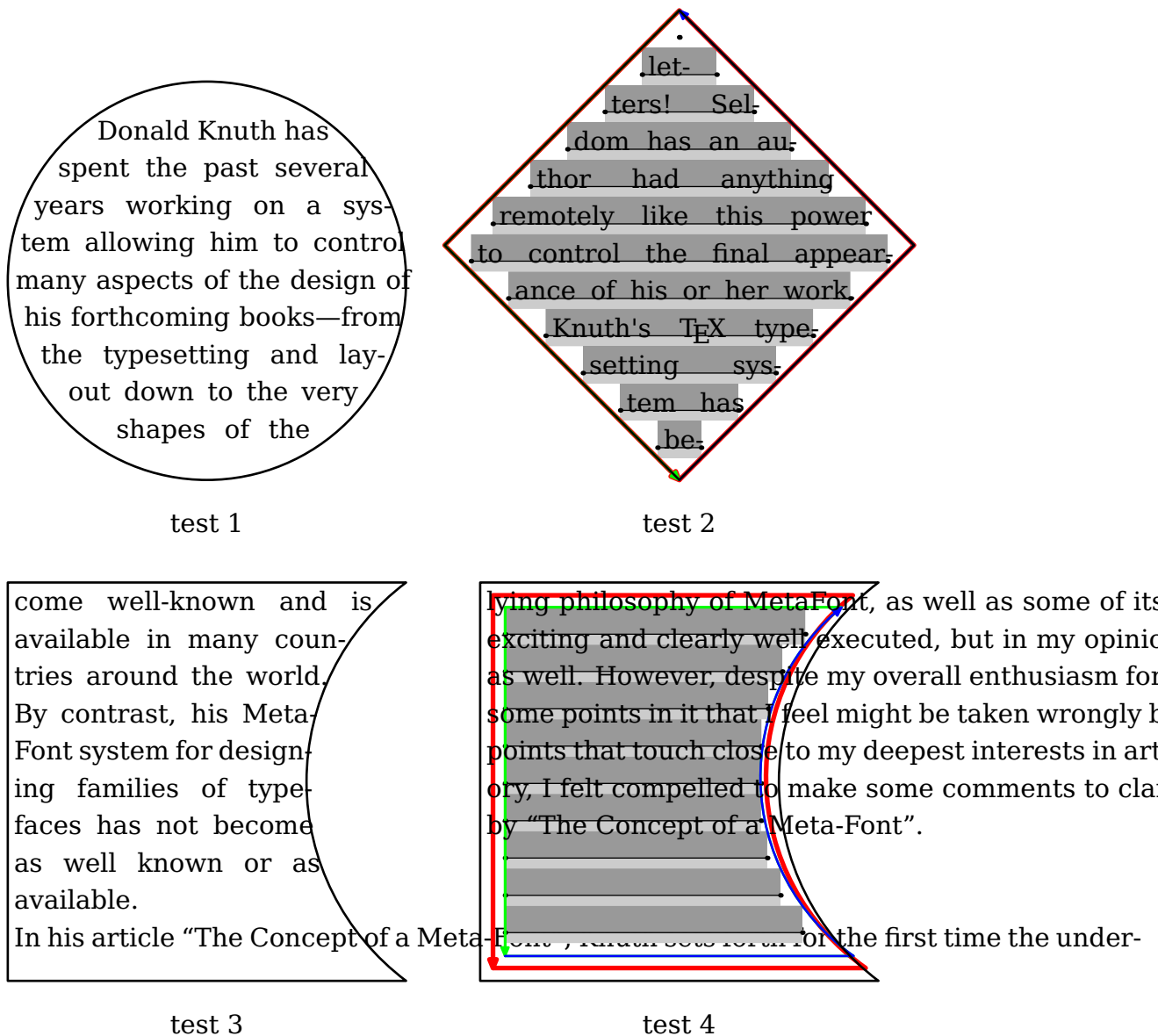
```

\startshapedparagraph[mp=circle,repeat=yes,method=cycle]%
  \setupalign[verytolerant,stretch,last]\dontcomplain
  {\darkred      \samplefile{tufte}}\par
  {\darkgreen    \samplefile{tufte}}\par
  {\darkblue     \samplefile{tufte}}\par
  {\darkcyan     \samplefile{tufte}}\par
  {\darkmagenta  \samplefile{tufte}}\par
\stopshapedparagraph

```

We get a multi-page shape:

We thrive in information-thick  
worlds because of our marvelous and every-  
day capacity to select, edit, single out, structure,



**Figure 11.6**

highlight, group, pair, merge, harmonize, synthesize, focus, organize, condense, reduce, boil down, choose, categorize, catalog, classify, list, abstract, scan, look into, idealize, isolate, discriminate, distinguish, screen, pigeonhole, pick over, sort, integrate, blend, inspect, filter, lump, skip, smooth, chunk, average, approximate, cluster, aggregate, outline, summarize, itemize, review, dip into, flip through, browse, glance into, leaf through, skim, refine, enumerate, glean, synopsise, winnow the wheat from the chaff and separate the sheep from the goats.

We thrive in information-thick worlds because of our marvelous and everyday capacity to select, edit, single out, structure, highlight, group, pair, merge, harmonize, synthe-

size, focus, organize, condense, reduce, boil down, choose, categorize, catalog, classify, list, abstract, scan, look into, idealize, isolate, discriminate, distinguish, screen, pigeon-hole, pick over, sort, integrate, blend, inspect, filter, lump, skip, smooth, chunk, average, approximate, cluster, aggregate, outline, summarize, itemize, review, dip into, flip through, browse, glance into, leaf through, skim, refine, enumerate, glean, synopsisize, winnow the wheat from the chaff and separate the sheep from the goats.

We thrive in information-thick worlds because of our marvelous and everyday capacity to select, edit, single out, structure, highlight, group, pair, merge, harmonize, synthesize, focus, organize, condense, reduce, boil down, choose, categorize, catalog, classify, list, abstract, scan, look into, idealize, isolate, discriminate, distinguish, screen, pigeon-hole, pick over, sort, integrate, blend, inspect, filter, lump, skip, smooth, chunk, average, approximate, cluster, aggregate, outline, summarize, itemize, review, dip into, flip through, browse, glance into, leaf through, skim, refine, enumerate, glean, synopsisize, winnow the wheat from the chaff and separate the sheep from the goats.

We thrive in information-thick worlds because of our marvelous and everyday capacity to select, edit, single out, structure, highlight, group, pair, merge, harmonize, synthesize, focus, organize, condense, reduce, boil down, choose, categorize, catalog, classify, list, abstract, scan, look into, idealize, isolate, discriminate, distinguish, screen, pigeon-hole, pick over, sort, integrate, blend, inspect, filter, lump, skip, smooth, chunk, average, approximate, cluster, aggregate, outline, summarize, itemize, review, dip into, flip through, browse, glance into, leaf through, skim, refine, enumerate, glean, synopsisize, winnow the wheat from the chaff and separate the sheep from the goats.

We thrive in information-thick worlds because of our marvelous and everyday capacity to select, edit, single out, structure, highlight, group, pair, merge, harmonize, synthesize, focus, organize, condense, reduce, boil down, choose, categorize, catalog, classify, list, abstract, scan, look into, idealize, isolate, discriminate, distinguish, screen, pigeon-hole, pick over, sort, integrate, blend, inspect, filter, lump, skip, smooth, chunk, average, approximate, cluster, aggregate, outline, summarize, itemize, review, dip into, flip through, browse, glance into, leaf through, skim, refine, enumerate, glean, synopsisize, winnow the wheat from the chaff and separate the sheep from the goats.

Compare this with:

```
\startshapedparagraph[mp=circle,repeat=yes,method=cycle]%
  \setupalign[verytolerant,stretch,last]\dontcomplain
  {\darkred      \samplefile{tufte}}
  {\darkgreen    \samplefile{tufte}}
  {\darkblue     \samplefile{tufte}}
  {\darkcyan     \samplefile{tufte}}
```

`{\darkmagenta \samplefile{tufte}}  
\stopshapedparagraph`

Which gives:

We thrive in information-thick worlds because of our marvelous and everyday capacity to select, edit, single out, structure, highlight, group, pair, merge, harmonize, synthesize, focus, organize, condense, reduce, boil down, choose, categorize, catalog, classify, list, abstract, scan, look into, idealize, isolate, discriminate, distinguish, screen, pigeonhole, pick over, sort, integrate, blend, inspect, filter, lump, skip, smooth, chunk, average, approximate, cluster, aggregate, outline, summarize, itemize, review, dip into, flip through, browse, glance into, leaf through, skim, refine, enumerate, glean, synthesize, winnow the wheat from the chaff and separate the sheep from the goats. We thrive in information-thick worlds because of our marvelous and everyday capacity to select, edit, single out, structure, highlight, group, pair, merge, harmonize, synthesize, focus, organize, condense, reduce, boil down, choose, categorize, catalog, classify, list, abstract, scan, look into, idealize, isolate, discriminate, distinguish, screen, pigeonhole, pick over, sort, integrate, blend, inspect, filter, lump, skip, smooth, chunk, average, approximate, cluster, aggregate, outline, summarize, itemize, review, dip into, flip through, browse, glance into, leaf through, skim, refine, enumerate, glean, synthesize, winnow the wheat from the chaff and separate the sheep from the goats. We thrive in information-thick worlds because of our marvelous and everyday capacity to select, edit, single out, structure, highlight, group, pair, merge, harmonize, synthesize, focus, organize, condense, reduce, boil down, choose, categorize, catalog, classify, list, abstract, scan, look into, idealize, isolate, discriminate, distinguish, screen, pigeonhole, pick over, sort, integrate, blend, inspect, filter, lump, skip, smooth, chunk, average, approximate, cluster, aggregate, outline, summarize, itemize, review, dip into, flip through, browse, glance into, leaf through, skim, refine, enumerate, glean, synthesize, winnow the wheat from the chaff and separate the sheep from the goats. We thrive in information-thick worlds because of our marvelous and everyday capacity to select, edit, single out, structure, highlight, group, pair, merge, harmonize, synthesize, focus, organize, condense, reduce, boil down, choose, categorize, catalog, classify, list, abstract, scan, look into, idealize, isolate, discriminate, distinguish, screen, pigeonhole, pick over, sort, integrate, blend, inspect, filter, lump, skip, smooth, chunk, average, approximate, cluster, aggregate, outline, summarize, itemize, review, dip into, flip through, browse, glance into, leaf through, skim, refine, enumerate, glean, synthesize, winnow the wheat from the chaff and separate the sheep from the goats. We thrive in information-thick worlds because of our marvelous and everyday capacity to select, edit, single out, structure, highlight, group, pair, merge, harmonize, synthesize, focus, organize, condense, reduce, boil down, choose, categorize, catalog, classify, list, abstract, scan, look into, idealize, isolate, discriminate, distinguish, screen, pigeonhole, pick over, sort, integrate, blend, inspect, filter, lump, skip, smooth, chunk, average, approximate, cluster, aggregate, outline, summarize, itemize, review, dip into, flip through, browse, glance into, leaf through, skim, refine, enumerate, glean, synthesize, winnow the wheat from the chaff and separate the sheep from the goats. We thrive in information-thick worlds because of our marvelous and everyday capacity to select, edit, single out, structure, highlight, group, pair, merge, harmonize, synthesize, focus, organize, condense, reduce, boil down, choose, categorize, catalog, classify, list, abstract, scan, look into, idealize, isolate, discriminate, distinguish, screen, pigeonhole, pick over, sort, integrate, blend, inspect, filter, lump, skip, smooth, chunk, average, approximate, cluster, aggregate, outline, summarize, itemize, review, dip into, flip through, browse, glance into, leaf through, skim, refine, enumerate, glean, synthesize, winnow the wheat from the chaff and separate the sheep from the goats.

classify, list, abstract, scan, look into, idealize, isolate, discriminate, distinguish, screen, pigeonhole, pick over, sort, integrate, blend, inspect, filter, lump, skip, smooth, chunk, average, approximate, cluster, aggregate, outline, summarize, itemize, review, dip into, flip through, browse, glance into, leaf through, skim, refine, enumerate, glean, synopsisize, winnow the wheat from the chaff and separate the sheep from the goats. We thrive in information-thick worlds because of our marvelous and everyday capacity to select, edit, single out, structure, highlight, group, pair, merge, harmonize, synthesize, focus, organize, condense, reduce, boil down, choose, categorize, catalog, classify, list, abstract, scan, look into, idealize, isolate, discriminate, distinguish, screen, pigeonhole, pick over, sort, integrate, blend, inspect, filter, lump, skip, smooth, chunk, average, approximate, cluster, aggregate, outline, summarize, itemize, review, dip into, flip through, browse, glance into, leaf through, skim, refine, enumerate, glean, synopsisize, winnow the wheat from the chaff and separate the sheep from the goats.

Here the `bottomskip` takes care of subtle rounding issues as well as discarding the last line in the shape so that we get nicer continuation. There is no full automated solution for all you can come up with.

Mixing a MetaPost specification into a regular one is also possible. The next example demonstrates this as well as the option to remove some lines from a specification:

```
\startparagraphshape[test]
  left 0em right 0em
  left 1em right 0em
  metapost {circle}
  delete 3
  metapost {circle,circle,circle}
  delete 7
  metapost {circle}
  repeat
\stopparagraphshape
```

You can combine a shape with narrowing a paragraph. Watch the `absolute` keyword in the next code. The result is shown in figure 11.7.

```
\startuseMPgraphic{circle}
  lmt_parshape [
    path      = fullcircle scaled TextWidth,
    bottomskip = - 1.5LineHeight,
  ] ;
\stopuseMPgraphic
```

```

\startparagraphshape[test-1]
  metapost {circle} repeat
\stopparagraphshape

\startparagraphshape[test-2]
  absolute left metapost {circle} repeat
\stopparagraphshape

\startparagraphshape[test-3]
  absolute right metapost {circle} repeat
\stopparagraphshape

\startparagraphshape[test-4]
  absolute both metapost {circle} repeat
\stopparagraphshape

\showframe

\startnarrower[4*left,2*right]
  \startshapedparagraph[list=test-1,repeat=yes,method=repeat]%
    \setupalign[verytolerant,stretch,last]\dontcomplain
    \dorecurse{3}{\samplefile{thuan}}
  \stopshapedparagraph
  \page
  \startshapedparagraph[list=test-2,repeat=yes,method=repeat]%
    \setupalign[verytolerant,stretch,last]\dontcomplain
    \dorecurse{3}{\samplefile{thuan}}
  \stopshapedparagraph
  \page
  \startshapedparagraph[list=test-3,repeat=yes,method=repeat]%
    \setupalign[verytolerant,stretch,last]\dontcomplain
    \dorecurse{3}{\samplefile{thuan}}
  \stopshapedparagraph
  \page
  \startshapedparagraph[list=test-4,repeat=yes,method=repeat]%
    \setupalign[verytolerant,stretch,last]\dontcomplain
    \dorecurse{3}{\samplefile{thuan}}
  \stopshapedparagraph
\stopnarrower

```





The shape mechanism has a few more tricks but these are really meant for usage in specific situations, where one knows what one deals with. The following examples are visualized in figure 11.8.

```

\useMPlibrary[dum]
\usemodule[article-basics]

\startbuffer
  \externalfigure[dummy][width=6cm]
\stopbuffer

\startshapedparagraph[text=\getbuffer]
  \dorecurse{3}{\samplefile{ward}\par}
\stopshapedparagraph

\page

\startshapedparagraph[text=\getbuffer,distance=1em]
  \dorecurse{3}{\samplefile{ward}\par}
\stopshapedparagraph

\page

\startshapedparagraph[text=\getbuffer,distance=1em,
  hoffset=-2em]
  \dorecurse{3}{\samplefile{ward}\par}
\stopshapedparagraph

\page

\startshapedparagraph[text=\getbuffer,distance=1em,
  voffset=-2ex,hoffset=-2em]
  \dorecurse{3}{\samplefile{ward}\par}
\stopshapedparagraph

\page

\startshapedparagraph[text=\getbuffer,distance=1em,
  voffset=-2ex,hoffset=-2em,lines=1]
  \dorecurse{3}{\samplefile{ward}\par}
\stopshapedparagraph

```

\page

\startshapedparagraph[width=4cm,lines=4]

\dorecurse{3}{\samplefile{ward}\par}

\stopshapedparagraph



Figure 11.8 Flow around something

## 11.7 Modes

*todo: some of the side effects of so called modes*

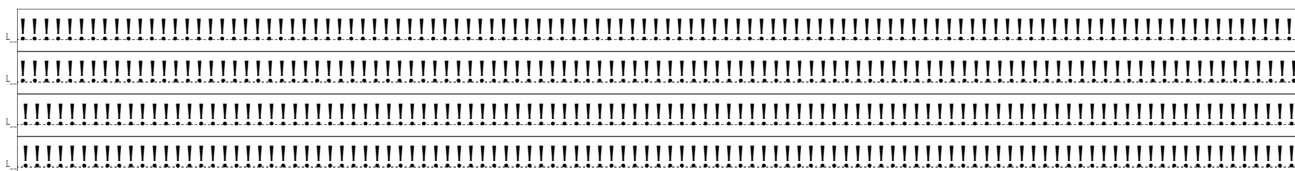
### Modes

## 11.8 Leaders

Leaders are a basic feature that users probably never run into directly. They repeat content till it fits the specified width which can be stretched out. The content is typeset once and it is the backend that does the real work of repetition.

```
\strut\leaders \hbox{!}\hfill\strut
\strut\xleaders\hbox{!}\hfill\strut
\strut\cleaders\hbox{!}\hfill\strut
\strut\gleaders\hbox{!}\hfill\strut
```

Here `\leaders` starts at the left edge and repeats the box as long as it fits, `\xleaders` spreads till the edges and `\cleaders` centers the lot. The `\gleaders` primitive (which is not in original  $\text{T}_{\text{E}}\text{X}$ ) takes the outer box as reference and further behaves like `\cleaders`.

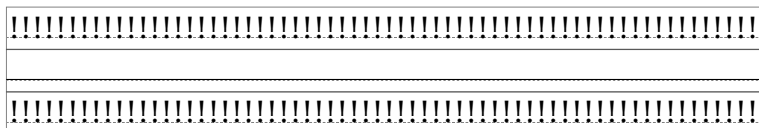


The leader primitives take box or rule but in  $\text{LuaMetaT}_{\text{E}}\text{X}$  a glyph can also be specified, which saves wrapping in a box.

```
\ruledvbox \bgroup \hsize 10cm
  \strut\cleaders\hbox{!}\hfill\strut
\egroup
```

```
\ruledvbox \bgroup \hsize 10cm
  \strut\cleaders\hrule\hfill\strut
\egroup
```

```
\ruledvbox \bgroup \hsize 10cm
  \strut\cleaders\glyph`!\hfill\strut
\egroup
```



The  $\text{LuaMetaT}_{\text{E}}\text{X}$  engine also introduced `\uleaders`

We show three boxes, a regular one first (red):

```
x xx xxx xxxx
\ruledhbox{L\hss R}\space
x xx xxx xxxx
```

The second one (blue) is also a box but one that stretches upto 100pt and is in a later stage, when the paragraph has been built, is repackaged to the effective width. The third example (green) leaves out the background.

```
x xx xxx xxxx L R x xx xxx xxxx x xx xxx xxxx L R x xx xxx xxxx x xx xxx xxxx L R x xx xxx xxxx x xx xxx xxxx
L R x xx xxx xxxx x xx xxx xxxx x xx xxx xxxx L R x xx xxx xxxx x xx xxx xxxx L R x xx xxx xxxx x xx xxx xxxx
xx xxx xxxx L R x xx xxx xxxx x xx xxx xxxx x xx xxx xxxx L R x xx xxx xxxx x xx xxx xxxx L R x xx xxx xxxx x xx xxx xxxx
x xx xxx xxxx L R x xx xxx xxxx x xx xxx xxxx L R x xx xxx xxxx x xx xxx xxxx L R x xx xxx xxxx x xx xxx xxxx
xx xxx xxxx L R x xx xxx xxxx x xx xxx xxxx L R x xx xxx xxxx x xx xxx xxxx L R x xx xxx xxxx x xx xxx xxxx
x xx xxx xxxx L R x xx xxx xxxx x xx xxx xxxx L R x xx xxx xxxx x xx xxx xxxx L R x xx xxx xxxx x xx xxx xxxx
```

In ConT<sub>E</sub>Xt we have wrapped this feature in the adaptive box mechanism, so here a few a few examples:

```
\startsetups adaptive:test:a
  \setbox\usedadaptivebox\vbox to \usedadaptivetotal \bgroup
  \externalfigure
  [cow.pdf]
  [width=\framedmaxwidth,
   frame=on,
   height=\usedadaptivetotal]%
  \egroup
\stopsetups
```

```
\startsetups adaptive:test:b
  \setbox\usedadaptivebox\vbox to \usedadaptivetotal \bgroup
  \externalfigure
  [cow.pdf]
  [width=\usedadaptivewidth,
   frame=on,
   height=\usedadaptivetotal]%
  \egroup
\stopsetups
```

We use this as follows (see figure 11.9 for the result):

```
\framed[height=18cm,align=middle,adaptive=yes,top=,bottom=] {%
  \begstrut \samplefile{tufte} \endstrut
```

```

\par
\adaptivevbox
  [strut=yes, setups=adaptive:test:a]
  {\showstruts\strut\hsize5cm\hss}%
\par
\adaptivevbox
  [strut=yes, setups=adaptive:test:b]
  {\showstruts\strut\hsize5cm\hss}%
\par
\begstrut \samplefile{tufte} \endstrut
}

```

Here is one that you can test yourself:

```

\startsetups adaptive:test
  \setbox\usedadaptivebox\vbox to \usedadaptivetotal \bgroup
  \externalfigure
    [cow.pdf]
    [width=\usedadaptivewidth,
     height=\usedadaptivetotal]%
  \egroup
\stopsetups

\ruledvbox to \textheight {
  \par \begstrut \samplefile{tufte} \endstrut \par
  \adaptivevbox[strut=yes, setups=adaptive:test]{\hsize\textwidth\hss}
  \par \begstrut \samplefile{tufte} \endstrut
}

```

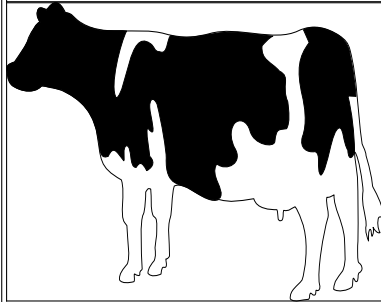
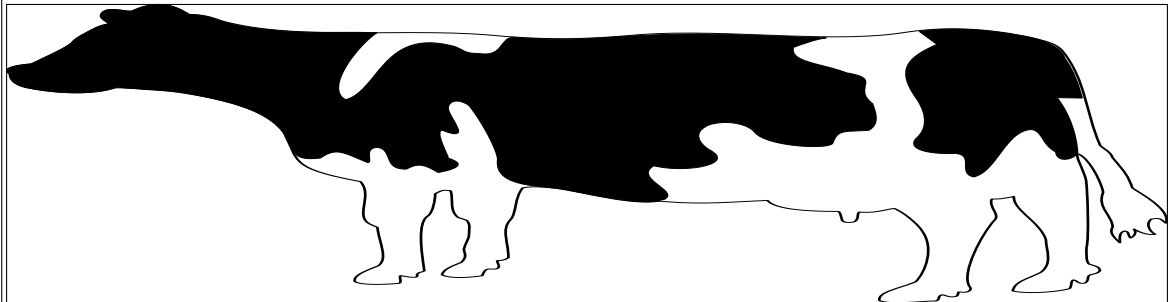
The next example comes from the test suite (where it runs over many pages in order to illustrate the idea):

```

\startMPdefinitions
  def TickTock =
    interim linecap := squared;
    save p ; path p ;
    p := fullsquare xysized(AdaptiveWidth, .9(AdaptiveHeight+AdaptiveDepth))
  ;
  fill p withcolor AdaptiveColor ;
  draw bottomboundary (p enlarged (-AdaptiveThickness) )
    withdashes (3*AdaptiveThickness)

```

We thrive in information-thick worlds because of our marvelous and everyday capacity to select, edit, single out, structure, highlight, group, pair, merge, harmonize, synthesize, focus, organize, condense, reduce, boil down, choose, categorize, catalog, classify, list, abstract, scan, look into, idealize, isolate, discriminate, distinguish, screen, pigeonhole, pick over, sort, integrate, blend, inspect, filter, lump, skip, smooth, chunk, average, approximate, cluster, aggregate, outline, summarize, itemize, review, dip into, flip through, browse, glance into, leaf through, skim, refine, enumerate, glean, synthesize, winnow the wheat from the chaff and separate the sheep from the goats.



We thrive in information-thick worlds because of our marvelous and everyday capacity to select, edit, single out, structure, highlight, group, pair, merge, harmonize, synthesize, focus, organize, condense, reduce, boil down, choose, categorize, catalog, classify, list, abstract, scan, look into, idealize, isolate, discriminate, distinguish, screen, pigeonhole, pick over, sort, integrate, blend, inspect, filter, lump, skip, smooth, chunk, average, approximate, cluster, aggregate, outline, summarize, itemize, review, dip into, flip through, browse, glance into, leaf through, skim, refine, enumerate, glean, synthesize, winnow the wheat from the chaff and separate the sheep from the goats.

**Figure 11.9**

```

withpen pencircle scaled AdaptiveThickness
withcolor white ;
enddef ;
\stopMPdefinitions

\startsetups adaptive:test

```

```

\setbox\usedadaptivebox\hbox
  to      \usedadaptivewidth
  yoffset -.9\usedadaptivedepth
\bggroup
  \hss
  \startMPcode
    TickTock ;
  \stopMPcode
  \hss
\egroup
\stopsetups

```

```

\definecolor[adaptive:tick][.25(blue,green)]
\definecolor[adaptive:tock][.75(blue,green)]

```

```

\defineadaptive
[tick]
[setups=adaptive:test,
 color=adaptive:tick,
 foregroundcolor=white,
 foregroundstyle=\infofont,
 strut=yes]

```

```

\defineadaptive
[tock]
[tick]
[color=adaptive:tock]

```

















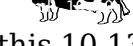
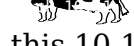
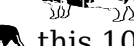



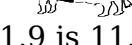
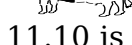
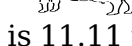
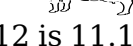
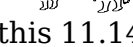
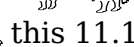
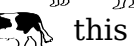
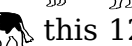
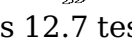
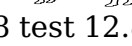
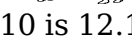
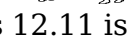
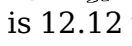
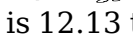
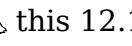

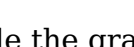
```

\dostepwiserecurse{8}{12}{1}{%
  \dostepwiserecurse{5}{15}{1}{%
    this~#1.##1 is~#1.##1 test~#1.##1
    \ifodd##1\relax
      \adaptivebox[tick]{\hss tick #1.##1\hss}
    \else
      \adaptivebox[tock]{\hss tock #1.##1\hss}
    \fi
  }
}

```

this 8.5 is 8.5 test 8.5 tick 8.5 this 8.6 is 8.6 test 8.6 tock 8.6 this 8.7 is 8.7 test 8.7  
 this 8.8 is 8.8 test 8.8 tock 8.8 this 8.9 is 8.9 test 8.9 tick 8.9 this 8.10 is 8.10 test 8.10



this 8.11 is 8.11 test 8.11  this 8.12 is 8.12 test 8.12  this 8.13 is 8.13  
 test 8.13  this 8.14 is 8.14 test 8.14  this 8.15 is 8.15 test 8.15  
 this 9.5 is 9.5 test 9.5  this 9.6 is 9.6 test 9.6  this 9.7 is 9.7 test 9.7  
 this 9.8 is 9.8 test 9.8  this 9.9 is 9.9 test 9.9  this 9.10 is 9.10 test 9.10  
 this 9.11 is 9.11 test 9.11  this 9.12 is 9.12 test 9.12  this 9.13 is 9.13  
 test 9.13  this 9.14 is 9.14 test 9.14  this 9.15 is 9.15 test 9.15  
 this 10.5 is 10.5 test 10.5  this 10.6 is 10.6 test 10.6  this 10.7 is 10.7  
 test 10.7  this 10.8 is 10.8 test 10.8  this 10.9 is 10.9 test 10.9  
 this 10.10 is 10.10 test 10.10  this 10.11 is 10.11 test 10.11  this 10.12  
 is 10.12 test 10.12  this 10.13 is 10.13 test 10.13  this 10.14 is 10.14  
 test 10.14  this 10.15 is 10.15 test 10.15  this 11.5 is 11.5 test 11.5  
 this 11.6 is 11.6 test 11.6  this 11.7 is 11.7 test 11.7  this 11.8 is 11.8  
 test 11.8  this 11.9 is 11.9 test 11.9  this 11.10 is 11.10 test 11.10  
 this 11.11 is 11.11 test 11.11  this 11.12 is 11.12 test 11.12  this 11.13  
 is 11.13 test 11.13  this 11.14 is 11.14 test 11.14  this 11.15 is 11.15  
 test 11.15  this 12.5 is 12.5 test 12.5  this 12.6 is 12.6 test 12.6  
 this 12.7 is 12.7 test 12.7  this 12.8 is 12.8 test 12.8  this 12.9 is 12.9  
 test 12.9  this 12.10 is 12.10 test 12.10  this 12.11 is 12.11 test 12.11  
 this 12.12 is 12.12 test 12.12  this 12.13 is 12.13 test 12.13  this 12.14  
 is 12.14 test 12.14  this 12.15 is 12.15 test 12.15









In the next example the graphics adapt to the available space:

```

\startsetups adaptive:test
  \setbox\usedadaptivebox\hbox
    to      \usedadaptivewidth
    yoffset -\usedadaptivedepth
  \bgroup
    \externalfigure
      [cow.pdf]
      [width=\usedadaptivewidth,
        height=\dimexpr\usedadaptivetotal\relax]%
  \egroup
\stopsetups

\dostepwiserecurse{1}{50}{1}{%
  this~#1 is~#1 test~#1
  {\adaptivebox[strut=yes,setup=adaptive:test]{} }
}

```

this 1 is 1 test 1  this 2 is 2 test 2  this 3 is 3 test 3  this 4 is 4 test 4  this 5  
 is 5 test 5  this 6 is 6 test 6  this 7 is 7 test 7  this 8 is 8 test 8  this 9 is 9

test 9 🐮 this 10 is 10 test 10 🐮 this 11 is 11 test 11 🐮 this 12 is 12 test 12 🐮 this 13  
 is 13 test 13 🐮 this 14 is 14 test 14 🐮 this 15 is 15 test 15 🐮 this 16 is 16 test 16  
 this 17 is 17 test 17 🐮 this 18 is 18 test 18 🐮 this 19 is 19 test 19 🐮 this 20 is 20  
 test 20 🐮 this 21 is 21 test 21 🐮 this 22 is 22 test 22 🐮 this 23 is 23 test 23  
 this 24 is 24 test 24 🐮 this 25 is 25 test 25 🐮 this 26 is 26 test 26 🐮 this 27 is 27  
 test 27 🐮 this 28 is 28 test 28 🐮 this 29 is 29 test 29 🐮 this 30 is 30 test 30  
 this 31 is 31 test 31 🐮 this 32 is 32 test 32 🐮 this 33 is 33 test 33 🐮 this 34 is 34  
 test 34 🐮 this 35 is 35 test 35 🐮 this 36 is 36 test 36 🐮 this 37 is 37 test 37  
 this 38 is 38 test 38 🐮 this 39 is 39 test 39 🐮 this 40 is 40 test 40 🐮 this 41 is 41  
 test 41 🐮 this 42 is 42 test 42 🐮 this 43 is 43 test 43 🐮 this 44 is 44 test 44  
 this 45 is 45 test 45 🐮 this 46 is 46 test 46 🐮 this 47 is 47 test 47 🐮 this 48 is 48  
 test 48 🐮 this 49 is 49 test 49 🐮 this 50 is 50 test 50

## 11.9 Prevdepth

The depth of a box is normally positive but rules can have a negative depth in order to get a rule above the baseline. When T<sub>E</sub>X was written the assumption was that a negative depth of more than 1000 point made no sense at all. The last depth on a vertical list is registered in the `\prevdepth` variable. This is basically a reference into the current list. In order to illustrate some interesting side effects of setting this `\prevdepth` and especially when we set it to `-1000pt`. In order to illustrate this this special value can be set to a different value in LuaMetaT<sub>E</sub>X. However, as dealing with the property is somewhat special in the engine you should not set it unless you know that the macro package is ware of it.

```
line 1\par line 2 \par \nointerlineskip line 3 \par
```

Assuming that we haven't set any inter paragraph spacing this gives:

line 1
line 2
line 3

Here `\nointerlineskip` is (normally) defined as:

```
\prevdepth-1000pt
```

although in ConT<sub>E</sub>Xt we use `\ignoredepthcriterion` instead of the hard coded dimension. We now give a more extensive example:

```
\ruledhbox \bgroup
  \PrevTest{-10.0pt}\quad
```

```

\PrevTest{-20.0pt}\quad
\PrevTest{-49.9pt}\quad
\PrevTest{-50.0pt}\quad
\PrevTest{-50.1pt}\quad
\PrevTest{-60.0pt}\quad
\PrevTest{-80.0pt}%

```

**\egroup**

In this example we set `\ignoredepthcriterion` to `-50.0pt` instead of the normal `-1000pt`. The helper is defined as:

```

\def\PrevTest#1%
{\setbox0\ruledhbox{\strut$\tf#1$}%
 \dp0=#1
 \vbox\bgroup\hsize4em
  FIRST\par
  \unhbox0\par
  LAST\par
 \egroup}

```

or

```

\def\PrevTest#1%
{\setbox0\ruledhbox{\strut$\tf#1$}%
 \dp0=#1
 \vbox\bgroup
  \ruledhbox{FIRST}\par
  \box0\par
  \ruledhbox{LAST}\par
 \egroup}

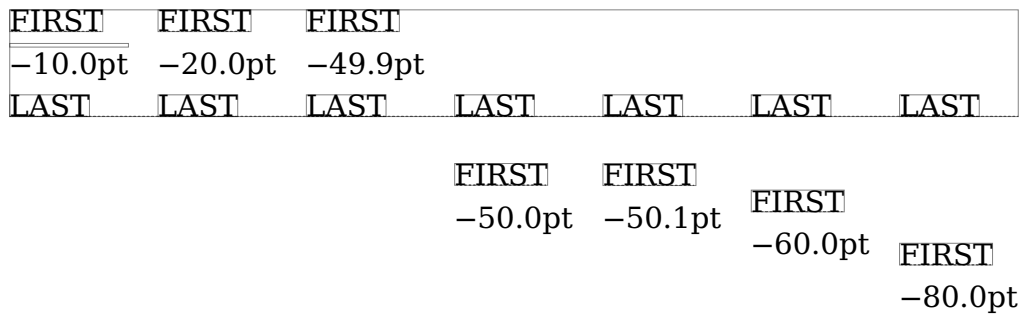
```

The result is shown in figures 11.10 and 11.11. The first case is what we normally have in text and we haven't set `\prevdepth` explicitly between lines so  $\text{\TeX}$  will just look at the depth of the lines. In the second case the depth is ignored when less than the criterion which is why, when we set the depth of the box to a negative value we get somewhat interesting skips.

FIRST	FIRST	FIRST	FIRST	FIRST	FIRST	FIRST
-10.0pt	-20.0pt	-49.9pt	-50.0pt	-50.1pt	-60.0pt	-80.0pt
LAST	LAST	LAST	LAST	LAST	LAST	LAST

**Figure 11.10**

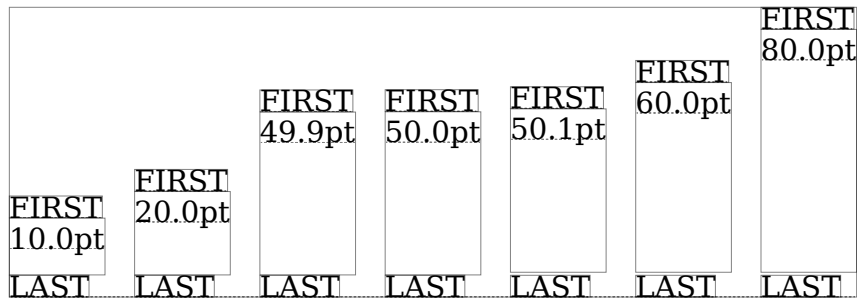
**Prevdepth**



**Figure 11.11**

I'm sure one can use this effect otherwise than intended but I doubt is any user is willing to do this but the fact that we can lower the criterion makes for nice experiments. Just for the record, in figure 11.12 you see what we get with positive values:

```
\ruledhbox \bgroup
  \PrevTest{10.0pt}\quad
  \PrevTest{20.0pt}\quad
  \PrevTest{49.9pt}\quad
  \PrevTest{50.0pt}\quad
  \PrevTest{50.1pt}\quad
  \PrevTest{60.0pt}\quad
  \PrevTest{80.0pt}%
\egroup
```



**Figure 11.12**

Watch the interline skip kicking in when we make the depth larger than in `\ignore-depthcriterion` being 50pt.

## 11.10 Normalization

*todo: users don't need to bother about this but it might be interesting anyway*

## 11.11 Dirty tricks

*todo: explain example for combining paragraphs*

## 11.12 Penalties

In figure 11.13 we demonstrate the (accumulated) effect of a few penalty arrays that you can set. Keep in mind that these are low level (primitive) commands that can interfere with other mechanisms. The interline penalties are applied to the current paragraph in the same way as `\looseness` is. This makes sense because using the club and widow penalties is more predictable. You have to zoom in to see how the penalties add up. As with `\looseness` the `\interlinepenalties` get reset before the paragraph, which is shown in the bottom cells of this figure.

```
\interlinepenalty 0
\clubpenalty      0
\widowpenalty     0
\orphanpenalty    0
\shapingpenalty   0

\clubpenalties    5  1000  2000  3000  4000  5000 % 6 -> 0
\widowpenalties   5   10    20    30    40    50 % 6 -> 0
\orphanpenalties  5    1     2     3     4     5 % 6 -> 0
\interlinepenalties 5 100000 200000 300000 400000 500000 % 6 -> 0
```

It actually makes sense to explicitly zero the last entry because as you can see in the figure the last one gets used when we run out of entries.

Can you guess what the next specification does?

```
\widowpenalties 3 options \largestspecificationoptioncode 3000 2000 1000
\clubpenalties  3 options \largestspecificationoptioncode   30   20   10
```

## 11.13 Par passes

Everything comes together in what we call par passes. Before we explain them first something about a feature that makes setting up for instance `\widowpenalties` easier. Here are a few definitions:

```
\specificationdef\strictwidowpenalties      \widowpenalties \plusthree
      \maxcount \maxcount \zerocount \relax
```

<pre>The Earth, as a habitat for animal life, is in old age and has a fatal illness. Sev- eral, in fact. It would be happening whether humans had ever evolved or not. But our presence is like the effect of an old- age patient who smokes many packs of cigarettes per day—and we humans are the cigarettes. The Earth, as a habitat for animal life, is in old age and has a fatal illness. Sev- eral, in fact. It would be happening whether humans had ever evolved or not. But our presence is like the effect of an old- age patient who smokes many packs of cigarettes per day—and we humans are the cigarettes.</pre>	<pre>The Earth, as a habitat for animal life, is in old age and has a fatal illness. Sev- eral, in fact. It would be happening whether humans had ever evolved or not. But our presence is like the effect of an old- age patient who smokes many packs of cigarettes per day—and we humans are the cigarettes. The Earth, as a habitat for animal life, is in old age and has a fatal illness. Sev- eral, in fact. It would be happening whether humans had ever evolved or not. But our presence is like the effect of an old- age patient who smokes many packs of cigarettes per day—and we humans are the cigarettes.</pre>
\normalizeparmode 8	\normalizeparmode 8

<pre>The Earth, as a habitat for animal life, is in old age and has a fatal illness. Sev- eral, in fact. It would be happening whether humans had ever evolved or not. But our presence is like the effect of an old- age patient who smokes many packs of cigarettes per day—and we humans are the cigarettes. The Earth, as a habitat for animal life, is in old age and has a fatal illness. Sev- eral, in fact. It would be happening whether humans had ever evolved or not. But our presence is like the effect of an old- age patient who smokes many packs of cigarettes per day—and we humans are the cigarettes.</pre>	<pre>The Earth, as a habitat for animal life, is in old age and has a fatal illness. Sev- eral, in fact. It would be happening whether humans had ever evolved or not. But our presence is like the effect of an old- age patient who smokes many packs of cigarettes per day—and we humans are the cigarettes. The Earth, as a habitat for animal life, is in old age and has a fatal illness. Sev- eral, in fact. It would be happening whether humans had ever evolved or not. But our presence is like the effect of an old- age patient who smokes many packs of cigarettes per day—and we humans are the cigarettes.</pre>
newline and \normalizeparmode 8	newline and \normalizeparmode 8

**Figure 11.13** Penalty lists

```
\specificationdef\strictwidowpenaltiestwo \widowpenalties \plustwo
\maxcount \zerocount \relax
\specificationdef\strictwidowpenaltiesthree \widowpenalties \plusthree
```

```

\maxcount \maxcount \zerocount \relax
\specificationdef\strictwidowpenaltiesfour \widowpenalties \plusfour
\maxcount \maxcount \maxcount \zerocount \relax

```

These are defined in the core and hooked into the alignment interface:

```

\installaligncommand{strictwidows} {\strictwidowpenalties }
\installaligncommand{strictwidows:2}{\strictwidowpenaltiestwo }
\installaligncommand{strictwidows:3}{\strictwidowpenaltiesthree}
\installaligncommand{strictwidows:4}{\strictwidowpenaltiesfour }

```

We also have four such ‘strict’ definitions for club but only one for broken penalties. For orphan penalties we have four ‘less’ orphan penalties but for widow, club and broken we have only one. So we end up with lessorphans, lessorphans:2, lessorphans:3, lessorphans:4, defaultwidows, defaultclubs, defaultbroken, strictwidows, strictwidows:2, strictwidows:3, strictwidows:4, strictclubs, strictclubs:2, strictclubs:3, strictclubs:4 and strictbroken.

You can also use `\specificationdef` for other constructs that have this multiple variable setup. Now to par passes. This is a mechanism unique to LuaMetaTeX that permits more than the usual upto three paragraph break steps: pretolerance, tolerance and emergency. How this works is explained in detail in the paragraphs chapter in the beyond document.

*todo: copy some from article when published*

## 11.13 Colofon

Author	Hans Hagen
ConTeXt	2025.07.04 21:26
LuaMetaTeX	2.11.07   20250705
Support	www.pragma-ade.com contextgarden.net ntg-context@ntg.nl

## **12 Alignments**



low level

TEX

alignments

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## 12.1 Introduction

$\text{\TeX}$  has a couple of subsystems and alignments is one of them. This mechanism is used to construct tables or alike. Because alignments use low level primitives to set up and construct a table, and because such a setup can be rather extensive, in most cases users will rely on macros that hide this.

```
\halign {
  \alignmark\hss \aligntab
  \hss\alignmark\hss \aligntab
  \hss\alignmark \cr
  1.1 \aligntab 2,2 \aligntab 3=3 \cr
  11.11 \aligntab 22,22 \aligntab 33=33 \cr
  111.111 \aligntab 222,222 \aligntab 333=333 \cr
}
```

That one doesn't look too complex and comes out as:

```
1.1      2,2      3=3
11.11    22,22    33=33
111.111  222,222 333=333
```

This is how the previous code comes out when we use one of the  $\text{Con}\text{\TeX}\text{t}$  table mechanism.

```
\starttabulate[|l|c|r|]
  \NC 1.1 \NC 2,2 \NC 3=3 \NC \NR
  \NC 11.11 \NC 22,22 \NC 33=33 \NC \NR
```

```
\NC 111.111 \NC 222,222 \NC 333=333 \NC \NR
\stoptabulate
```

```
1.1      2,2      3=3
11.11    22,22    33=33
111.111  222,222  333=333
```

That one looks a bit different with respect to spaces, so let's go back to the low level variant:

```
\halign {
    \alignmark\hss \aligntab
    \hss\alignmark\hss \aligntab
    \hss\alignmark      \cr
1.1\aligntab      2,2\aligntab      3=3\cr
11.11\aligntab    22,22\aligntab    33=33\cr
111.111\aligntab  222,222\aligntab  333=333\cr
}
```

Here we don't have spaces in the content part and therefore also no spaces in the result:

```
1.1      2,2      3=3
11.11    22,22    33=33
111.111222,222333=333
```

You can automate dealing with unwanted spacing:

```
\halign {
    \ignorespaces\alignmark\unskip\hss \aligntab
    \hss\ignorespaces\alignmark\unskip\hss \aligntab
    \hss\ignorespaces\alignmark\unskip      \cr
1.1      \aligntab 2,2      \aligntab 3=3      \cr
11.11    \aligntab 22,22    \aligntab 33=33    \cr
111.111  \aligntab 222,222  \aligntab 333=333  \cr
}
```

We get:

```
1.1      2,2      3=3
11.11    22,22    33=33
111.111222,222333=333
```

By moving the space skipping and cleanup to the so called preamble we don't need to deal with it in the content part. We can also deal with inter-column spacing there:

```

\halign {
  \ignorespaces\alignmark\unskip\hss \tabskip 1em \aligntab
  \hss\ignorespaces\alignmark\unskip\hss \tabskip 1em \aligntab
  \hss\ignorespaces\alignmark\unskip \tabskip Opt \cr
  1.1 \aligntab 2,2 \aligntab 3=3 \cr
  11.11 \aligntab 22,22 \aligntab 33=33 \cr
  111.111 \aligntab 222,222 \aligntab 333=333 \cr
}

```

```

1.1      2,2      3=3
11.11    22,22    33=33
111.111  222,222 333=333

```

If for the moment we forget about spanning columns (`\span`) and locally ignoring preamble entries (`\omit`) these basic commands are not that complex to deal with. Here we use `\alignmark` but that is just a primitive that we use instead of `#` while `\aligntab` is the same as `&`, but using the characters instead also assumes that they have the catcode that relates to a parameter and alignment tab (and in `ConTeXt` that is not the case). The `TeXbook` has plenty alignment examples so if you really want to learn about them, consult that must-have-book.

## 12.2 Between the lines

The individual rows of a horizontal alignment are treated as lines. This means that, as we see in the previous section, the interline spacing is okay. However, that also means that when we mix the lines with rules, the normal `TeX` habits kick in. Take this:

```

\halign {
  \ignorespaces\alignmark\unskip\hss \tabskip 1em \aligntab
  \hss\ignorespaces\alignmark\unskip\hss \tabskip 1em \aligntab
  \hss\ignorespaces\alignmark\unskip \tabskip Opt \cr
  \noalign{\hrule}
  1.1 \aligntab 2,2 \aligntab 3=3 \cr
  \noalign{\hrule}
  11.11 \aligntab 22,22 \aligntab 33=33 \cr
  \noalign{\hrule}
  111.111 \aligntab 222,222 \aligntab 333=333 \cr
  \noalign{\hrule}
}

```

The result doesn't look pretty and actually, when you see documents produced by T<sub>E</sub>X using alignments you should not be surprised to notice rather ugly spacing. The user (or the macropackage) should deal with that explicitly, and this is not always the case.

1.1	2,2	3=3
11.11	22,22	33=33
111.111	222,222	333=333

The solution is often easy:

```
\halign {
  \ignorespaces\strut\alignmark\unskip\hss \tabskip 1em \aligntab
  \hss\ignorespaces\strut\alignmark\unskip\hss \tabskip 1em \aligntab
  \hss\ignorespaces\strut\alignmark\unskip \tabskip 0pt \cr
  \noalign{\hrule}
  1.1 \aligntab 2,2 \aligntab 3=3 \cr
  \noalign{\hrule}
  11.11 \aligntab 22,22 \aligntab 33=33 \cr
  \noalign{\hrule}
  111.111 \aligntab 222,222 \aligntab 333=333 \cr
  \noalign{\hrule}
}
```

1.1	2,2	3=3
11.11	22,22	33=33
111.111	222,222	333=333

The user will not notice it but alignments put some pressure on the general T<sub>E</sub>X scanner. Actually, the scanner is either scanning an alignment or it expects regular text (including math). When you look at the previous example you see `\noalign`. When the preamble is read, T<sub>E</sub>X will pick up rows till it finds the final brace. Each row is added to a temporary list and the `\noalign` will enter a mode where other stuff gets added to that list. It all involves subtle look ahead but with minimal overhead. When the whole alignment is collected a final pass over that list will package the cells and rows (lines) in the appropriate way using information collected (like the maximum width of a cell and width of the current cell. It will also deal with spanning cells then.

So let's summarize what happens:

1. scan the preamble that defines the cells (where the last one is repeated when needed)
2. check for `\cr`, `\noalign` or a right brace; when a row is entered scan for cells in parallel the preamble so that cell specifications can be applied (then start again)
3. package the preamble based on information with regards to the cells in a column

4. apply the preamble packaging information to the columns and also deal with pending cell spans
5. flush the result to the current list, unless packages in a box a `\halign` is seen as paragraph and rows as lines (such a table can split)

The second (repeated) step is complicated by the fact that the scanner has to look ahead for a `\noalign`, `\cr`, `\omit` or `\span` and when doing that it has to expand what comes. This can give side effects and often results in obscure error messages. When for instance an `\if` is seen and expanded, the wrong branch can be entered. And when you use protected macros embedded alignment commands are not seen at all; of course they still need to produce valid operations in the current context.

All these side effects are to be handled in a macro package when it wraps alignments in a high level interface and Con<sub>T</sub>E<sub>X</sub>t does that for you. But because the code doesn't always look pretty then, in LuaMeta<sub>T</sub>E<sub>X</sub> the alignment mechanism has been extended a bit over time.

Nesting `\noalign` is normally not permitted (but one can redefine this primitive such that a macro package nevertheless handles it). The first extension permits nested usage of `\noalign`. This has resulted of a little reorganization of the code. A next extension showed up when overload protection was introduced and extra prefixes were added. We can signal the scanner that a macro is actually a `\noalign` variant:<sup>19</sup>

```
\noaligned\protected\def\InBetween{\noalign{...}}
```

Here the `\InBetween` macro will get the same treatment as `\noalign` and it will not trigger an error. This extension resulted in a second bit of reorganization (think of internal command codes and such) but still the original processing of alignments was there.

A third overhaul of the code actually did lead to some adaptations in the way alignments are constructed so let's move on to that.

## 12.3 Pre-, inter- and post-tab skips

The basic structure of a preamble and row is actually not that complex: it is a mix of tab skip glue and cells (that are just boxes):

```
\tabskip 10pt
```

---

<sup>19</sup> One can argue for using the name `\peekaligned` because in the meantime other alignment primitives also can use this property.

```

\halign {
  \strut\alignmark\tabskip 12pt\aligntab
  \strut\alignmark\tabskip 14pt\aligntab
  \strut\alignmark\tabskip 16pt\cr
  \noalign{\hrule}
  cell 1.1\aligntab cell 1.2\aligntab cell 1.3\cr
  \noalign{\hrule}
  cell 2.1\aligntab cell 2.2\aligntab cell 2.3\cr
  \noalign{\hrule}
}

```

The tab skips are set in advance and apply to the next cell (or after the last one).

---

<span style="font-family: monospace; font-size: small; color: gray;">TB:10.000</span> <span style="font-size: x-large; font-weight: bold;">cell 1.1</span> <span style="font-family: monospace; font-size: small; color: gray;">SP:3.497</span> <span style="font-size: x-large; font-weight: bold;">cell 1.2</span> <span style="font-family: monospace; font-size: small; color: gray;">TB:12.000</span> <span style="font-size: x-large; font-weight: bold;">cell 1.3</span> <span style="font-family: monospace; font-size: small; color: gray;">SP:3.497</span> <span style="font-size: x-large; font-weight: bold;">cell 1.3</span> <span style="font-family: monospace; font-size: small; color: gray;">TB:16.000</span>
<span style="font-family: monospace; font-size: small; color: gray;">TB:10.000</span> <span style="font-size: x-large; font-weight: bold;">cell 2.1</span> <span style="font-family: monospace; font-size: small; color: gray;">SP:3.497</span> <span style="font-size: x-large; font-weight: bold;">cell 2.2</span> <span style="font-family: monospace; font-size: small; color: gray;">TB:12.000</span> <span style="font-size: x-large; font-weight: bold;">cell 2.3</span> <span style="font-family: monospace; font-size: small; color: gray;">SP:3.497</span> <span style="font-size: x-large; font-weight: bold;">cell 2.3</span> <span style="font-family: monospace; font-size: small; color: gray;">TB:16.000</span>

---

In the ConT<sub>E</sub>Xt table mechanisms the value of `\tabskip` is zero in most cases. As in:

```

\tabskip 0pt
\halign {
  \strut\alignmark\aligntab
  \strut\alignmark\aligntab
  \strut\alignmark\cr
  \noalign{\hrule}
  cell 1.1\aligntab cell 1.2\aligntab cell 1.3\cr
  \noalign{\hrule}
  cell 2.1\aligntab cell 2.2\aligntab cell 2.3\cr
  \noalign{\hrule}
}

```

When these skips are zero, they still show up in the end:

---

cell 1.1 cell 1.2 cell 1.3

---

cell 2.1 cell 2.2 cell 2.3

---

Normally, in order to achieve certain effects there will be more align entries in the preamble than cells in the table, for instance because you want vertical lines between cells. When these are not used, you can get quite a bit of empty boxes and zero skips. Now, of course this is seldom a problem, but when you have a test document where you want to show font properties in a table and that font supports a script with some ten thousand glyphs, you can imagine that it accumulates and in LuaT<sub>E</sub>X (and LuaMetaT<sub>E</sub>X) nodes are larger so it is one of these cases where in ConT<sub>E</sub>Xt we get messages on the console that node memory is bumped.<sup>20</sup>

After playing a bit with stripping zero tab skips I found that the code would not really benefit from such a feature: lots of extra tests made it quite ugly. As a result a first alternative was to just strip zero skips before an alignment got flushed. At least we're then a bit leaner in the processes that come after it. This feature is now available as one of the normalizer bits.

But, as we moved on, a more natural approach was to keep the skips in the preamble, because that is where a guaranteed alternating skip/box is assumed. It also makes that the original documentation is still valid. However, in the rows construction we can be lean. This is driven by a keyword to `\halign`:

```
\tabskip 0pt
\halign noskips {
  \strut\alignmark\aligntab
  \strut\alignmark\aligntab
  \strut\alignmark\cr
  \noalign{\hrule}
  cell 1.1\aligntab cell 1.2\aligntab cell 1.3\cr
  \noalign{\hrule}
```

---

<sup>20</sup> I suppose it was a coincidence that a few weeks after these features came available a user consulted the mailing list about a few thousand page table that made the engine run out of memory, something that could be cured by enabling these new features.



```
cell 2.1\aligntab cell 2.2\aligntab cell 2.3\cr
\noalign{\hrule}
}
```

No zero tab skips show up here:

---

cell<sub>SP:3.497</sub>.1 cell<sub>SP:3.497</sub>.2 cell<sub>SP:3.497</sub>.3

---

cell<sub>SP:2.497</sub>.1 cell<sub>SP:2.497</sub>.2 cell<sub>SP:2.497</sub>.3

---

When playing with all this the LuaMetaTeX engine also got a tracing option for alignments. We already had one that showed some of the `\noalign` side effects, but showing the preamble was not yet there. This is what `\tracingalignments = 2` results in:

```
<preamble>
\glue[ignored][...] 0.0pt
\alignrecord
..\{\strut }
..<content>
..\{\endtemplate }
\glue[ignored][...] 0.0pt
\alignrecord
..\{\strut }
..<content>
..\{\endtemplate }
\glue[ignored][...] 0.0pt
\alignrecord
..\{\strut }
..<content>
..\{\endtemplate }
\glue[ignored][...] 0.0pt
```

The ignored subtype is (currently) only used for these alignment tab skips and it triggers a check later on when the rows are constructed. The `<content>` is what get injected in the cell (represented by `\alignmark`). The pseudo primitives are internal and not public.

## 12.4 Cell widths

Imagine this:

```
\halign {
  x\hbox to 3cm{\strut \alignmark\hss}\aligntab
  x\hbox to 3cm{\strut\hss\alignmark\hss}\aligntab
  x\hbox to 3cm{\strut\hss\alignmark } \cr
  cell 1.1\aligntab cell 1.2\aligntab cell 1.3\cr
  cell 2.1\aligntab cell 2.2\aligntab cell 2.3\cr
}
```

which renders as:

xcell 1.1	x	cell 1.2	x	cell 1.3
xcell 2.1	x	cell 2.2	x	cell 2.3

A reason to have boxes here is that it enforces a cell width but that is done at the cost of an extra wrapper. In LuaMetaTeX the `hlist` nodes are rather large because we have more options than in original TeX, for instance offsets and orientation. In a table with 10K rows of 4 cells yet get 40K extra `hlist` nodes allocated. Now, one can argue that we have plenty of memory but being lazy is not really a sign of proper programming.

```
\halign {
  x\tabsize 3cm\strut \alignmark\hss\aligntab
  x\tabsize 3cm\strut\hss\alignmark\aligntab
  x\tabsize 3cm\strut\hss\alignmark\hss\cr
  cell 1.1\aligntab cell 1.2\aligntab cell 1.3\cr
  cell 2.1\aligntab cell 2.2\aligntab cell 2.3\cr
}
```

If you look carefully you will see that this time we don't have the embedded boxes:

xcell 1.1	x	cell 1.2	x	cell 1.3
xcell 2.1	x	cell 2.2	x	cell 2.3

So, both the `sparse skip` and new `\tabsize` feature help to make these extreme tables (spanning hundreds of pages) not consume irrelevant memory and also make that later on we don't have to consult useless nodes.

## 12.5 Plugins

Yet another LuaMetaTeX extension is a callback that kicks in between the preamble pre-roll and finalizing the alignment. Initially as test and demonstration a basic character alignment feature was written but that works so well that in some places it can replace (or compliment) the already existing features in the ConTeXt table mechanisms.

```
\starttabulate[|lG{.}|cG{,}|rG{=}|cG{x}||]
\NC 1.1      \NC 2,2      \NC 3=3      \NC a 0xFF   \NC \NR
\NC 11.11   \NC 22,22   \NC 33=33   \NC b 0xFFF  \NC \NR
\NC 111.111 \NC 222,222 \NC 333=333 \NC c 0xFFFF \NC \NR
\stoptabulate
```

The tabulate mechanism in ConTeXt is rather old and stable and it is the preferred way to deal with tabular content in the text flow. However, adding the G specifier (as variant of the g one) could be done without interference or drop in performance. This new G specifier tells the tabulate mechanism that in that column the given character is used to vertically align the content that has this character.

```
1.1      2,2      3=3      a 0xFF
11.11   22,22   33=33   b 0xFFF
111.111 222,222 333=333 c 0xFFFF
```

Let's make clear that this is *not* an engine feature but a ConTeXt one. It is however made easy by this callback mechanism. We can of course use this feature with the low level alignment primitives, assuming that you tell the machinery that the plugin is to be kicked in.

```
\halign noskips \alignmentcharactertrigger \bgroup
\tabskip2em
\setalignmentcharacter.\ignorespaces\alignmark\unskip\hss \aligntab
\hss\setalignmentcharacter.\ignorespaces\alignmark\unskip\hss \aligntab
\hss\setalignmentcharacter=\ignorespaces\alignmark\unskip \aligntab
\hss \ignorespaces\alignmark\unskip\hss \cr
1.1 \aligntab 2,2 \aligntab 3=3 \aligntab \setalignmentcharacter{.}\relax 4.4\cr
11.11 \aligntab 22,22 \aligntab 33=33 \aligntab \setalignmentcharacter{,}\relax 44,44\cr
111.111 \aligntab 222,222 \aligntab 333=333 \aligntab \setalignmentcharacter{!}\relax 444!444\cr
x \aligntab x \aligntab x \aligntab \setalignmentcharacter{/}\relax /\cr
.1 \aligntab ,2 \aligntab =3 \aligntab \setalignmentcharacter{?}\relax ?4\cr
.111 \aligntab ,222 \aligntab =333 \aligntab \setalignmentcharacter{=}\relax 44=444\cr
\egroup
```

This rather verbose setup renders as:

```
1.1      2,2      3=3      4.4
11.11   22,22   33=33   44,44
```

111.111	222,222	333=333	444!444
x	x	x	/
.1	,2	=3	?4
.111	,222	=333	44=444

Using a high level interface makes sense but local control over such alignment too, so here follow some more examples. Here we use different alignment characters:

```
\starttabulate[|lG{.}|cG{,}|rG{=}|cG{x}|]
\NC 1.1      \NC 2,2      \NC 3=3      \NC a 0xFF    \NC \NR
\NC 11.11   \NC 22,22   \NC 33=33   \NC b 0xFFF   \NC \NR
\NC 111.111 \NC 222,222 \NC 333=333 \NC c 0xFFFF \NC \NR
\stoptabulate
```

1.1	2,2	3=3	a 0xFF
11.11	22,22	33=33	b 0xFFF
111.111	222,222	333=333	c 0xFFFF

In this example we specify the characters in the cells. We still need to add a specifier in the preamble definition because that will trigger the plugin.

```
\starttabulate[|lG{}|rG{}|]
\NC left      \NC right      \NC \NR
\NC \showglyphs \setalignmentcharacter{.}1.1 \NC \setalignmentcharacter{.}1.1 \NC \NR
\NC \showglyphs \setalignmentcharacter{,}11,11 \NC \setalignmentcharacter{,}11,11 \NC \NR
\NC \showglyphs \setalignmentcharacter{=}111=111 \NC \setalignmentcharacter{=}111=111 \NC \NR
\stoptabulate
```

left	right
<u>1</u> . <u>1</u>	1 . 1
<u>11</u> , <u>11</u>	11 , 11
<u>111</u> = <u>111</u>	111 = 111

You can mix these approaches:

```
\starttabulate[|lG{.}|rG{}|]
\NC left      \NC right      \NC \NR
\NC 1.1      \NC \setalignmentcharacter{.}1.1 \NC \NR
\NC 11.11   \NC \setalignmentcharacter{.}11.11 \NC \NR
\NC 111.111 \NC \setalignmentcharacter{.}111.111 \NC \NR
\stoptabulate
```

left	right
1.1	1.1

```
11.11    11.11
111.111  111.111
```

Here the already present alignment feature, that at some point in `tabulate` might use this new feature, is meant for numbers, but here we can go wild with words, although of course you need to keep in mind that we deal with typeset text, so there may be no match.

```
\starttabulate[|lG{.}|rG{.}|]
\NC foo.bar \NC foo.bar \NC \NR
\NC oo.ba   \NC oo.ba   \NC \NR
\NC o.b     \NC o.b     \NC \NR
\stoptabulate
```

```
foo.bar  foo.bar
oo.ba    oo.ba
o.b      o.b
```

This feature will only be used in know situations and those seldom involve advanced typesetting. However, the following does work:<sup>21</sup>

```
\starttabulate[|cG{d}|]
\NC \smallcaps abcdefgh \NC \NR
\NC          xdy        \NC \NR
\NC \sl          xdy    \NC \NR
\NC \tttf        xdy    \NC \NR
\NC \tfd         d      \NC \NR
\stoptabulate
```

```
abc d efgh
 x d y
 x d y
 x d y
  d
```

As always with such mechanisms, the question is “Where to stop?” But it makes for nice demos and as long as little code is needed it doesn't hurt.

<sup>21</sup> Should this be an option instead?

## 12.6 Pitfalls and tricks

The next example mixes bidirectional typesetting. It might look weird at first sight but the result conforms to what we discussed in previous paragraphs.

```
\starttabulate[|lG{.}|lG{}|]
\NC \righttoleft 1.1 \NC \righttoleft \setalignmentcharacter{.}1.1 \NC\NR
\NC 1.1 \NC \setalignmentcharacter{.}1.1 \NC\NR
\NC \righttoleft 1.11 \NC \righttoleft \setalignmentcharacter{.}1.11 \NC\NR
\NC 1.11 \NC \setalignmentcharacter{.}1.11 \NC\NR
\NC \righttoleft 1.111 \NC \righttoleft \setalignmentcharacter{.}1.111 \NC\NR
\NC 1.111 \NC \setalignmentcharacter{.}1.111 \NC\NR
\stoptabulate
```

```
1.1 1.1
1.1 1.1
11.1 11.1
1.11 1.11
111.1 111.1
1.111 1.111
```

In case of doubt, look at this:

```
\starttabulate[|lG{.}|lG{}|lG{.}|lG{}|]
\NC \righttoleft 1.1 \NC \righttoleft \setalignmentcharacter{.}1.1 \NC
1.1 \NC \setalignmentcharacter{.}1.1 \NC\NR
\NC \righttoleft 1.11 \NC \righttoleft \setalignmentcharacter{.}1.11 \NC
1.11 \NC \setalignmentcharacter{.}1.11 \NC\NR
\NC \righttoleft 1.111 \NC \righttoleft \setalignmentcharacter{.}1.111 \NC
1.111 \NC \setalignmentcharacter{.}1.111 \NC\NR
\stoptabulate
```

```
1.1 1.1 1.1 1.1
11.1 11.1 1.11 1.11
111.1 111.1 1.111 1.111
```

The next example shows the effect of `\omit` and `\span`. The first one makes that in this cell the preamble template is ignored.

```
\halign \bgroup
\tabsize 2cm\relax [\alignmark]\hss \aligntab
\tabsize 2cm\relax \hss[\alignmark]\hss \aligntab
\tabsize 2cm\relax \hss[\alignmark]\cr
1\aligntab 2\aligntab 3\cr
\omit 1\aligntab \omit 2\aligntab \omit 3\cr
1\aligntab 2\span 3\cr
1\span 2\aligntab 3\cr
```

```

1\span          2\span          3\cr
1\span \omit   2\span \omit   3\cr
\omit 1\span \omit 2\span \omit 3\cr

```

`\egroup`

Spans are applied at the end so you see a mix of templates applied.

[1]	[2]	[3]
1	2	3
[1]	[2]	[3]
[1]	[2]	[3]
[1]	[2]	[3]
[1]	[2]	[3]
[1]	23	
123		

When you define an alignment inside a macro, you need to duplicate the `\alignmark` signals. This is similar to embedded macro definitions. But in LuaMetaTeX we can get around that by using `\aligncontent`. Keep in mind that when the preamble is scanned there is no doesn't expand with the exception of the token after `\span`.

`\halign \bgroup`

```

\tabsize 2cm\relax \aligncontent\hss \aligntab
\tabsize 2cm\relax \hss\aligncontent\hss \aligntab
\tabsize 2cm\relax \hss\aligncontent\cr
1\aligntab 2\aligntab 3\cr
A\aligntab B\aligntab C\cr

```

`\egroup`

1	2	3
A	B	C

In this example we still have to be verbose in the way we align but we can do this:

```
\halign \bgroup
  \tabsize 2cm\relax \aligncontentleft \aligntab
  \tabsize 2cm\relax \aligncontentmiddle\aligntab
  \tabsize 2cm\relax \aligncontentright \cr
  1\aligntab 2\aligntab 3\cr
  A\aligntab B\aligntab C\cr
\egroup
```

Where the helpers are defined as:

```
\noaligned\protected\def\aligncontentleft
  {\ignorespaces\aligncontent\unskip\hss}

\noaligned\protected\def\aligncontentmiddle
  {\hss\ignorespaces\aligncontent\unskip\hss}

\noaligned\protected\def\aligncontentright
  {\hss\ignorespaces\aligncontent\unskip}
```

The preamble scanner see such macros as candidates for a single level expansion so it will inject the meaning and see the `\aligncontent` eventually.

1	2	3
A	B	C

The same effect could be achieved by using the `\span` prefix:

```
\def\aligncontentleft{\ignorespaces\aligncontent\unskip\hss}

\halign { ... \span\aligncontentleft ... }
```

One of the reasons for not directly using the low level `\halign` command is that it's a lot of work but by providing a set of helpers like here might change that a bit. Keep in mind that much of the above is not new in the sense that we could not achieve the same already, it's just a bit programmer friendly.

## 12.7 Rows

Alignment support is what the documented source calls 'interwoven'. When the engine scans for input it processing text, math or alignment content. While doing alignments it collects rows, and inside these cells but also deals with material that ends up in



between. In LuaMetaTeX I tried to isolate the bits and pieces as good as possible but it remains complicated (for all good reasons). Cells as well as rows are finalized after the whole alignment has been collected and processed. In the end cells and rows are boxes but till we're done they are in an 'unset' state.

Scanning starts with interpreting the preamble, and then grabbing rows. There is some nasty lookahead involved for `\noalign`, `\span`, `\omit`, `\cr` and `\crr` and that is not code one wants to tweak too much (although we did in LuaMetaTeX). This means for instance that adding 'let's start a row here' primitive is sort of tricky (but it might happen some day) which in turn means that it is not really possible to set row properties. As an experiment we can set some properties now by hijacking `\noalign` and storing them on the alignment stack (indeed: at the cost of some extra overhead and memory). This permits the following:

```
\halign {
  \hss
  \ignorespaces \alignmark \removeunwantedspaces
  \hss
  \quad \aligntab \quad
  \hss
  \ignorespaces \alignmark \removeunwantedspaces
  \hss
  \cr
  \noalign xoffset 40pt {}
  {\darkred cell one} \aligntab {\darkgray cell one} \cr
  \noalign orientation "002 {}
  {\darkgreen cell one} \aligntab {\darkblue cell one} \cr
  \noalign xoffset 40pt {}
  {\darkred cell two} \aligntab {\darkgray cell two} \cr
  \noalign orientation "002 {}
  {\darkgreen cell two} \aligntab {\darkblue cell two} \cr
  \noalign xoffset 40pt {}
  {\darkred cell three} \aligntab {\darkgray cell three} \cr
  \noalign orientation "002 {}
  {\darkgreen cell three} \aligntab {\darkblue cell three} \cr
  \noalign xoffset 40pt {}
  {\darkred cell four} \aligntab {\darkgray cell four} \cr
  \noalign orientation "002 {}
  {\darkgreen cell four} \aligntab {\darkblue cell four} \cr
}
```

cell one	cell one
cell one	cell one
cell two	cell two
cell two	cell two
cell three	cell three
cell three	cell three
cell four	cell four
cell four	cell four

The supported keywords are similar to those for boxes: source, target, anchor, orientation, shift, xoffset, yoffset, xmove and ymove. The dimensions can be prefixed by add and reset wipes all. Here is another example:

```
\halign {
  \hss
  \ignorespaces \alignmark \removeunwantedspaces
  \hss
  \quad \aligntab \quad
  \hss
  \ignorespaces \alignmark \removeunwantedspaces
  \hss
  \cr
  \noalign xmove 40pt {}
  {\darkred cell one} \aligntab {\darkgray cell one} \cr
  {\darkgreen cell one} \aligntab {\darkblue cell one} \cr
  \noalign xmove 20pt {}
  {\darkred cell two} \aligntab {\darkgray cell two} \cr
  {\darkgreen cell two} \aligntab {\darkblue cell two} \cr
  \noalign xmove 40pt {}
  {\darkred cell three} \aligntab {\darkgray cell three} \cr
  {\darkgreen cell three} \aligntab {\darkblue cell three} \cr
  \noalign xmove 20pt {}
  {\darkred cell four} \aligntab {\darkgray cell four} \cr
  {\darkgreen cell four} \aligntab {\darkblue cell four} \cr
}
```

cell one	cell one
cell one	cell one
cell two	cell two
cell two	cell two
cell three	cell three
cell three	cell three
cell four	cell four
cell four	cell four

Some more features might be added in the future as is it an interesting playground. It is to be seen how this ends up in ConT<sub>E</sub>Xt high level interfaces like tabulate.

## 12.8 Templates

The `\omit` command signals that the template should not be applied. But what if we actually want something at the left and right of the content? Here is how it's done:

```
\tabskip10pt \showboxes
```

```
\halign\bgroup
```

```
[\hss\aligncontent\hss]\aligntab
```

```
[\hss\aligncontent\hss]\aligntab
```

```
[\hss\aligncontent\hss]\cr
```

```
x\aligntab
```

```
x\aligntab
```

```
x\cr
```

```
xx\aligntab
```

```
xx\aligntab
```

```
xx\cr
```

```
xxx\aligntab
```

```
xxx\aligntab
```

```
xxx\cr
```

```
\omit oo\aligntab\omit
```

```
oo\aligntab\omit
```

```
oo\cr
```

```
xx\aligntab\realign{\hss({})\hss}xx\aligntab
```

```
xx\cr
```

```
\realign{\hss({})\hss}xx\aligntab xx\aligntab xx\cr
```

```
\egroup
```

The `\realign` command is like an `omit` but it expects two token lists that will for this cell be used instead of the ones from the preamble. While `\omit` also skips insertion of `\everytab`, here it is inserted, just like with normal preambles.

X	X	X
XX	XX	XX
XXX	XXX	XXX
oo	oo	oo
XX	(XX)	XX
(XX)	XX	XX

It will probably take a while before I'll apply this in ConT<sub>E</sub>Xt because changing existing (stable) table environment is not something done lightly.

## 12.9 Pitfalls

Alignment have a few properties that can catch you off-guard. One is the use of `\everycr`. The next example demonstrates that it is also injected after the preamble definition.

```
\everycr{\noalign{\hrule}}
\halign\bgroup \hsize 5cm \strut \alignmark\cr one\cr two\cr\egroup
```

This makes sense because it is one way to make sure that for instance a rule gets the width of the cell.

one
two

The same is of course true for a vertical align:

```
\everycr{\noalign{\vrule}}
\valign\bgroup \hsize 4cm \strut \aligncontent\cr one\cr two\cr\egroup
```

We set the width because otherwise the current text width is used.

one	two	
-----	-----	--

Something similar happens with a `\tabskip`: the value set before the alignment is used left of the first cell.

```
\tabskip10pt
\halign\bgroup \tabskip20pt\relax\aligncontent\cr x\cr x\cr \egroup
```

X
X

The `\tabskip` outside the alignment is an internal glue register so you can for instance use the prefix `\global`. However, in a preamble it is more a directive: the given value is stored with the cell. This means that the next code is invalid:

```
\tabskip10pt
\halign\bgroup \global\tabskip20pt\relax\aligncontent\cr x\cr x\cr \egroup
```

The parser looks at tokens in the preamble, sees the `\global` and appends it to the current pre-part of the cell's template. Then it sees a `\tabskip` and assigns the value

after it to the cell's skip. The token and its value just disappear, they are not appended to the template. Now, when the template is injected (and interpreted) this `\global` expects a variable next and in our case the `x` doesn't qualify. The next snippet however works okay:

```
\scratchcounter0
\halign\bgroup
  \global\tabskip40pt\relax\advance\scratchcounter\plusone\aligncontent
  \cr
  x:\the\scratchcounter\cr
  x:\the\scratchcounter\cr
  x:\the\scratchcounter\cr
\egroup
```

Here the `\global` is applied to the `advance` because the `skip` definition is *not* in the preamble.

```
x:1
x:2
x:3
```

Here is a variant:

```
\scratchcounter0
\halign\bgroup
  \global\tabskip10pt\relax\aligncontent\cr
  \advance\scratchcounter\plusone x:\the\scratchcounter\cr
  \advance\scratchcounter\plusone x:\the\scratchcounter\cr
  \advance\scratchcounter\plusone x:\the\scratchcounter\cr
\egroup
```

Again the `\global` stays and this time it ends up before the content which starts with an `\advance`.

```
x:1
x:2
x:3
```

Normally you will not need the next trickery but it shows that cells are grouped:

```
\halign\bgroup\aligncontent\cr
  1\atendofgrouped{A}\atendofgrouped{B}\cr
  2\aftergrouped {A}\aftergrouped {B}\cr
  3 \cr
```

**\egroup**

Maybe at some point I'll add something a bit more tuned for dealing with cells, but here is what you get with the above:

1AB

2

AB3

**12.10 Remark**

It can be that the way alignments are interfaced with respect to attributes is a bit different between LuaT<sub>E</sub>X and LuaMetaT<sub>E</sub>X but because the former is frozen (in order not to interfere with current usage patterns) this is something that we will deal with deep down in ConT<sub>E</sub>Xt LMTX.

In principle we can have hooks into the rows for pre and post material but it doesn't really pay of as grouping will still interfere. So for now I decided not to add these.

**12.10 Colofon**

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ConT <sub>E</sub> Xt	2025.07.04 21:26
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**13 Marks**

low level

TEX

marks



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### 13.1 Introduction

Marks are one of the subsystems of  $\text{T}_{\text{E}}\text{X}$ , as are for instance alignments and math as well as inserts which they share some properties with. Both inserts and marks put signals in the list that later on get intercepted and can be used to access stored information. In the case of inserts this is typeset materials, like footnotes, and in the case of marks it's token lists. Inserts are taken into account when breaking pages, and marks show up when a page has been broken and is presented to the output routine. Marks are used for running headers but other applications are possible.

In  $\text{MkII}$  marks are used to keep track of colors, transparencies and more properties that work across page boundaries. It permits picking up at the top of a page from where one left at the bottom of the preceding one. When  $\text{MkII}$  was written there was only one mark so on top of that a multiple mark mechanism was implemented that filtered specific marks from a collection. Later,  $\varepsilon\text{-T}_{\text{E}}\text{X}$  provided mark classes so that mechanism could be simplified. Although it is not that hard to do, this extension to  $\text{T}_{\text{E}}\text{X}$  didn't add any further features, so we can assume that there was no real demand for that.<sup>22</sup>

But, marks have some nasty limitations, so from the  $\text{ConT}_{\text{E}}\text{Xt}$  perspective there always was something to wish for. When you hide marks in boxes they will not be seen (the same is true for inserts). You cannot really reset them either. Okay, you can set them to nothing, but even then already present marks are still there. The  $\text{LuaT}_{\text{E}}\text{X}$  engine has a `\clearmarks` primitive but that works global. In  $\text{LuaMetaT}_{\text{E}}\text{X}$  a proper mark flusher is available. That engine also can work around the deeply nested disappearing marks. In addition, the current state of a mark can be queried and we have some tracing facilities.

In  $\text{MkIV}$  the engine's marks were not used at all and an alternative mechanism was written using Lua. It actually is one of the older  $\text{MkIV}$  features. It doesn't have the side

---

<sup>22</sup> This is probably true for most  $\text{LuaT}_{\text{E}}\text{X}$  and  $\text{LuaMetaT}_{\text{E}}\text{X}$  extensions, maybe example usage create retrospective demand. But one reason for picking up on engine development is that in the  $\text{ConT}_{\text{E}}\text{Xt}$  perspective we actually had some demands.

effects that native marks have but it comes at the price of more overhead, although that is bearable.

In this document we discuss marks but assume that LuaMetaT<sub>E</sub>X is used with ConT<sub>E</sub>Xt LMTX. There we experiment with using the native marks, complemented by a few Lua mechanisms, but it is to be seen if that will be either a replacement or an alternative.

## 13.2 The basics

Although the original T<sub>E</sub>X primitives are there, the plural  $\varepsilon$ -T<sub>E</sub>X mark commands are to be used. Marks, signals with token lists, are set with:

```
\marks0{This is mark 0} % equivalent to: \mark{This is mark 0}
\marks4{This is mark 4}
```

When a page has been split off, you can (normally this only makes sense in the output routine) access marks with:

```
\topmarks 4
\firstmarks4
\botmarks 4
```

A ‘top’ mark is the last one on the previous page(s), the ‘first’ and ‘bottom’ refer to the current page. A mark is a so called node, something that ends up in the current list and the token list is stored with it. The accessors are just commands and they fetch the token list from a separately managed storage. When you set or access a mark that has not yet been used, the storage is bumped to the right size, so it doesn't make sense to use e.g. `\marks 999` when there are no 998 ones too: it not only takes memory, it also makes T<sub>E</sub>X run over all these mark stores when synchronization happens. The best way to make sure that you are sparse is:

```
\newmarks\MyMark
```

Currently the first 16 marks are skipped so this makes `\MyMark` become mark 17. The reason is that we want to make sure that users who experiment with marks have some scratch marks available and don't overload system defined ones. Future versions of ConT<sub>E</sub>Xt might become more restrictive.

Marks can be cleared with:

```
\clearmarks 4
```

which clears the storage that keeps the top, first and bot marks. This happens immediately. You can delay this by putting a signal in the list:

**`\flushmarks`** 4

This (LuaMetaT<sub>E</sub>X) feature makes it for instance easy to reset marks that keep track of section (and lower) titles when a new chapter starts. Of course it still means that one has to implement some mechanism that deals with this but ConT<sub>E</sub>Xt always had that.

The current, latest assigned, value of a mark is available too:

**`\currentmarks`** 4

Using this value in for instance headers and footers makes no sense because the last node set can be on a following page.

## 13.3 Migration

In the introduction we mentioned that LuaMetaT<sub>E</sub>X has migration built in. In MkIV we have this as option too, but there it is delegated to Lua. It permits deeply nested inserts (notes) and marks (but we don't use native marks in MkIV).

Migrated marks end up in the postmigrated sublist of a box. In other lowlevel manuals we discuss these pre- and postmigrated sublists. As example we use this definition:

```
\setbox0\ vbox\ bgroup
test \marks 4 {mark 4.1}\par
test \marks 4 {mark 4.1}\par
test \marks 4 {mark 4.1}\par
\egroup
```

When we turn migration on (officially the second bit):

**`\automigrationmode"FF \showbox0`**

we get this:

```
> \box0=
2:4: \vbox[normal][...], width 483.69687, height 63.43475, depth 0.15576, direction l2r
2:4: .\list
2:4: ..\hbox[line][...], width 483.69687, height 7.48193, depth 0.15576, glue 459.20468fil, direction l2r
2:4: ...list
2:4: ....\glue[left hang][...] 0.0pt
2:4: ....\glue[left][...] 0.0pt
2:4: ....\glue[parfillsleft][...] 0.0pt
2:4: ....\par[newgraf][...], hangafter 1, hsize 483.69687, pretolerance 100, tolerance 200, adjdemerits 10000, linepenalty 10, doublehyphendemerits 10000,
  finallyhyphendemerits 5000, clubpenalty 2000, widowpenalty 2000, brokenpenalty 100, parfillskip 0.0pt plus 1.0fil, hyphenationmode 499519
2:4: ....\glue[indent][...] 0.0pt
2:4: ....\glyph[32768][...], language (n=1,l=2,r=3), hyphenationmode "79F3F, options "80, font <8: DejaVuSerif @ 11.0pt>, glyph U+000074 t
2:4: ....\glyph[32768][...], language (n=1,l=2,r=3), hyphenationmode "79F3F, options "80, font <8: DejaVuSerif @ 11.0pt>, glyph U+000065 e
2:4: ....\glyph[32768][...], language (n=1,l=2,r=3), hyphenationmode "79F3F, options "80, font <8: DejaVuSerif @ 11.0pt>, glyph U+000073 s
```

```

2:4: ....\glyph[32768][...], language (n=1,l=2,r=3), hyphenationmode "79F3F, options "80, font <8: DejaVuSerif @ 11.0pt>, glyph U+000074 t
2:4: ....\glue[space][...] 3.49658pt plus 1.74829pt minus 1.16553pt, font 8
2:4: ....\penalty[line][...] 10000
2:4: ....\glue[parfill][...] 0.0pt plus 1.0fil
2:4: ....\glue[right][...] 0.0pt
2:4: ....\glue[right hang][...] 0.0pt
2:4: ....\glue[par][...] 11.98988pt plus 3.99663pt minus 3.99663pt
2:4: ....\glue[baseline][...] 8.34883pt
2:4: ....\hbox[line][...], width 483.69687, height 7.48193, depth 0.15576, glue 459.20468fil, direction l2r
2:4: ...list
2:4: ....\glue[left hang][...] 0.0pt
2:4: ....\glue[left][...] 0.0pt
2:4: ....\glue[parfillleft][...] 0.0pt
2:4: ....\par[newgraf][...], hangafter 1, hsize 483.69687, pretolerance 100, tolerance 200, adjdemerits 10000, linepenalty 10, doublehyphendemerits 10000,
  finalhyphendemerits 5000, clubpenalty 2000, widowpenalty 2000, brokenpenalty 100, parfillskip 0.0pt plus 1.0fil, hyphenationmode 499519
2:4: ....\glue[indent][...] 0.0pt
2:4: ....\glyph[32768][...], language (n=1,l=2,r=3), hyphenationmode "79F3F, options "80, font <8: DejaVuSerif @ 11.0pt>, glyph U+000074 t
2:4: ....\glyph[32768][...], language (n=1,l=2,r=3), hyphenationmode "79F3F, options "80, font <8: DejaVuSerif @ 11.0pt>, glyph U+000065 e
2:4: ....\glyph[32768][...], language (n=1,l=2,r=3), hyphenationmode "79F3F, options "80, font <8: DejaVuSerif @ 11.0pt>, glyph U+000073 s
2:4: ....\glyph[32768][...], language (n=1,l=2,r=3), hyphenationmode "79F3F, options "80, font <8: DejaVuSerif @ 11.0pt>, glyph U+000074 t
2:4: ....\glue[space][...] 3.49658pt plus 1.74829pt minus 1.16553pt, font 8
2:4: ....\penalty[line][...] 10000
2:4: ....\glue[parfill][...] 0.0pt plus 1.0fil
2:4: ....\glue[right][...] 0.0pt
2:4: ....\glue[right hang][...] 0.0pt
2:4: ....\glue[par][...] 11.98988pt plus 3.99663pt minus 3.99663pt
2:4: ....\glue[baseline][...] 8.34883pt
2:4: ....\hbox[line][...], width 483.69687, height 7.48193, depth 0.15576, glue 459.20468fil, direction l2r
2:4: ...list
2:4: ....\glue[left hang][...] 0.0pt
2:4: ....\glue[left][...] 0.0pt
2:4: ....\glue[parfillleft][...] 0.0pt
2:4: ....\par[newgraf][...], hangafter 1, hsize 483.69687, pretolerance 100, tolerance 200, adjdemerits 10000, linepenalty 10, doublehyphendemerits 10000,
  finalhyphendemerits 5000, clubpenalty 2000, widowpenalty 2000, brokenpenalty 100, parfillskip 0.0pt plus 1.0fil, hyphenationmode 499519
2:4: ....\glue[indent][...] 0.0pt
2:4: ....\glyph[32768][...], language (n=1,l=2,r=3), hyphenationmode "79F3F, options "80, font <8: DejaVuSerif @ 11.0pt>, glyph U+000074 t
2:4: ....\glyph[32768][...], language (n=1,l=2,r=3), hyphenationmode "79F3F, options "80, font <8: DejaVuSerif @ 11.0pt>, glyph U+000065 e
2:4: ....\glyph[32768][...], language (n=1,l=2,r=3), hyphenationmode "79F3F, options "80, font <8: DejaVuSerif @ 11.0pt>, glyph U+000073 s
2:4: ....\glyph[32768][...], language (n=1,l=2,r=3), hyphenationmode "79F3F, options "80, font <8: DejaVuSerif @ 11.0pt>, glyph U+000074 t
2:4: ....\glue[space][...] 3.49658pt plus 1.74829pt minus 1.16553pt, font 8
2:4: ....\penalty[line][...] 10000
2:4: ....\glue[parfill][...] 0.0pt plus 1.0fil
2:4: ....\glue[right][...] 0.0pt
2:4: ....\glue[right hang][...] 0.0pt
2:4: \postmigrated
2:4: \mark[4][...]
2:4: ..{mark 4.1}
2:4: \mark[4][...]
2:4: ..{mark 4.1}
2:4: \mark[4][...]
2:4: ..{mark 4.1}

```

When we don't migrate, enforced with:

`\automigrationmode"00 \showbox0`

the result is:

```

> \box0=
2:4: \vbox[normal][...], width 483.69687, height 63.43475, depth 0.15576, direction l2r
2:4: \list
2:4: \hbox[line][...], width 483.69687, height 7.48193, depth 0.15576, glue 459.20468fil, direction l2r
2:4: ...list
2:4: ....\glue[left hang][...] 0.0pt
2:4: ....\glue[left][...] 0.0pt
2:4: ....\glue[parfillleft][...] 0.0pt
2:4: ....\par[newgraf][...], hangafter 1, hsize 483.69687, pretolerance 100, tolerance 200, adjdemerits 10000, linepenalty 10, doublehyphendemerits 10000,
  finalhyphendemerits 5000, clubpenalty 2000, widowpenalty 2000, brokenpenalty 100, parfillskip 0.0pt plus 1.0fil, hyphenationmode 499519
2:4: ....\glue[indent][...] 0.0pt
2:4: ....\glyph[32768][...], language (n=1,l=2,r=3), hyphenationmode "79F3F, options "80, font <8: DejaVuSerif @ 11.0pt>, glyph U+000074 t
2:4: ....\glyph[32768][...], language (n=1,l=2,r=3), hyphenationmode "79F3F, options "80, font <8: DejaVuSerif @ 11.0pt>, glyph U+000065 e
2:4: ....\glyph[32768][...], language (n=1,l=2,r=3), hyphenationmode "79F3F, options "80, font <8: DejaVuSerif @ 11.0pt>, glyph U+000073 s
2:4: ....\glue[space][...] 3.49658pt plus 1.74829pt minus 1.16553pt, font 8
2:4: ....\penalty[line][...] 10000
2:4: ....\glue[parfill][...] 0.0pt plus 1.0fil
2:4: ....\glue[right][...] 0.0pt
2:4: ....\glue[right hang][...] 0.0pt
2:4: \mark[4][...]
2:4: ..{mark 4.1}
2:4: \glue[par][...] 11.98988pt plus 3.99663pt minus 3.99663pt
2:4: \glue[baseline][...] 8.34883pt
2:4: \hbox[line][...], width 483.69687, height 7.48193, depth 0.15576, glue 459.20468fil, direction l2r

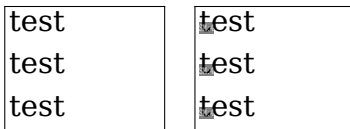
```

```

2:4: ...\list
2:4: ...\glue[left hang][...] 0.0pt
2:4: ...\glue[left][...] 0.0pt
2:4: ...\glue[parfillleft][...] 0.0pt
2:4: ...\par[newgraf][...], hangafter 1, hsize 483.69687, pretolerance 100, tolerance 200, adjdemerits 10000, linepenalty 10, doublehyphendemerits 10000,
  finalhyphendemerits 5000, clubpenalty 2000, widowpenalty 2000, brokenpenalty 100, parfillskip 0.0pt plus 1.0fil, hyphenationmode 499519
2:4: ...\glue[indent][...] 0.0pt
2:4: ...\glyph[32768][...], language (n=1,l=2,r=3), hyphenationmode "79F3F, options "80, font <8: DejaVuSerif @ 11.0pt>, glyph U+000074 t
2:4: ...\glyph[32768][...], language (n=1,l=2,r=3), hyphenationmode "79F3F, options "80, font <8: DejaVuSerif @ 11.0pt>, glyph U+000065 e
2:4: ...\glyph[32768][...], language (n=1,l=2,r=3), hyphenationmode "79F3F, options "80, font <8: DejaVuSerif @ 11.0pt>, glyph U+000073 s
2:4: ...\glyph[32768][...], language (n=1,l=2,r=3), hyphenationmode "79F3F, options "80, font <8: DejaVuSerif @ 11.0pt>, glyph U+000074 t
2:4: ...\glue[space][...] 3.49658pt plus 1.74829pt minus 1.16553pt, font 8
2:4: ...\penalty[line][...] 10000
2:4: ...\glue[parfill][...] 0.0pt plus 1.0fil
2:4: ...\glue[right][...] 0.0pt
2:4: ...\glue[right hang][...] 0.0pt
2:4: ...\mark[4][...]
2:4: ...{mark 4.1}
2:4: ...\glue[par][...] 11.98988pt plus 3.99663pt minus 3.99663pt
2:4: ...\glue[baseline][...] 8.34883pt
2:4: ...\hbox[line][...], width 483.69687, height 7.48193, depth 0.15576, glue 459.20468fil, direction l2r
2:4: ...\list
2:4: ...\glue[left hang][...] 0.0pt
2:4: ...\glue[left][...] 0.0pt
2:4: ...\glue[parfillleft][...] 0.0pt
2:4: ...\par[newgraf][...], hangafter 1, hsize 483.69687, pretolerance 100, tolerance 200, adjdemerits 10000, linepenalty 10, doublehyphendemerits 10000,
  finalhyphendemerits 5000, clubpenalty 2000, widowpenalty 2000, brokenpenalty 100, parfillskip 0.0pt plus 1.0fil, hyphenationmode 499519
2:4: ...\glue[indent][...] 0.0pt
2:4: ...\glyph[32768][...], language (n=1,l=2,r=3), hyphenationmode "79F3F, options "80, font <8: DejaVuSerif @ 11.0pt>, glyph U+000074 t
2:4: ...\glyph[32768][...], language (n=1,l=2,r=3), hyphenationmode "79F3F, options "80, font <8: DejaVuSerif @ 11.0pt>, glyph U+000065 e
2:4: ...\glyph[32768][...], language (n=1,l=2,r=3), hyphenationmode "79F3F, options "80, font <8: DejaVuSerif @ 11.0pt>, glyph U+000073 s
2:4: ...\glyph[32768][...], language (n=1,l=2,r=3), hyphenationmode "79F3F, options "80, font <8: DejaVuSerif @ 11.0pt>, glyph U+000074 t
2:4: ...\glue[space][...] 3.49658pt plus 1.74829pt minus 1.16553pt, font 8
2:4: ...\penalty[line][...] 10000
2:4: ...\glue[parfill][...] 0.0pt plus 1.0fil
2:4: ...\glue[right][...] 0.0pt
2:4: ...\glue[right hang][...] 0.0pt
2:4: ...\mark[4][...]
2:4: ...{mark 4.1}

```

When you say `\showmakeup` or in this case `\showmakeup[mark]` the marks are visualized:



enabled

disabled

Here `sm` means ‘set mark’ while `rm` would indicate a ‘reset mark’. Of course migrated marks don't show up because these are bound to the box and thereby have become a specific box property as can be seen in the above trace.

## 13.4 Tracing

The LuaMetaTeX engine has a dedicated tracing option for marks. The fact that the traditional engine doesn't have this can be seen as indication that this is seldom needed.

`\tracingmarks1`

`\tracingonline2`

When tracing is set to 1 we get a list of marks for the just split of page:

```
2:7: <mark class 51, top := bot>
```

```

2:7: ..{sample 9.1}
2:7: <mark class 51: first := mark>
2:7: ..{sample 10.1}
2:7: <mark class 51: bot := mark>
2:7: ..{sample 10.1}
2:7: <mark class 51, page state>
2:7: ..top {sample 9.1}
2:7: ..first {sample 10.1}
2:7: ..bot {sample 10.1}

```

When tracing is set to 2 you also get details we get a list of marks of the analysis:

```

1:9: <mark class 51, top := bot>
1:9: ..{sample 5.1}
1:9: <mark class 51: first := mark>
1:9: ..{sample 6.1}
1:9: <mark class 51: bot := mark>
1:9: ..{sample 6.1}
1:9: <mark class 51: bot := mark>
1:9: ..{sample 7.1}
1:9: <mark class 51: bot := mark>
1:9: ..{sample 8.1}
1:9: <mark class 51: bot := mark>
1:9: ..{sample 9.1}
1:9: <mark class 51, page state>
1:9: ..top {sample 5.1}
1:9: ..first {sample 6.1}
1:9: ..bot {sample 9.1}

```

## 13.5 High level commands

I think that not that many users define their own marks. They are useful for showing section related titles in headers and footers but the implementation of that is hidden. The native mark references are `top`, `first` and `bottom` but in the `ConTeXt` interface we use different keywords.

<b>ConTeXt</b>	<b>TeX</b>	<b>column</b>	<b>page</b>
previous	top	last before sync	last on previous page
top	first	first in sync	first on page
bottom	bot	last in sync	last on page

first	top	first not top in sync	first on page
last	bot	last not bottom in sync	last on page
<hr/>			
default	the same as first		
current	the last set value		
<hr/>			

In order to separate marks in ConT<sub>E</sub>Xt from those in T<sub>E</sub>X, the term ‘marking’ is used. In MkIV the regular marks mechanism is of course there but, as mentioned, not used. By using a different namespace we could make the transition from MkII to MkIV (the same is true for some more mechanisms).

A marking is defined with

```
\definemarking[MyMark]
```

A defined marking can be set with two equivalent commands:

```
\setmarking[MyMark]{content}
\marking [MyMark]{content}
```

The content is not typeset but stored as token list. In the sectioning mechanism that uses markings we don't even store titles, we store a reference to a title. In order to use that (deep down) we hook in a filter command. By default that command does nothing:

```
\setupmarking[MyMark] [filtercommand=\firstofoneargument]
```

The token list does *not* get expanded by default, unless you set it up:

```
\setupmarking[MyMark] [expansion=yes]
```

The current state of a marking can be cleared with:

```
\clearmarking[MyMark]
```

but because that en is not synchronized the real deal is:

```
\resetmarking[MyMark]
```

Be aware that it introduces a node in the list. You can test if a marking is defined with (as usual) a test macro. Contrary to (most) other test macros this one is fully expandable:

```
\doifelsemarking {MyMark} {
  defined
} {
  undefined
```

```
}
```

Because there can be a chain involved, we can relate markings. Think of sections below chapters and subsections below sections:

```
\relatemarking[MyMark][YourMark]
```

When a marking is set its relatives are also reset, so setting YourMark will reset MyMark. It is this kind of features that made for marks being wrapped into high level commands very early in the ConT<sub>E</sub>Xt development (and one can even argue that this is why a package like ConT<sub>E</sub>Xt exists in the first place).

The rest of the (relatively small) repertoire of commands has to do with fetching markings. The general command is `\getmarking` that takes two or three arguments:

```
\getmarking[MyMarking][first]
\getmarking[MyMarking][page][first]
\getmarking[MyMarking][page][first]
\getmarking[MyMarking][column:1][first]
```

There are (normally) three marks that can be fetched so we have three commands that do just that:

```
\fetchonemark [MyMarking][which one]
\fetchtwomarks[MyMarking]
\fetchallmarks[MyMarking]
```

You can setup a separator key which by default is:

```
\setupmarking[MyMarking][separator=\space\emdash\space]
```

Injection is enabled by default due to this default:

```
\setupmarking[MyMarking][state=start]
```

The following three variants are (what is called) fully expandable:

```
\fetchonemarking [MyMarking][which one]
\fetchtwomarkings[MyMarking]
\fetchallmarkings[MyMarking]
```



## 13.6 Pitfalls

The main pitfall is that a (re)setting a mark will inject a node which in vertical mode can interfere with spacing. In for instance section commands we wrap them with the title so there it should work out okay.

## 13.6 Colofon

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ConT <sub>E</sub> Xt	2025.07.04 21:26
LuaMetaT <sub>E</sub> X	2.11.07   20250705
Support	<a href="http://www.pragma-ade.com">www.pragma-ade.com</a> <a href="http://contextgarden.net">contextgarden.net</a> <a href="mailto:ntg-context@ntg.nl">ntg-context@ntg.nl</a>

## **14 Inserts**

low level

TEX

inserts

## Contents

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### 14.1 Introduction

This document is a mixed bag. We do discuss inserts but also touch elements of the page builder because inserts and regular page content are handled there. Examples of mechanisms that use inserts are footnotes. These have an anchor in the running text and some content that ends up (normally) at the bottom of the page. When considering a page break the engine tries to make sure that the anchor (reference) and the content end up on the same page. When there is too much, it will distribute (split) the content over pages.

We can discuss page breaks in a (pseudo) scientific way and explore how to optimize this process, taking into accounts also inserts that contain images but it doesn't make much sense to do that because in practice we can encounter all kind of interferences. Theory and practice are too different because a document can contain a wild mix of text, figures, formulas, notes, have backgrounds and location dependent processing. It get seven more complex when we are dealing with columns because  $\text{T}_{\text{E}}\text{X}$  doesn't really know that concept.

I will therefore stick to some practical aspects and the main reason for this document is that I sort of document engine features and at the same time give an impression of what we deal with. I will do that in the perspective of  $\text{LuaMetaT}_{\text{E}}\text{X}$ , which has a few more options and tracing than other engines.

*Currently this document is mostly for myself to keep track of the state of inserts and the page builder in  $\text{LuaMetaT}_{\text{E}}\text{X}$  and  $\text{ConT}_{\text{E}}\text{Xt LMTX}$ . The text is not yet corrected and can have errors.*

### 14.2 The page builder

When your document is processed content eventually gets added to the so called main vertical list (mvl). Content first get appended to the list of contributions and at specific

moments it will be handed over to the mvl. This process is called page building. There we can encounter the following elements (nodes):

glue	a vertical skip
penalty	a vertical penalty
kern	a vertical kern
vlist	a a vertical box
hlist	a horizontal box (often a line)
rule	a horizontal rule
boundary	a boundary node
whatsit	a node that is used by user code (often some extension)
mark	a token list (as used for running headers)
insert	a node list (as used for notes)

The engine itself will not insert anything other than this but Lua code can mess up the contribution list and the mvl and that can trigger an error. Handing over the contributions is done by the page builder and that one kicks in in several places:

- When a penalty gets inserted it is part of evaluating if the output routine should be triggered. This triggering can be enforced by values equal or below 10.000 that then can be checked in the set routine.
- The builder is *not* exercised when a glue or kern is injected so there can be multiple of them before another element triggers the builder.
- Adding a box triggers the builder as does the result of an alignment which can be a list of boxes.
- When the output routine is finished the builder is executed because the routine can have pushed back content.
- When a new paragraph is triggered by the `\par` command the builder kicks in but only when the engine was able to enter vertical mode.
- When the job is finished the builder will make sure that pending content is handled.
- An insert and `vadjust can` trigger the builder but only at the nesting level zero which normally is not the case (I need an example).
- At the beginning of a paragraph (like text), before display math is entered, and when display math ends the builder is also activated.

At the  $\text{T}_{\text{E}}\text{X}$  the builder is triggered automatically in the mentioned cases but at the Lua end you can use `tex.triggerbuildpage()` to flush the pending contributions.

The properties that relate to the page look like counter and dimension registers but they are not. These variables are global and managed differently.

<code>\pagegoal</code>	the available space
<code>\pagetotal</code>	the accumulated space
<code>\pagestretch</code>	the possible zero order stretch
<code>\pagefilstretch</code>	the possible one order stretch
<code>\pagefillstretch</code>	the possible second order stretch
<code>\pagefilllstretch</code>	the possible third order stretch
<code>\pageshrink</code>	the possible shrink
<code>\pagedepth</code>	the current page depth
<code>\pagevsize</code>	the initial page goal

When the first content is added to an empty page the `\pagegoal` gets the value of `\vsize` and gets frozen but the value is diminished by the space needed by left over inserts. These inserts are managed via a separate list so they don't interfere with the page that itself of course can have additional inserts. The `\pagevsize` is just a (LuaMeta-TeX) status variable that hold the initial `\pagegoal` but it might play a role in future extensions.

Another variable is `\deadcycles` that registers the number of times the output routine is called without returning result.

### 14.3 Inserts

We now come to inserts. In traditional TeX an insert is a data structure that runs on top of registers: a box, count, dimension and skip. An insert is accessed by a number so for instance `insert 123` will use the four registers of that number. Because TeX only offers a command alias mechanism for registers (like `\countdef`) a macro package will implement some allocator management subsystem (like `\newcount`). A `\newinsert` has to be defined in a way that the four registers are not clashing with other allocators. When you start with TeX seeing code that deals with in (in plain TeX) can be puzzling but it follows from the way TeX is set up. But inserts are probably not what you start exploring right away away.

In LuaMetaTeX you can set `\insertmode` to 1 and that is what we do in ConTeXt. In that mode inserts are taken from a pool instead of registers. A side effect is that like the page properties the insert properties are global too but that is normally no problem and can be managed well by a macro package (that probably would assign register the values globally too). The insert pool will grow dynamically on demand so one can just start at 1; in ConTeXt MkIV we use the range 127 upto 255 in order to avoid a clash with registers. In LMTX we start at 1 because there are no clashes.

A consequence of this approach is that we use dedicated commands to set the insert properties:

<code>\insertdistance</code>	glue	the space before the first instance (on a page)
<code>\insertmultiplier</code>	count	a factor that is used to calculate the height used
<code>\insertlimit</code>	dimen	the maximum amount of space on a page to be taken
<code>\insertpenalty</code>	count	the floating penalty (used when set)
<code>\insertmaxdepth</code>	dimen	the maximum split depth (used when set)
<code>\insertstorage</code>	count	signals that the insert has to be stored for later
<code>\insertheight</code>	dimen	the accumulated height of the inserts so far
<code>\insertdepth</code>	dimen	the current depth of the inserts so far
<code>\insertwidth</code>	dimen	the width of the inserts

These commands take a number and an integer, dimension or glue specification. They can be set and queried but setting the dimensions can have side effects. The accumulated height of the inserts is available in `\insertheights` (which can be set too). The `\floatingpenalty` variable determines the penalty applied when a split is needed.

In the output routine the original  $\TeX$  variable `\insertpenalties` is a counter that keeps the number of insertions that didn't fit on the page while otherwise it has the accumulated penalties of the split insertions. When `\holdinginserts` is non zero the inserts in the list are not collected for output, which permits the list to be fed back for reprocessing.

The LuaMeta $\TeX$  specific storage mode `\insertstoring` variable is explained in the next section.

## 14.4 Storing

This feature is kind of special and still experimental. When `\insertstoring` is set 1, all inserts that have their storage flag set will be saved. Think of a multi column setup where inserts have to end up in the last column. If there are three columns, the first two will store inserts. Then when the last column is dealt with `\insertstoring` can be set to 2 and that will signal the builder that we will inject the inserts. In both cases, the value of this register will be set to zero so that it doesn't influence further processing.

## 14.5 Synchronizing

The page builder can be triggered by (for instance) a penalty but you can also use `\pageboundary`. This will trigger the page builder but not leave anything behind. (This is experimental.)

## 14.6 Migration

*Todo, nothing new there, so no hurry.*

## 14.7 Callbacks

*Todo, nothing new there, so no hurry.*

## 14.7 Colofon

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ConT <sub>E</sub> Xt	2025.07.04 21:26
LuaMetaT <sub>E</sub> X	2.11.07   20250705
Support	<a href="http://www.pragma-ade.com">www.pragma-ade.com</a> <a href="http://contextgarden.net">contextgarden.net</a> <a href="mailto:ntg-context@ntg.nl">ntg-context@ntg.nl</a>



## **15 Localboxes**

low level

TEX

localboxes

## Contents

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### 15.1 Introduction

The LuaTeX engine inherited a few features from other engines and adding local boxes to paragraphs is one of them. This concept comes from Omega but over time it has been made a bit more robust, also by using native par nodes instead of whatsit nodes that are used to support TeX's extensions. In another low level manual we discuss paragraph properties and these local boxes are also part of that so you might want to catch up on that. Local boxes are stored in an initial par node with an adequate subtype but users won't notice this (unless they mess around in Lua). The inline par nodes have a different subtype and are injected with the `\localinterlinepenalty`, `\localbrokenpenalty`, `\localleftbox`, `\localrightbox` and LuaMetaTeX specific `\localmiddlebox` primitives. When these primitives are used in vertical mode they just set registers.

The original (Omega) idea was that local boxes are used for repetitive punctuation (like quotes) at the left and/or right end of the lines that make up a paragraph. That means that when these primitives inject nodes they actually introduce states so that a stretch of text can be marked.

When this mechanism was cleaned up in LuaMetaTeX I decided to investigate if other usage made sense. After all, it is a feature that introduces some extra code and it then pays off to use it when possible. Among the extensions are a callback that is triggered when the left and right boxes get added and experiments with that showed some potential but in order to retain performance as well as limit extensive node memory usage (par nodes are large) a system of indices was added. All this will be illustrated below. Warning: the mechanism in LuaMetaTeX is not compatible with LuaTeX.

*This is a preliminary, uncorrected manual.*

### 15.2 The basics

This mechanism uses a mix of setting (pseudo horizontal) box registers that get associated with (positions in a) paragraph. When the lines resulting from breaking the list gets packaged into an horizontal (line) box, the local left and right boxes get prepended

and appended to the textual part (inside the left, right and parfills kips and left or right hanging margins). When assigning the current local boxes to the paragraph node(s) references to the pseudo registers are used and the packaging actually copies them. This mix of referencing and copying is somewhat tricky but the engine does it best to hide this for the user.

This mechanism is rather useless when not wrapped into some high level mechanism because by default setting these boxes wipes the existing value. In LuaMetaT<sub>E</sub>X you can actually access the boxes so prepending and appending is possible but experiments showed that this could come with a huge performance hit when the lists are not cleaned up during a run. This is why we have introduced indices: when you assign local boxes using the index option that specific index will be replaced and therefore we have a more sparse solution. So, contrary to LuaT<sub>E</sub>X, in LuaMetaT<sub>E</sub>X the local box registers have a linked lists of local boxes tagged by index. Unless you manipulate in Lua, this is hidden from the user. One can access the boxes from the T<sub>E</sub>X the but there can be no confusion with LuaT<sub>E</sub>X here because there we don't have access. This is why usage as in LuaT<sub>E</sub>X will also work in LuaMetaT<sub>E</sub>X.

This mechanism obeys grouping as is demonstrated in the next three examples. The first example is:

```
\start
  \dorecurse{10}{test #1.1 }
  \localleftbox{\blackrule[width=2em,color=darkred] }
  \dorecurse{20}{test #1.2 }
  \removeunwantedspaces
  \localrightbox{ \blackrule[width=3em,color=darkblue]}
  \dorecurse{20}{test #1.3 }
\stop
  \dorecurse{20}{test #1.4 }
  % par ends here
```

The next example differs in a subtle way: watch the keep keyword, it makes the setting retain after the group ends.

```
\start
  \start
    \dorecurse{10}{test #1.1 }
    \localleftbox keep {\blackrule[width=2em,color=darkred] }
    \dorecurse{20}{test #1.2 }
    \removeunwantedspaces
    \localrightbox { \blackrule[width=3em,color=darkblue]}
```

```

    \dorecurse{20}{test #1.3 }
  \stop
    \dorecurse{20}{test #1.4 }
\stop
% par ends here

```

The third example has two times keep. This option is LuaMetaTeX specific.

```

\start
  \start
    \dorecurse{10}{test #1.1 }
    \localleftbox keep { \blackrule[width=2em,color=darkred] }
    \dorecurse{20}{test #1.2 }
    \removeunwantedspaces
    \localrightbox keep { \blackrule[width=3em,color=darkblue] }
    \dorecurse{20}{test #1.3 }
  \stop
    \dorecurse{20}{test #1.4 }
\stop
% par ends here

```

One (nasty) side effect is that when you set these boxes ungrouped they are applied to whatever follows, which is why resetting them is built in the relevant parts of ConTeXt. The next examples are typeset grouped and demonstrate the use of indices:

```

\dorecurse{20}{before #1 }
\localleftbox{\bf \darkred L 1 }%
\localleftbox{\bf \darkred L 2 }%
\dorecurse{20}{after #1 }

```

before 1 before 2 before 3 before 4 before 5 before 6 before 7 before 8 before 9 before  
 10 before 11 before 12 before 13 before 14 before 15 before 16 before 17 before 18  
 before 19 before 20 after 1 after 2 after 3 after 4 after 5 after 6 after 7 after 8 after 9  
**L 2** after 10 after 11 after 12 after 13 after 14 after 15 after 16 after 17 after 18 after  
**L 2** 19 after 20

Indices can be set for both sides:

```

\dorecurse{5}{\localrightbox index #1{ \bf \darkgreen R #1}}%
\dorecurse{20}{before #1 }
\dorecurse{5}{\localleftbox index #1{\bf \darkred L #1 }}%
\dorecurse{20}{after #1 }

```

test 1.1 test 2.1 test 3.1 test 4.1 test 5.1 test 6.1 test 7.1 test 8.1 test 9.1 test 10.1 test  
 ■ 1.2 test 2.2 test 3.2 test 4.2 test 5.2 test 6.2 test 7.2 test 8.2 test 9.2 test 10.2 test  
 ■ 11.2 test 12.2 test 13.2 test 14.2 test 15.2 test 16.2 test 17.2 test 18.2 test 19.2  
 ■ test 20.2 test 1.3 test 2.3 test 3.3 test 4.3 test 5.3 test 6.3 test 7.3 test 8.3 ■  
 ■ test 9.3 test 10.3 test 11.3 test 12.3 test 13.3 test 14.3 test 15.3 test 16.3 ■  
 ■ test 17.3 test 18.3 test 19.3 test 20.3 test 1.4 test 2.4 test 3.4 test 4.4 test 5.4 test  
 6.4 test 7.4 test 8.4 test 9.4 test 10.4 test 11.4 test 12.4 test 13.4 test 14.4 test 15.4 test  
 16.4 test 17.4 test 18.4 test 19.4 test 20.4

### Example 1

test 1.1 test 2.1 test 3.1 test 4.1 test 5.1 test 6.1 test 7.1 test 8.1 test 9.1 test 10.1 test  
 ■ 1.2 test 2.2 test 3.2 test 4.2 test 5.2 test 6.2 test 7.2 test 8.2 test 9.2 test 10.2 test  
 ■ 11.2 test 12.2 test 13.2 test 14.2 test 15.2 test 16.2 test 17.2 test 18.2 test 19.2  
 ■ test 20.2 test 1.3 test 2.3 test 3.3 test 4.3 test 5.3 test 6.3 test 7.3 test 8.3 ■  
 ■ test 9.3 test 10.3 test 11.3 test 12.3 test 13.3 test 14.3 test 15.3 test 16.3 ■  
 ■ test 17.3 test 18.3 test 19.3 test 20.3 test 1.4 test 2.4 test 3.4 test 4.4 test 5.4 test  
 6.4 test 7.4 test 8.4 test 9.4 test 10.4 test 11.4 test 12.4 test 13.4 test 14.4 test 15.4 test  
 16.4 test 17.4 test 18.4 test 19.4 test 20.4

### Example 2

test 1.1 test 2.1 test 3.1 test 4.1 test 5.1 test 6.1 test 7.1 test 8.1 test 9.1 test 10.1 test  
 ■ 1.2 test 2.2 test 3.2 test 4.2 test 5.2 test 6.2 test 7.2 test 8.2 test 9.2 test 10.2 test  
 ■ 11.2 test 12.2 test 13.2 test 14.2 test 15.2 test 16.2 test 17.2 test 18.2 test 19.2 test  
 ■ 20.2 test 1.3 test 2.3 test 3.3 test 4.3 test 5.3 test 6.3 test 7.3 test 8.3 test ■  
 ■ 9.3 test 10.3 test 11.3 test 12.3 test 13.3 test 14.3 test 15.3 test 16.3 test ■  
 ■ 17.3 test 18.3 test 19.3 test 20.3 test 1.4 test 2.4 test 3.4 test 4.4 test 5.4 ■  
 ■ test 6.4 test 7.4 test 8.4 test 9.4 test 10.4 test 11.4 test 12.4 test 13.4 test ■  
 ■ 14.4 test 15.4 test 16.4 test 17.4 test 18.4 test 19.4 test 20.4 ■

### Example 3

#### Figure 15.1

before 1 before 2 before 3 before 4 before 5 before 6 before 7 **R 1 R 2 R 3 R 4 R 5**  
 before 8 before 9 before 10 before 11 before 12 before 13 before **R 1 R 2 R 3 R 4 R 5**  
 14 before 15 before 16 before 17 before 18 before 19 before 20 **R 1 R 2 R 3 R 4 R 5**  
 after 1 after 2 after 3 after 4 after 5 after 6 after 7 after 8 after **R 1 R 2 R 3 R 4 R 5**  
**L 1 L 2 L 3 L 4 L 5** 9 after 10 after 11 after 12 after 13 after 14 **R 1 R 2 R 3 R 4 R 5**  
**L 1 L 2 L 3 L 4 L 5** after 15 after 16 after 17 after 18 after 19 **R 1 R 2 R 3 R 4 R 5**  
**L 1 L 2 L 3 L 4 L 5** after 20 **R 1 R 2 R 3 R 4 R 5**

We can instruct this mechanism to hook the local box into the main par node by using the `par` keyword. Keep in mind that these local boxes only come into play when the lines are broken, so till then changing them is possible.

```
\dorecurse{3}{\localrightbox index #1{ \bf \darkgreen R #1}}%
\dorecurse{20}{before #1 }
\dorecurse{2}{\localleftbox par index #1{\bf \darkred L #1 }}%
\dorecurse{20}{after #1 }
```

```
L 1 L 2 before 1 before 2 before 3 before 4 before 5 before 6 before 7 before R 1 R 2 R 3
L 1 L 2 8 before 9 before 10 before 11 before 12 before 13 before 14 before R 1 R 2 R 3
L 1 L 2 15 before 16 before 17 before 18 before 19 before 20 after 1 after 2 R 1 R 2 R 3
L 1 L 2 after 3 after 4 after 5 after 6 after 7 after 8 after 9 after 10 after 11 R 1 R 2 R 3
L 1 L 2 after 12 after 13 after 14 after 15 after 16 after 17 after 18 after 19 R 1 R 2 R 3
L 1 L 2 after 20 R 1 R 2 R 3
```

## 15.3 The interface

*The interface described here is experimental.*

Because it is hard to foresee if this mechanism will be used at all the ConT<sub>E</sub>Xt interface is somewhat low level: one can build functionality on top of it. In the previous section we saw examples of local boxes being part of the text but one reason for extending the interface was to see if we can also use this engine feature for efficiently placing marginal content.

```
\definelocalboxes
  [lefttext]
  [location=lefttext,width=3em,color=darkblue]
\definelocalboxes
  [lefttextx]
  [location=lefttext,width=3em,color=darkblue]

\definelocalboxes
  [righttext]
  [location=righttext,width=3em,color=darkyellow]
\definelocalboxes
  [righttextx]
  [location=righttext,width=3em,color=darkyellow]
```

The order of definition matters! Here the `x` variants have a larger index number. There can (currently) be at most 256 indices. The defined local boxes are triggered with `\localbox`:

```

\startnarrower
\dorecurse{20}{before #1 }%
\localbox[lefttext]{[L] }%
\localbox[lefttextx]{[LL] }%
\localbox[righttext]{ [RR]}%
\localbox[righttextx]{ [R]}%
\dorecurse{20}{ after #1}%
\stopnarrower

```

Watch how we obey the margins:

```

before 1 before 2 before 3 before 4 before 5 before 6 before 7 before 8 before 9
before 10 before 11 before 12 before 13 before 14 before 15 before 16 before 17
before 18 before 19 before 20 after 1 after 2 after 3 after 4 after 5 after [RR] [R]
[L] [LL] 6 after 7 after 8 after 9 after 10 after 11 after 12 after 13 after [RR] [R]
[L] [LL] 14 after 15 after 16 after 17 after 18 after 19 after 20 [RR] [R]

```

Here these local boxes have dimensions. The predefined margin variants are virtual. Here we set up the style and color:

```

\setuplocalboxes
[leftmargin]
[style=\bs,
color=darkgreen]
\setuplocalboxes
[rightmargin]
[style=\bs,
color=darkred]

\dorecurse{2}{
\dorecurse{10}{some text #1.##1 }%
KEY#1.1%
\localmargintext[leftmargin]{L #1.1}%
\localmargintext[rightmargin]{R #1.1}%
\dorecurse{10}{some text #1.##1 }%
KEY#1.2%
\localmargintext[leftmargin]{L #1.2}%
\localmargintext[rightmargin]{R #1.2}%
\dorecurse{10}{some text #1.##1 }%
\blank
}

```



You can also use `leftedge` and `rightedge` but using them here would put them outside the page.

**L 1.1** some text 1.1 some text 1.2 some text 1.3 some text 1.4 some text 1.5 some text 1.6  
 some text 1.7 some text 1.8 some text 1.9 some text 1.10 KEY1.1some text 1.1 some **R 1.1**  
 text 1.2 some text 1.3 some text 1.4 some text 1.5 some text 1.6 some text 1.7 some

**L 1.2** text 1.8 some text 1.9 some text 1.10 KEY1.2some text 1.1 some text 1.2 some text 1.3 **R 1.2**  
 some text 1.4 some text 1.5 some text 1.6 some text 1.7 some text 1.8 some text 1.9  
 some text 1.10

**L 2.1** some text 2.1 some text 2.2 some text 2.3 some text 2.4 some text 2.5 some text 2.6  
 some text 2.7 some text 2.8 some text 2.9 some text 2.10 KEY2.1some text 2.1 some **R 2.1**  
 text 2.2 some text 2.3 some text 2.4 some text 2.5 some text 2.6 some text 2.7 some

**L 2.2** text 2.8 some text 2.9 some text 2.10 KEY2.2some text 2.1 some text 2.2 some text 2.3 **R 2.2**  
 some text 2.4 some text 2.5 some text 2.6 some text 2.7 some text 2.8 some text 2.9  
 some text 2.10

In previous examples you can see that setting something at the left will lag behind so deep down we use another trick here: `\localmiddlebox`. When these boxes get placed a callback can be triggered and in `ConTeXt` we use that to move these middle boxes to the margins.

Next we implement line numbers. Watch out: this will not replace the existing mechanisms, it's just an alternative as we have alternative table mechanisms. We have a repertoire of helpers for constructing the result:

### `\definelocalboxes`

```
[linenumberleft]
[command=\LeftNumber,
  location=middle,
  distance=\leftmargindistance,
  width=3em,
  style=\bs,
  color=darkred]
```

### `\definelocalboxes`

```
[linenumberright] % [linenumberleft]
[command=\RightNumber,
  location=middle,
  distance=\rightmargindistance,
  width=3em,
  style=\bf,
```

```

color=darkgreen]

\definecounter[MyLineNumberL]
\definecounter[MyLineNumberR]

\setupcounter
  [MyLineNumberL]
  [numberconversion=characters]

\setupcounter
  [MyLineNumberR]
  [numberconversion=romannumerals]

\def\LineNumberL
  {\incrementcounter[MyLineNumberL]%
   \convertedcounter[MyLineNumberL]}

\def\LineNumberR
  {\incrementcounter[MyLineNumberR]%
   \convertedcounter[MyLineNumberR]}

\protected\def\LeftNumber
  {\setbox\localboxcontentbox\hbox
   to \localboxesparameter{width}
   {(\LineNumberL\hss\strut)}%
   \localmarginlefttext\zeropoint}

\protected\def\RightNumber
  {\setbox\localboxcontentbox\hbox
   to \localboxesparameter{width}
   {(\strut\hss\LineNumberR)}%
   \localmarginrighttext\zeropoint}

\localbox[linenumberleft]{}%
\localbox[linenumberright]{}%
\dorecurse{2}{
  \samplefile{tufte}
  \par
}
\resetlocalbox[linenumberleft]%
\resetlocalbox[linenumberright]%

```

We use our tufte example to illustrate the usage:

<b>(15.d)</b>	We thrive in information-thick worlds because of our marvelous and everyday capacity	<b>(15.i)</b>
<b>(15.h)</b>	to select, edit, single out, structure, highlight, group, pair, merge, harmonize, synthe-	<b>(15.ii)</b>
<b>(15.d)</b>	size, focus, organize, condense, reduce, boil down, choose, categorize, catalog, classify,	<b>(15.iii)</b>
<b>(15.d)</b>	list, abstract, scan, look into, idealize, isolate, discriminate, distinguish, screen, pigeon-	<b>(15.iv)</b>
<b>(15.e)</b>	hole, pick over, sort, integrate, blend, inspect, filter, lump, skip, smooth, chunk, aver-	<b>(15.v)</b>
<b>(15.f)</b>	age, approximate, cluster, aggregate, outline, summarize, itemize, review, dip into, flip	<b>(15.vi)</b>
<b>(15.g)</b>	through, browse, glance into, leaf through, skim, refine, enumerate, glean, synopsise,	<b>(16.vii)</b>
<b>(15.h)</b>	winnow the wheat from the chaff and separate the sheep from the goats.	<b>(15.viii)</b>
<b>(15.i)</b>	We thrive in information-thick worlds because of our marvelous and everyday capacity	<b>(15.ix)</b>
<b>(15.j)</b>	to select, edit, single out, structure, highlight, group, pair, merge, harmonize, synthe-	<b>(15.x)</b>
<b>(15.k)</b>	size, focus, organize, condense, reduce, boil down, choose, categorize, catalog, classify,	<b>(15.xi)</b>
<b>(15.l)</b>	list, abstract, scan, look into, idealize, isolate, discriminate, distinguish, screen, pigeon-	<b>(16.xii)</b>
<b>(15.m)</b>	hole, pick over, sort, integrate, blend, inspect, filter, lump, skip, smooth, chunk, aver-	<b>(15.xiii)</b>
<b>(15.n)</b>	age, approximate, cluster, aggregate, outline, summarize, itemize, review, dip into, flip	<b>(15.xiv)</b>
<b>(15.o)</b>	through, browse, glance into, leaf through, skim, refine, enumerate, glean, synopsise,	<b>(15.xv)</b>
<b>(15.p)</b>	winnow the wheat from the chaff and separate the sheep from the goats.	<b>(15.xvi)</b>

For convenience we support ranges like this (we've reset the line number counters here):

```

\startlocalboxrange[linenumberleft]%
\startlocalboxrange[linenumberright]%
\dorecurse{2}{
  \samplefile{tufte}
  \par
}
\stoplocalboxrange
\stoplocalboxrange

```

<b>(15.d)</b>	We thrive in information-thick worlds because of our marvelous and everyday capacity	<b>(15.i)</b>
<b>(15.h)</b>	to select, edit, single out, structure, highlight, group, pair, merge, harmonize, synthe-	<b>(15.ii)</b>
<b>(15.d)</b>	size, focus, organize, condense, reduce, boil down, choose, categorize, catalog, classify,	<b>(15.iii)</b>
<b>(15.d)</b>	list, abstract, scan, look into, idealize, isolate, discriminate, distinguish, screen, pigeon-	<b>(15.iv)</b>
<b>(15.e)</b>	hole, pick over, sort, integrate, blend, inspect, filter, lump, skip, smooth, chunk, aver-	<b>(15.v)</b>
<b>(15.f)</b>	age, approximate, cluster, aggregate, outline, summarize, itemize, review, dip into, flip	<b>(15.vi)</b>
<b>(15.g)</b>	through, browse, glance into, leaf through, skim, refine, enumerate, glean, synopsise,	<b>(16.vii)</b>
<b>(15.h)</b>	winnow the wheat from the chaff and separate the sheep from the goats.	<b>(15.viii)</b>
<b>(15.i)</b>	We thrive in information-thick worlds because of our marvelous and everyday capacity	<b>(15.ix)</b>
<b>(15.j)</b>	to select, edit, single out, structure, highlight, group, pair, merge, harmonize, synthe-	<b>(15.x)</b>
<b>(15.k)</b>	size, focus, organize, condense, reduce, boil down, choose, categorize, catalog, classify,	<b>(15.xi)</b>

- (15.l)** list, abstract, scan, look into, idealize, isolate, discriminate, distinguish, screen, pigeon- **16.xii)**  
**(15.m)** hole, pick over, sort, integrate, blend, inspect, filter, lump, skip, smooth, chunk, aver- **15.xiii)**  
**(15.n)** age, approximate, cluster, aggregate, outline, summarize, itemize, review, dip into, flip **15.xiv)**  
**(15.o)** through, browse, glance into, leaf through, skim, refine, enumerate, glean, synopsise, **15.xv)**  
**(15.p)** winnow the wheat from the chaff and separate the sheep from the goats. **15.xvi)**

## 15.4 The helpers

For the moment we have these helpers:

<code>\localboxindex</code>	integer
<code>\localboxlinenumber</code>	integer
	<code>\localboxlinewidth</code> dimension
<code>\localboxlocalwidth</code>	dimension
<code>\localboxprogress</code>	dimension
<code>\localboxleftoffset</code>	dimension
<code>\localboxrightoffset</code>	dimension
	<code>\localboxleftskip</code> dimension
<code>\localboxrightskip</code>	dimension
<code>\localboxlefthang</code>	dimension
<code>\localboxrighthang</code>	dimension
	<code>\localboxindent</code> dimension
<code>\localboxparfillleftskip</code>	dimension
<code>\localboxparfillrightskip</code>	dimension
<code>\localboxovershoot</code>	dimension

The progress and offsets are accumulated values of the normalized indent, hangs, skips etc. The line number is the position in the paragraph. In the callback we set the box register `\localboxcontentbox` and use it after the command has been applied. In the line number example you can see how we set its final content, so these boxes are sort of dynamic. Normally in the middle case no content is passed and in the par builder a middle is not taken into account when calculating the line width.

## 15.4 Colofon

Author	Hans Hagen
ConT <sub>E</sub> Xt	2025.07.04 21:26
LuaMetaT <sub>E</sub> X	2.11.07   20250705
Support	<a href="http://www.pragma-ade.com">www.pragma-ade.com</a> <a href="http://contextgarden.net">contextgarden.net</a> <a href="mailto:ntg-context@ntg.nl">ntg-context@ntg.nl</a>

## 16 Loops

low level

TEX

loops

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### 16.1 Introduction

I have hesitated long before I finally decided to implement native loops in LuaMetaTeX. Among the reasons against such a feature is that one can define macros that do loops (preferably using tail recursion). When you don't need an expandable loop, counters can be used, otherwise there are dirty and obscure tricks that can be of help. This is often the area where tex programmers can show off but the fact remains that we're using side effects of the expansion machinery and specific primitives like `\romannumeral` magic. In LuaMetaTeX it is actually possible to use the local control mechanism to hide loop counter advance and checking but that comes with at a performance hit. And, no matter what tricks one used, tracing becomes pretty much cluttered.

In the next sections we describe the new native loop primitives in LuaMetaTeX as well as the more traditional ConTeXt loop helpers.

### 16.2 Primitives

Because MetaPost, which is also a macro language, has native loops, it makes sense to also have native loops in TeX and in LuaMetaTeX it was not that hard to add it. One variant uses the local control mechanism which is reflected in its name and two others collect expanded bodies. In the local loop content gets injected as we go, so this one doesn't work well in for instance an `\edef`. The macro takes the usual three loop numbers as well as a token list:

```
\localcontrolledloop 1 100000 1 {%
  % body
}
```

Here is an example of usage:



```

\localcontrolledloop 1 5 1 {%
  [\number\currentloopiterator]
  \localcontrolledloop 1 10 1 {%
    (\number\currentloopiterator)
  }%
  [\number\currentloopiterator]
\par
}

```

The `\currentloopiterator` is a numeric token so you need to explicitly serialize it with `\number` or `\the` if you want it to be typeset:

```

[1] (1) (2) (3) (4) (5) (6) (7) (8) (9) (10) [1]
[2] (1) (2) (3) (4) (5) (6) (7) (8) (9) (10) [2]
[3] (1) (2) (3) (4) (5) (6) (7) (8) (9) (10) [3]
[4] (1) (2) (3) (4) (5) (6) (7) (8) (9) (10) [4]
[5] (1) (2) (3) (4) (5) (6) (7) (8) (9) (10) [5]

```

Here is another example. This time we also show the current nesting:

```

\localcontrolledloop 1 100 1 {%
  \ifnum\currentloopiterator>6\relax
    \quitloop
  \else
    [\number\currentloopnesting:\number\currentloopiterator]
    \localcontrolledloop 1 8 1 {%
      (\number\currentloopnesting:\number\currentloopiterator)
    }\par
  \fi
}

```

Watch the `\quitloop`: it will end the loop at the *next* iteration so any content after it will show up. Normally this one will be issued in a condition and we want to end that properly.

```

[1:1] (2:1) (2:2) (2:3) (2:4) (2:5) (2:6) (2:7) (2:8)
[1:2] (2:1) (2:2) (2:3) (2:4) (2:5) (2:6) (2:7) (2:8)
[1:3] (2:1) (2:2) (2:3) (2:4) (2:5) (2:6) (2:7) (2:8)
[1:4] (2:1) (2:2) (2:3) (2:4) (2:5) (2:6) (2:7) (2:8)
[1:5] (2:1) (2:2) (2:3) (2:4) (2:5) (2:6) (2:7) (2:8)
[1:6] (2:1) (2:2) (2:3) (2:4) (2:5) (2:6) (2:7) (2:8)

```

The three loop variants all perform differently:

```
l:\testfeatureonce {1000} {\localcontrolledloop 1 2000 1 {\relax}} %
  \elapsedtime
e:\testfeatureonce {1000} {\expandedloop          1 2000 1 {\relax}} %
  \elapsedtime
u:\testfeatureonce {1000} {\unexpandedloop        1 2000 1 {\relax}} %
  \elapsedtime
```

An unexpanded loop is (of course) the fastest because it only collects and then feeds back the lot. In an expanded loop each cycle does an expansion of the body and collects the result which is then injected afterwards, and the controlled loop just expands the body each iteration.

```
l: 0.093
e: 0.090
u: 0.031
```

The different behavior is best illustrated with the following example:

```
\edef\TestA{\localcontrolledloop 1 5 1 {A}} % out of order
\edef\TestB{\expandedloop          1 5 1 {B}}
\edef\TestC{\unexpandedloop        1 5 1 {C\relax}}
```

We can show the effective definition:

```
\meaningasis\TestA
\meaningasis\TestB
\meaningasis\TestC
```

```
A: \TestA
B: \TestB
C: \TestC
```

Watch how the first test pushes the content in the main input stream:

```
AAAAA
\def \TestA {}
\def \TestB {BBBBB}
\def \TestC {C\relax C\relax C\relax C\relax C\relax }

A:
B: BBBBB
```

C: CCCCC

Here are some examples that show what gets expanded and what not:

```
\edef\whatever
  {\expandedloop 1 10 1
   {(\number\currentloopiterator)
    \scratchcounter=\number\currentloopiterator\relax}}
```

```
\meaningasis\whatever
```

```
\def \whatever {(1) \scratchcounter =1\relax (2) \scratchcounter =2\relax (3)
\scratchcounter =3\relax (4) \scratchcounter =4\relax (5) \scratchcounter =5\relax (6)
\scratchcounter =6\relax (7) \scratchcounter =7\relax (8) \scratchcounter =8\relax (9)
\scratchcounter =9\relax (10) \scratchcounter =10\relax }
```

A local control encapsulation hides the assignment:

```
\edef\whatever
  {\expandedloop 1 10 1
   {(\number\currentloopiterator)
    \beginlocalcontrol
     \scratchcounter=\number\currentloopiterator\relax
    \endlocalcontrol}}
```

```
\meaningasis\whatever
```

```
\def \whatever {(1) (2) (3) (4) (5) (6) (7) (8) (9) (10) }
```

Here we see the assignment being retained but with changing values:

```
\edef\whatever
  {\unexpandedloop 1 10 1
   {\scratchcounter=1\relax}}
```

```
\meaningasis\whatever
```

```
\def \whatever {\scratchcounter =1\relax \scratchcounter =1\relax \scratchcounter =1\relax
\scratchcounter =1\relax \scratchcounter =1\relax \scratchcounter =1\relax \scratchcounter
=1\relax \scratchcounter =1\relax \scratchcounter =1\relax \scratchcounter =1\relax }
```

We get no expansion at all:

```
\edef\whatever
```

```
{\unexpandedloop 1 10 1
  {\scratchcounter=\the\currentloopiterator\relax}}
```

```
\meaningasis\whatever
```

```
\def \whatever {\scratchcounter =0\relax \scratchcounter =0\relax \scratchcounter =0\relax
\scratchcounter =0\relax \scratchcounter =0\relax \scratchcounter =0\relax \scratchcounter
=0\relax \scratchcounter =0\relax \scratchcounter =0\relax \scratchcounter =0\relax }
```

And here we have a mix:

```
\edef\whatever
  {\expandedloop 1 10 1
   {\scratchcounter=\the\currentloopiterator\relax}}
```

```
\meaningasis\whatever
```

```
\def \whatever {\scratchcounter =1\relax \scratchcounter =2\relax \scratchcounter =3\relax
\scratchcounter =4\relax \scratchcounter =5\relax \scratchcounter =6\relax \scratchcounter
=7\relax \scratchcounter =8\relax \scratchcounter =9\relax \scratchcounter =10\relax }
```

There is one feature worth noting. When you feed three numbers in a row, like here, there is a danger of them being seen as one:

```
\expandedloop
  \number\dimexpr1pt
  \number\dimexpr2pt
  \number\dimexpr1pt
  {}
```

This gives an error because a too large number is seen. Therefore, these loops permit leading equal signs, as in assignments (we could support keywords but it doesn't make much sense):

```
\expandedloop =\number\dimexpr1pt =\number\dimexpr2pt =\number\dimexpr1pt{}
```

## 16.3 Wrappers

We always had loop helpers in ConTEXt and the question is: “What we will gain when we replace the definitions with ones using the above?”. The answer is: “We have little performance but not as much as one expects!”. This has to do with the fact that we support #1 as iterator and #2 as (verbose) nesting values and that comes with some

overhead. It is also the reason why these loop macros are protected (unexpandable). However, using the primitives might look somewhat more natural in low level T<sub>E</sub>X code.

Also, replacing their definitions can have side effects because the primitives are (and will be) still experimental so it's typically a patch that I will run on my machine for a while.

Here is an example of two loops. The inner state variables have one hash, the outer one extra:

```
\dorecurse{2}{
  \dostepwiserecurse{1}{10}{2}{
    (#1:#2) [##1:##2]
  }\par
}
```

We get this:

```
(1:1) [1:2] (1:1) [3:2] (1:1) [5:2] (1:1) [7:2] (1:1) [9:2]
(2:1) [1:2] (2:1) [3:2] (2:1) [5:2] (2:1) [7:2] (2:1) [9:2]
```

We can also use two state macro but here we would have to store the outer ones:

```
\dorecurse {2} {
  /\recursedepth:\recurselevel/
  \dostepwiserecurse {1} {10} {2} {
    <\recursedepth:\recurselevel>
  }\par
}
```

That gives us:

```
/1:1/ <2:1> <2:3> <2:5> <2:7> <2:9>
/1:2/ <2:1> <2:3> <2:5> <2:7> <2:9>
```

An endless loop works as follows:

```
\doloop {
  ...
  \ifsomeconditionismet
  ...
  \exitloop
\else
```

```

    ...
    \fi
    % \exitloopnow
    ...
}

```

Because of the way we quit there will not be a new implementation in terms of the loop primitives. You need to make sure that you don't leave in the middle of an ongoing condition. The second exit is immediate.

We also have a (simple) expanded variant:

```
\edef\TestX{\doexpandedrecurse{10}{!}} \meaningasis\TestX
```

This helper can be implemented in terms of the loop primitives which makes them a bit faster, but these are not critical:

```
\def\TestX {!!!!!!!!!!!!}
```

A variant that supports #1 is the following:

```
\edef\TestX{\doexpandedrecursed{10}{#1}} \meaningasis\TestX
```

So:

```
\def\TestX {12345678910}
```

## 16.4 About quitting

You can quit a local and expanded loop at the next iteration using `\quitloop`. With `\quitloopnow` you immediately leave the loop but you need to beware of side effects, like not ending a condition properly. Keep in mind that a macro language like  $\text{\TeX}$  is not that friendly towards loops so the implementation is a bit hairy.

## 16.5 Simple repeaters

For simple iterations we have `\localcontrolledrepeat`, `\expandedrepeat`, `\unexpandedrepeat`. These take one integer instead of three: the final iterator value.

## 16.6 Endless loops

There are three endless loop primitives: `\localcontrolledendless`, `\expandedendless`, `\unexpandedendless`. These will keep running till you quit them. The loop counter can overflow the maximum integer value and will then start again at 1.

## 16.7 Loop variables

The following example shows how we can access the current, parent and grand parent loop iterator values using a parameter like syntax:

```
\localcontrolledloop 1 4 1 {%
  \localcontrolledloop 1 3 1 {%
    \localcontrolledloop 1 2 1 {%
      \edef\foo{[#G,#P,#I]}\foo
      \def \oof{<#G,#P,#I>}\oof
        (#G,#P,#I)\space
      }
    \par
  }
}
```

```
[1,1,1]<1,1,1>(1,1,1) [1,1,2]<1,1,2>(1,1,2)
[1,2,1]<1,2,1>(1,2,1) [1,2,2]<1,2,2>(1,2,2)
[1,3,1]<1,3,1>(1,3,1) [1,3,2]<1,3,2>(1,3,2)
[2,1,1]<2,1,1>(2,1,1) [2,1,2]<2,1,2>(2,1,2)
[2,2,1]<2,2,1>(2,2,1) [2,2,2]<2,2,2>(2,2,2)
[2,3,1]<2,3,1>(2,3,1) [2,3,2]<2,3,2>(2,3,2)
[3,1,1]<3,1,1>(3,1,1) [3,1,2]<3,1,2>(3,1,2)
[3,2,1]<3,2,1>(3,2,1) [3,2,2]<3,2,2>(3,2,2)
[3,3,1]<3,3,1>(3,3,1) [3,3,2]<3,3,2>(3,3,2)
[4,1,1]<4,1,1>(4,1,1) [4,1,2]<4,1,2>(4,1,2)
[4,2,1]<4,2,1>(4,2,1) [4,2,2]<4,2,2>(4,2,2)
[4,3,1]<4,3,1>(4,3,1) [4,3,2]<4,3,2>(4,3,2)
```

Another way to access a(ny) parent is:

```
\localcontrolledloop 1 4 1 {%
  \localcontrolledloop 1 3 1 {%
    \localcontrolledloop 1 2 1 {%
      (\the\previousloopiterator2,%
      \the\previousloopiterator1,%
      \the\currentloopiterator)
    }
  \par
}
}
```

These methods make that one doesn't have to store the outer loop variables for usage inside the inner loop. Watch out with the `\edef`:

```
\edef\foo{[#G,#P,#I]}
\def \oof{<#G,#P,#I>}

\localcontrolledloop 1 4 1 {%
  \localcontrolledloop 1 3 1 {%
    \localcontrolledloop 1 2 1 {%
      %
      % I iterator    \currentloopiterator
      % P parent      \previousloopiterator1
      % G grandparent \previousloopiterator2
      %
      \edef\ofo{[#G,#P,#I]}%
      \foo\oof\ofo(#G,#P,#I)\space
    }
  }
}
\par
}
```

```
[0,0,0]<1,1,1>[1,1,1](1,1,1) [0,0,0]<1,1,2>[1,1,2](1,1,2)
[0,0,0]<1,2,1>[1,2,1](1,2,1) [0,0,0]<1,2,2>[1,2,2](1,2,2)
[0,0,0]<1,3,1>[1,3,1](1,3,1) [0,0,0]<1,3,2>[1,3,2](1,3,2)
[0,0,0]<2,1,1>[2,1,1](2,1,1) [0,0,0]<2,1,2>[2,1,2](2,1,2)
[0,0,0]<2,2,1>[2,2,1](2,2,1) [0,0,0]<2,2,2>[2,2,2](2,2,2)
[0,0,0]<2,3,1>[2,3,1](2,3,1) [0,0,0]<2,3,2>[2,3,2](2,3,2)
[0,0,0]<3,1,1>[3,1,1](3,1,1) [0,0,0]<3,1,2>[3,1,2](3,1,2)
[0,0,0]<3,2,1>[3,2,1](3,2,1) [0,0,0]<3,2,2>[3,2,2](3,2,2)
[0,0,0]<3,3,1>[3,3,1](3,3,1) [0,0,0]<3,3,2>[3,3,2](3,3,2)
[0,0,0]<4,1,1>[4,1,1](4,1,1) [0,0,0]<4,1,2>[4,1,2](4,1,2)
[0,0,0]<4,2,1>[4,2,1](4,2,1) [0,0,0]<4,2,2>[4,2,2](4,2,2)
[0,0,0]<4,3,1>[4,3,1](4,3,1) [0,0,0]<4,3,2>[4,3,2](4,3,2)
```



## 16.7 Colofon

Author           Hans Hagen  
ConT<sub>E</sub>Xt        2025.07.04 21:26  
LuaMetaT<sub>E</sub>X    2.11.07 | 20250705  
Support         www.pragma-ade.com  
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## **17 Tokens**

low level

TEX

tokens

## Contents

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### 17.1 Introduction

Most users don't need to know anything about tokens but it happens that when T<sub>E</sub>Xies meet in person (users group meetings), or online (support platforms) there always seem to pop up folks who love token speak. When you try to explain something to a user it makes sense to talk in terms of characters but then those token speakers can jump in and start correcting you. In the past I have been puzzled by this because, when one can write a decent macro that does the job well, it really doesn't matter if one knows about tokens. Of course one should never make the assumption that token speakers really know T<sub>E</sub>X that well or can come up with better solutions than users but that is another matter.<sup>23</sup>

That said, because in documents about T<sub>E</sub>X the word 'token' does pop up I will try to give a little insight here. But for using T<sub>E</sub>X it's mostly irrelevant. The descriptions below for sure won't match the proper token speak criteria which is why at a presentation for the 2020 user meeting I used the title "Tokens as I see them."

### 17.2 What are tokens

Both the words 'node' and 'token' are quite common in programming and also rather old which is proven by the fact that they also are used in the T<sub>E</sub>X source. A node is a storage container that is part of a linked list. When you input the characters tex the three characters become part of the current linked list. They become 'character' nodes (or in LuaT<sub>E</sub>X speak 'glyph' nodes) with properties like the font and the character referred to. But before that happens, the three characters in the input t, e and x, are interpreted as in this case being just that: characters. When you enter \TeX the input processors first sees a backslash and because that has a special meaning in T<sub>E</sub>X it will

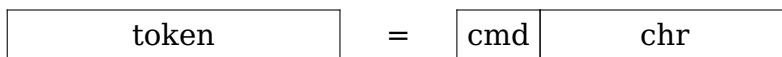
---

<sup>23</sup> Talking about fashion: it would be more impressive to talk about T<sub>E</sub>X and friends as a software stack than calling it a distribution. Today, it's all about marketing.

read following characters and when done does a lookup in it's internal hash table to see what it actually is: a macro that assembled the word  $\TeX$  in uppercase with special kerning and a shifted (therefore boxed) 'E'. When you enter  $\$ \TeX$  will look ahead for a second one in order to determine display math, push back the found token when there is no match and then enter inline math mode.

A token is internally just a 32 bit number that encodes what  $\TeX$  has seen. It is the assembled token that travels through the system, get stored, interpreted and often discarded afterwards. So, the character 'e' in our example gets tagged as such and encoded in this number in a way that the intention can be derived later on.

Now, the way  $\TeX$  looks at these tokens can differ. In some cases it will just look at this (32 bit) number, for instance when checking for a specific token, which is fast, but sometimes it needs to know some detail. The mentioned integer actually encodes a command (opcode) and a so called char code (operand). The second name is somewhat confusing because in many cases that code is not representing a character but that is not that relevant here. When you look at the source code of a  $\TeX$  engine it is enough to know that a char can also be a sub command.



Back to the three characters: these become tokens where the command code indicates that it is a letter and the char code stores what letter we have at hand and in the case of  $\text{Lua}\TeX$  and  $\text{LuaMeta}\TeX$  these are Unicode values. Contrary to the traditional 8 bit  $\TeX$  engine, in the Unicode engines an utf sequence is read, but these multiple bytes still become one number that will be encoded in the token number. In order to determine that something is a letter the engine has to be told (which is what a macro package does when it sets up the engine). For instance, digits are so called other characters and the backslash is called escape. Every  $\TeX$  user knows that curly braces are special and so are dollar symbols and hashes. If this rings a bell, and you relate this to catcodes, you can indeed assume that the command codes of these tokens have the same numbers as the catcodes. Given that Unicode has plenty of characters slots you can imagine that combining 16 catcode commands with all the possible Unicode values makes a large repertoire of tokens.

There are more commands than the 16 basic characters related ones, in  $\text{LuaMeta}\TeX$  we have just over 150 command codes ( $\text{Lua}\TeX$  has a few more but they are also organized differently). Each of these codes can have a sub command, For instance the primitives  $\backslash\text{vbox}$  and  $\backslash\text{hbox}$  are both a `make_box_cmd` (we use the symbolic name here) and in  $\text{LuaMeta}\TeX$  the first one has sub command code 9 (`vbox_code`) and the second one has code 10 (`hbox_code`). There are twelve primitives that are in the same category.

### What are tokens

The many primitives that make up the core of the engine are grouped in a way that permits processing similar ones with one function and also makes it possible to distinguish between the way commands are handled, for instance with respect to expansion.

Now, before we move on it is important to know that all these codes are in fact abstract numbers. Although it is quite likely that engines that are derived from each other have similar numbers (just more) this is not the case for LuaMetaT<sub>E</sub>X. Because the internals have been opened up (even more than in LuaT<sub>E</sub>X) the command and char codes have been reorganized in a such a way that exposure is consistent. We could not use some of the reuse and remap tricks that the other engines use because it would simply be too confusing (and demand real in depth knowledge of the internals). This is also the reason why development took some time. You probably won't notice it from the current source but it was a very stepwise process. We not only had to make sure that all kept working (ConT<sub>E</sub>Xt LMTX and LuaMetaT<sub>E</sub>X were pretty useable during the process), but also had to (re)consider intermediate choices.

So, input is converted into tokens, in most cases one-by-one. When a token is assembled, it either gets stored (deliberately or as part of some look ahead scanning), or it immediately gets (what is called:) expanded. Depending on what the command is, some action is triggered. For instance, a character gets appended to the node list immediately. An `\hbox` command will start assembling a box which its own node list that then gets some treatment: if this primitive was a follow up on `\setbox` it will get stored, otherwise it might end up in the current node list as so called hlist node. Commands that relate to registers have `0xFFFF` char codes because that is how many registers we have per category.

When a token gets stored for later processing it becomes part of a larger data structure, a so called memory word. These memory words are taken from a large pool of words and they store a token and additional properties. The info field contains the token value, the mentioned command and char. When there is no linked list, the link can actually be used to store a value, something that in LuaMetaT<sub>E</sub>X we actually do.

1	info	link
2	info	link
3	info	link
n	info	link

When for instance we say `\toks 0 {tex}` the scanner sees an escape, followed by 4 letters (toks) and the escape triggers a lookup of the primitive (or macro or ...) with that name, in this case a primitive assignment command. The found primitive (its property gets stored in the token) triggers scanning for a number and when that is successful

scanning of a brace delimited token list starts. The three characters become three letter tokens and these are a linked list of the mentioned memory words. This list then gets stored in token register zero. The input sequence `\the\toks 0` will push back a copy of this list into the input.

In addition to the token memory pool, there is also a table of equivalents. That one is part of a larger table of memory words where  $\text{\TeX}$  stores all it needs to store. The 16 groups of character commands are virtual, storing these makes no sense, so the first real entries are all these registers (count, dimension, skip, box, etc). The rest is taken up by possible hash entries.

main hash	null control sequence
	128K hash entries
	frozen control sequences
	special sequences (undefined)
registers	17 internal & 64K user glues
	4 internal & 64K user mu glues
	12 internal & 64K user tokens
	2 internal & 64K user boxes
	116 internal & 64K user integers
	0 internal & 64K user attribute
	22 internal & 64K user dimensions
specifications	5 internal & 0 user
extra hash	additional entries (grows dynamic)

So, a letter token `t` is just that, a token. A token referring to a register is again just a number, but its char code points to a slot in the equivalents table. A macro, which we haven't discussed yet, is actually just a token list. When a name lookup happens the hash table is consulted and that one runs in parallel to part of the table of equivalents. When there is a match, the corresponding entry in the equivalents table points to a token list.

1	string index	equivalents or (next > n) index
2	string index	equivalents or (next > n) index
n	string index	equivalents or (next > n) index
n + 1	string index	equivalents or (next > n) index
n + 2	string index	equivalents or (next > n) index
n + m	string index	equivalents or (next > n) index

It sounds complex and it actually also is somewhat complex. It is not made easier by the fact that we also track information related to grouping (saving and restoring), need reference counts for copies of macros and token lists, sometimes store information directly instead of via links to token lists, etc. And again one cannot compare LuaMetaTeX with the other engines. Because we did away with some of the limitations of the traditional engine we not only could save some memory but in the end also simplify matters (we're 32/64 bit after all). On the one hand some traditional speedups were removed but these have been compensated by improvements elsewhere, so overall processing is more efficient.

1	level	type	flag	value
2	level	type	flag	value
3	level	type	flag	value
n	level	type	flag	value

So, here LuaMetaTeX differs from other engines because it combines two tables, which is possible because we have at least 32 bits. There are at most 0xFFFF levels but we need at most 0xFF types. In LuaMetaTeX macros can have extra properties (flags) and these also need one byte. Contrary to the other engines, `\protected` macros are native and have their own command code, but `\tolerant` macros duplicate that (so we have four distinct macro commands). All other properties, like the `\permanent` ones are stored in the flags.

Because a macro starts with a reference count we have some room in the info field to store information about it having arguments or not. It is these details that make LuaMetaTeX a bit more efficient in terms of memory usage and performance than its ancestor LuaTeX. But as with the other changes, it was a very stepwise process in order to keep the system compatible and working.

### 17.3 Some implementation details

Sometimes there is a special head token at the start. This makes for easier appending of extra tokens. In traditional TeX node lists are forward linked, in LuaTeX they are



double linked<sup>24</sup>. Token lists are always forward linked. Shared token lists use the head node for a reference count.

For various reasons original  $\text{T}_{\text{E}}\text{X}$  uses global variables temporary lists. This is for instance needed when we expand (nested) and need to report issues. But in  $\text{LuaT}_{\text{E}}\text{X}$  we often just serialize lists and using local variables makes more sense. One of the first things done in  $\text{LuaMetaT}_{\text{E}}\text{X}$  was to group all global variables in (still global) structures but well isolated. That also made it possible to actually get rid of some globals.

Because  $\text{T}_{\text{E}}\text{X}$  had to run on machines that we nowadays consider rather limited, it had to be sparse and efficient. There are quite some optimizations to limit code and memory consumption. The engine also does its own memory management. Freed token memory words are collected in a cache and reused but they can get scattered which is not that bad, apart from maybe cache hits. In  $\text{LuaMetaT}_{\text{E}}\text{X}$  we stay as close to original  $\text{T}_{\text{E}}\text{X}$  as possible but there have been some improvements. The Lua interfaces force us to occasionally divert from the original, and that in fact might lead to some retrofit but the original documentation still mostly applies. However, keep in mind that in  $\text{LuaT}_{\text{E}}\text{X}$  we store much more in nodes (each has a prev pointer and an attribute list pointer and for instance glyph nodes have some 20 extra fields compared to traditional  $\text{T}_{\text{E}}\text{X}$  character nodes).

## 17.4 Other data management

There is plenty going on in  $\text{T}_{\text{E}}\text{X}$  when it processes your input, just to mention a few:

- Grouping is handled by a nesting stack.
- Nested conditionals (`\if . . .`) have their own stack.
- The values before assignments are saved on the save stack.
- Also other local changes (housekeeping) ends up in the save stack.
- Token lists and macro aliases have references pointers (reuse).
- Attributes, being linked node lists, have their own management.

In all these subsystems tokens or references to tokens can play a role. Reading a single character from the input can trigger a lot of action. A curly brace tagged as begin group command will push the grouping level and from then on registers and some other quantities that are changed will be stored on the save stack so that after the group ends they can be restored. When primitives take keywords, and no match happens, tokens are pushed back into the input which introduces a new input level (also some stack).

---

<sup>24</sup> On the agenda of  $\text{LuaMetaT}_{\text{E}}\text{X}$  is to use this property in the underlying code, that doesn't yet profit from this and therefore keep previous pointers in store.

When numbers are read a token that represents no digit is pushed back too and macro packages use numbers and dimensions a lot. It is a surprise that T<sub>E</sub>X is so fast.

## 17.5 Macros

There is a distinction between primitives, the build in commands, and macros, the commands defined by users. A primitive relates to a command code and char code but macros are, unless they are made an alias to something else, like a `\countdef` or `\let` does, basically pointers to a token list. There is some additional data stored that makes it possible to parse and grab arguments.

When we have a control sequence (macro) `\controlsequence` the name is looked up in the hash table. When found its value will point to the table of equivalents. As mentioned, that table keeps track of the cmd and points to a token list (the meaning). We saw that this table also stores the current level of grouping and flags.

If we say, in the input, `\hbox to 10pt {x\hss}`, the box is assembled as we go and when it is appended to the current node list there are no tokens left. When scanning this, the engine literally sees a backslash and the four letters `hbox`. However when we have this:

```
\def\MyMacro{\hbox to 10pt {x\hss}}
```

the `\hbox` has become one memory word which has a token representing the `\hbox` primitive plus a link to the next token. The space after a control sequence is gobbled so the next two tokens, again stored in a linked memory word, are letter tokens, followed by two other and two letter tokens for the dimensions. Then we have a space, a brace, a letter, a primitive and a brace. The about 20 characters in the input became a dozen memory words each two times four bytes, so in terms of memory usage we end up with quite a bit more. However, when T<sub>E</sub>X runs over that list it only has to interpret the token values because the scanning and conversion already happened. So, the space that a macro takes is more than compensated by efficient reprocessing.

## 17.6 Looking at tokens

When you say `\tracingall` you will see what the engine does: read input, expand primitives and macros, typesetting etc. You might need to set `\tracingonline` to get a bit more output on the console. One way to look at macros is to use the `\meaning` command, so if we have:

```
\permanent\protected\def\MyMacro#1#2{Do #1 or #2!}
```

we can say this:

```
\meaning    \MyMacro
\meaningless\MyMacro
\meaningfull\MyMacro
```

and get:

```
protected macro:#1#2->Do #1 or #2!
#1#2->Do #1 or #2!
permanent protected macro:#1#2->Do #1 or #2!
```

You get less when you ask for the meaning of a primitive, just its name. The `\meaningfull` primitive gives the most information. In LuaMetaTeX protected macros are first class commands: they have their own command code. In the other engines they are just regular macros with an initial token indicating that they are protected. There are specific command codes for `\outer` and `\long` macros but we dropped that in LuaMetaTeX. Instead we have `\tolerant` macros but that's another story. The flags that were mentioned can mark macros in a way that permits overload protection as well as some special treatment in otherwise tricky cases (like alignments). The overload related flags permits a rather granular way to prevent users from redefining macros and such. They are set via prefixes, and add to that repertoire: we have 14 prefixes but only some eight deal with flags (we can add more if really needed). The probably most well known prefix is `\global` and that one will not become a flag: it has immediate effect.

For the above definition, the `\showluatokens` command will show a meaning on the console.

```
\showluatokens\MyMacro
```

This gives the next list, where the first column is the address of the token, the second one the command code, and the third one the char code. When there are arguments involved, the list of what needs to get matched is shown.

```
permanent protected control sequence: MyMacro
501263  19   49  match                argument 1
501087  19   50  match                argument 2
385528  20    0  end match
-----
501090  11   68  letter              D (U+00044)
 30833  11  111  letter              o (U+0006F)
500776  10   32  spacer
385540  21    1  parameter reference
```

112057	10	32	spacer	
431886	11	111	letter	o (U+0006F)
30830	11	114	letter	r (U+00072)
30805	10	32	spacer	
500787	21	2	parameter reference	
213412	12	33	other char	! (U+00021)

In the next subsections I will give some examples. This time we use helper defined in a module:

```
\usemodule[system-tokens]
```

### 17.6.1 Example 1: in the input

```
\luatokenable{1 \bf{2} 3\what {!}}
```

---

given token list:

<no tokens>

---

### 17.6.2 Example 2: in the input

```
\luatokenable{a \the\scratchcounter b \the\parindent \hbox to 10pt{x}}
```

---

given token list:

<no tokens>

---

### 17.6.3 Example 3: user registers

```
\scratchtoks{foo \framed{\red 123}456}
```

```
\luatokenable\scratchtoks
```

---

token register: scratchtoks

<no tokens>

---

### 17.6.4 Example 4: internal variables

```
\luatokenable\everypar
```

---

internal token variable: everypar

---

---

<no tokens>

---

### 17.6.5 Example 5: macro definitions

```
\protected\def\whatever#1[#2](#3)\relax
  {oops #1 and #2 & #3 done ## error}
```

```
\luatokenable\whatever
```

---

**protected control sequence: whatever**

---

601588	19	49	match		argument 1
598513	12	91	other char	[ U+0005B	
599582	19	50	match		argument 2
597688	12	93	other char	] U+0005D	
594830	12	40	other char	( U+00028	
596423	19	51	match		argument 3
597514	12	41	other char	) U+00029	
598713	16	0	relax		relax
600456	20	0	end match		

---

600816	11	111	letter	o U+0006F	
598702	11	101	letter	e U+00065	
599302	11	112	letter	p U+00070	
597510	11	115	letter	s U+00073	
597084	10	32	spacer		
599022	21	1	parameter reference		
598724	10	32	spacer		
597011	11	97	letter	a U+00061	
599292	11	110	letter	n U+0006E	
596652	11	100	letter	d U+00064	
16596	10	32	spacer		
599181	21	2	parameter reference		
597874	10	32	spacer		
599856	12	38	other char	& U+00026	
598410	10	32	spacer		
598481	21	3	parameter reference		
600361	10	32	spacer		
599147	11	100	letter	d U+00064	
599693	11	111	letter	o U+0006F	
599414	11	110	letter	n U+0006E	
597061	11	101	letter	e U+00065	
599672	10	32	spacer		
597651	6	35	parameter		
598035	10	32	spacer		
600537	11	101	letter	e U+00065	
596792	11	114	letter	r U+00072	
598143	11	114	letter	r U+00072	
596924	11	111	letter	o U+0006F	
595380	11	114	letter	r U+00072	

---

## 17.6.6 Example 6: commands

`\luatokenable\startitemize`

`\luatokenable\stopitemize`

---

**frozen instance protected control sequence: startitemize**

---

523223	147	0	tolerant protected call		startitemgroup
523224	12	91	other char	[	U+0005B
523225	11	105	letter	i	U+00069
523226	11	116	letter	t	U+00074
523227	11	101	letter	e	U+00065
523228	11	109	letter	m	U+0006D
523229	11	105	letter	i	U+00069
523230	11	122	letter	z	U+0007A
523231	11	101	letter	e	U+00065
523232	12	93	other char	]	U+0005D

---

**frozen instance protected control sequence: stopitemize**

433433	143	0	protected call		stopitemgroup
--------	-----	---	----------------	--	---------------

---

## 17.6.7 Example 7: commands

`\luatokenable\doifelse`

---

**permanent protected control sequence: doifelse**

---

55521	19	49	match		argument 1
55522	19	50	match		argument 2
55523	20	0	end match		
55524	137	29	if test		iftok
55525	1	123	left brace		
55526	21	1	parameter reference		
55527	2	125	right brace		
55528	1	123	left brace		
55529	21	2	parameter reference		
55530	2	125	right brace		
55531	131	0	expand after		expandafter
55532	142	0	call		firstoftwoarguments
55533	137	3	if test		else
55534	131	0	expand after		expandafter
55535	142	0	call		secondoftwoarguments
55536	137	2	if test		fi

---

## 17.6.8 Example 8: nothing

`\luatokenable\relax`

---

primitive control sequence: relax

---

---

<no tokens>

---

### 17.6.9 Example 9: hashes

```
\edef\foo#1#2{(#1)(\letterhash)(#2)} \luatokenable\foo
```

---

**control sequence: foo**


---

599292	19	49	match		argument 1
597504	19	50	match		argument 2
593449	20	0	end match		
598535	12	40	other char	(	U+00028
596956	21	1	parameter reference		
481196	12	41	other char	)	U+00029
597423	12	40	other char	(	U+00028
600676	12	35	other char	#	U+00023
600779	12	41	other char	)	U+00029
599449	12	40	other char	(	U+00028
596652	21	2	parameter reference		
596618	12	41	other char	)	U+00029

---

### 17.6.10 Example 10: nesting

```
\def\foo#1{\def\foo##1{(#1)(##1)}} \luatokenable\foo
```

---

**control sequence: foo**


---

601207	19	49	match		argument 1
599510	20	0	end match		
598513	128	1	def		def
595530	142	0	call		foo
599556	6	35	parameter		
596626	12	49	other char	1	U+00031
598394	1	123	left brace		
601259	12	40	other char	(	U+00028
596833	21	1	parameter reference		
597354	12	41	other char	)	U+00029
599542	12	40	other char	(	U+00028
595380	6	35	parameter		
600292	12	49	other char	1	U+00031
600409	12	41	other char	)	U+00029
598136	2	125	right brace		

---

### 17.6.11 Remark

In all these examples the numbers are to be seen as abstractions. Some command codes and sub command codes might change as the engine evolves. This is why the Lua-MetaTeX engine has lots of Lua functions that provide information about what number represents what command.

## 17.6.11 Colofon

Author	Hans Hagen
ConT <sub>E</sub> Xt	2025.07.04 21:26
LuaMetaT <sub>E</sub> X	2.11.07   20250705
Support	<a href="http://www.pragma-ade.com">www.pragma-ade.com</a> <a href="http://contextgarden.net">contextgarden.net</a> <a href="mailto:ntg-context@ntg.nl">ntg-context@ntg.nl</a>



## 18 Buffers

low level

TEX

buffers

## Contents

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### 18.1 Preamble

Buffers are not that low level but it makes sense to discuss them in this perspective because it relates to tokenization, internal representation and manipulating.

*In due time we can describe some more commands and details here. This is a start. Feel free to tell me what needs to be explained.*

### 18.2 Encoding

Normally processing a document starts with reading from file. In the past we were talking single bytes that were then maps onto a specific input encoding that itself matches the encoding of a font. When you enter an 'a' its (normally ascii) number 97 becomes the index into a font. That same number is also used in the hyphenator which is why font encoding and hyphenation are strongly related. If in an eight bit T<sub>E</sub>X engine you need a precomposed 'ä' you have to use an encoding that has that character in some slot with again matching fonts and patterns. The actually used font can have the *shapes* in different slots and remapping is then done in the backend code using encoding and mapping files. When OpenType fonts are used the relationship between characters (input) and glyphs (rendering) also depends on the application of font features.

In eight bit environments all this brings a bit of a resource management nightmare along with complex installation of new fonts. It also puts strain on the macro package, especially when you want to mix different input encodings onto different font encodings and thereby pattern encodings in the same document. You can compare this with code pages in operating system, but imagine them potentially being mixed in one document,

which can happen when you mix multiple languages where the accumulated number of different characters exceeds 256. You end up switching between encodings. One way to deal with it is making special characters active and let their meaning differ per situation. That is for instance how in MkII we handled utf8 and thereby got around distributing multiple pattern files per language as we only needed to encoding them in utf and then remap them to the required encoding when loading patterns. A mental exercise is wondering how to support cjk scripts in an eight bit MkII, something that actually can be done with some effort.

The good news is that when we moved from MkII to MkIV we went exclusively utf8 because that is what the Lua $\TeX$  engine expects. Upto four bytes are read in and translated into one Unicode character. The internal representation is a 32 bit integer (four bytes) instead of a single byte. That also means that in the transition we got rid of quite some encoding related low level font and pattern handling. We still support input encodings (called regimes in Con $\TeX$ t) but I'm pretty sure that nowadays no one uses input other than utf8. While Con $\TeX$ t is normally quite upward compatible this is one area where there were fundamental changes.

There is still some interpretation going on when reading from file: for instance, we need to normalize the Unicode input, and we feed the engine separate lines on demand. Apart from that, some characters like the backslash, dollar sign and curly braces have special meaning so for accessing them as characters we have to use commands that inject those characters. That didn't change when we went from MkII to MkIV. In practice it's never really a problem unless you find yourself in one of the following situations:

- *Example code has to be typeset as-is, so braces etc. are just that.* This means that we have to change the way characters are interpreted. Typesetting code is needed when you want to document  $\TeX$  and macros which is why mechanisms for that have to be present right from the start.
- *Content is collected and used later.* A separation of content and usage later on often helps making a source look cleaner. Examples are “wrapping a table in a buffer” and “including that buffer when a table is placed” using the placement macros.
- *Embedded MetaPost and Lua code.* These languages come with different interpretation of some characters and especially MetaPost code is often stored first and used (processed) later.

- *The content comes from a different source.* Examples are xml files where angle brackets are special but for instance braces aren't. The data is interpreted as a stream or as a structured tree.
- *The content is generated.* It can for instance come from Lua, where bytes (representing utf) is just text and no special characters are to be intercepted. Or it can come from a database (using a library).

For these reasons ConT<sub>E</sub>Xt always had ways to store data in ways that makes this possible. The details on how that is done might have changed over versions, been optimized, extended with additional interfaces and features but given where we come from most has been there from the start.

## 18.3 Performance

When T<sub>E</sub>X came around, the bottlenecks in running T<sub>E</sub>X were the processor, memory and disks and depending on the way one used it the speed of the console or terminal; so, basically the whole system. One could sit there and wait for the page counters ([1] [2] . . . to show up. It was possible to run T<sub>E</sub>X on a personal computer but it was somewhat resource hungry: one needed a decent disk (a 10 MB hard disk was huge and with today's phone camera snapshots that sounds crazy). One could use memory extenders to get around the 640K limitation (keep in mind that the programs and operating systems also took space). This all meant that one could not afford to store too many tokens in memory but even using files for all kind of (multi-pass) trickery was demanding.

When processors became faster and memory plenty the disk became the bottleneck, but that changed when ssd's showed up. Combined with already present file caching that had some impact. We are now in a situation that cpu cores don't get that much faster (at least not twice as fast per iteration) and with T<sub>E</sub>X being a single core byte cruncher we're more or less in a situation where performance has to come from efficient programming. That means that, given enough memory, in some cases storing in tokens wins over storing in files, but it is no rule. In practice there is not much difference so one can even more than yesterday choose for the most convenient method. Just assume that the ConT<sub>E</sub>Xt code, combined with LuaMetaT<sub>E</sub>X will give you what you need with a reasonable performance. When in doubt, test with simple test files and it that works out well compared to the real code, try to figure out where 'mistakes' are made. Inefficient Lua and T<sub>E</sub>X code has way more impact than storing a few more tokens or using some files.

## 18.4 Files

Nearly always files are read once per run. The content (mixed with commands) is scanned and macros are expanded and/or text is typeset as we go. Internally the LuaMetaTeX engine is in “scanning from file”, “scanning from token lists”, or “scanning from Lua output” mode. The first mode is (in principle) the slowest because utf sequences are converted to tokens (numbers) but there is no way around it. The second method is fast because we already have these numbers, but we need to take into account where the linked list of tokens comes from. If it is converted runtime from for instance file input or macro expansion we need to add the involved overhead. But scanning a stored macro body is pretty efficient especially when the macro is part of the loaded macro package (format file). The third method is comparable with reading from file but here we need to add the overhead involved with storing the Lua output into data structures suitable for TeX's input mechanism, which can involve memory allocation outside the reserved pool of tokens. On modern systems that is not really a problem. It is good to keep in mind that when TeX was written much attention was paid to optimization and in LuaMetaTeX we even went a bit further, also because we know what kind of input, processing and output we're dealing with.

When reading from file or Lua output we interpret bytes turned utf numbers and that is when catcode regimes kick in: characters are interpreted according to the catcode properties: escape character (backslash), curly braces (grouping and arguments), dollars (math), etc. While with reading from token lists these catcodes are already taken care of and we're basically interpreting meanings instead of characters. By changing the catcode regime we can for instance typeset content verbatim from files and Lua strings but when reading from token lists we're sort of frozen. There are tricks to reinterpret the token list but that comes with overhead and limitations.

## 18.5 Macros

A macro can be seen as a named token with a meaning attached. In LuaMetaTeX macros can take up to 15 arguments (six more than regular TeX) that can be separated by so called delimiters. A token has a command property (operator) and a value (operand). Because a Unicode character doesn't need all four bytes of an integer and because in the engine numbers, dimensions and pointers are limited in size we can store all of these efficiently with the command code. Here the body of `\foo` is a list of three tokens:

```
\def\foo{abc} \foo \foo \foo
```

When the engine fetches a token from a list it will interpret the command and when it fetches from file it will create tokens on the fly and then interpret those. When a file or

list is exhausted the engine pops the stack and continues at the previous level. Because macros are already tokenized they are more efficient than file input. For more about macros you can consult the low level document about them.

The more you use a macro, the more it pays off compared to a file. However don't overestimate this, because in the end the typesetting and expanding all kind of other involved macros might reduce the file overhead to noise.

## 18.6 Token lists

A token list is like a macro but is part of the variable (register) system. It is just a list (so no arguments) and you can append and prepend to that list.

```
\toks123={abc}    \the\toks123
\scratchtoks{abc} \the\scratchtoks
```

Here `\scratchtoks` is defined with `\newtoks` which creates an efficient reference to a list so that, contrary to the first line, no register number has to be scanned. There are low level manuals about tokens and registers that you can read if you want to know more about this. As with macros the list in this example is three tokens long. Contrary to macros there is no macro overhead as there is no need to check for arguments.<sup>25</sup>

Because they use more or less the same storage method macros and token list registers perform the same. The power of registers comes from some additional manipulators in Lua $\TeX$  (and LuaMeta $\TeX$ ) and the fact that one can control expansion with `\the`, although that later advantage is compensated with extensions to the macro language (like `\protected` macro definitions).

## 18.7 Buffers

Buffers are something specific for Con $\TeX$ t and they have always been part of this system. A buffer is defined as follows:

```
\startbuffer[one]
line 1
line 2
\stopbuffer
```

Among the operations on buffers the next two are used most often:

<sup>25</sup> In LuaMeta $\TeX$  a macro without arguments is also quite efficient.

```
\typebuffer[one]
\getbuffer[one]
```

Scanning a buffer at the T<sub>E</sub>X end takes a little effort because when we start reading the catcodes are ignored and for instance backslashes and curly braces are retained. Hardly any interpretation takes place. The same is true for spacing, so multiple spaces are not collapsed and newlines stay. The tokenized content of a buffer is converted back to a string and that content is then read in as a pseudo file when we need it. So, basically buffers are files! In MkII they actually were files (in the `\jobname` name space and suffix `tmp`), but in MkIV they are stored in and managed by Lua. That also means that you can set them very efficiently at the Lua end:

```
\startluacode
buffers.assign("one",[
line 1
line 2
])
\stopluacode
```

Always keep in mind that buffers eventually are read as files: character by character, and at that time the content gets (as with other files) tokenized. A buffer name is optional. You can nest buffers, with and without names.

Because ConT<sub>E</sub>Xt is very much about re-use of content and selective processing we have an (already old) subsystem for defining named blocks of text (using `\begin...` and `\end...` tagging. These blocks are stored just like buffers but selective flushing is part of the concept. Think of coding an educational document with explanations, questions, answers and then typesetting only the explanations, or the explanation along with some questions. Other components can be typeset later so one can make for instance a special book(let) with answers that either of not repeats the questions. Here we need features like synchronization of numbers so that's why we cannot really use buffers. An alternative is to use xml and filter from that.

The `\definebuffer` command defines a new buffer environment. When you set buffers in Lua you don't need to define a buffer because likely you don't need the `\start` and `\stop` commands. Instead of `\getbuffer` you can also use `\getdefinedbuffer` with defined buffers. In that case the `before` and `after` keys of that specific instance are used.

The `\getinlinebuffer` command, which like the getters takes a list of buffer names, ignores leading and trailing spaces. When multiple buffers are flushed this way, spacing between buffers is retained.



The most important aspect of buffers is that the content is *not* interpreted and tokenized: the bytes stay as they are.

```
\definebuffer[MyBuffer]

\startMyBuffer
\bold{this is
a buffer}
\stopMyBuffer

\typeMyBuffer \getMyBuffer
```

These commands result in:

```
\bold{this is
a buffer}
```

**this is a buffer**

There are not that many parameters that can be set: before, after and strip (when set to no leading and trailing spacing will be kept. The `\stop...` command, in our example `\stopMyBuffer`, can be defined independent to do something after the buffer has been read and stored but by default nothing is done.

You can test if a buffer exists with `\doifelsebuffer` (expandable) and `\doifelsebufferempty` (unexpandable). A buffer is kept in memory unless it gets wiped clean with `resetbuffer`.

```
\savebuffer      [MyBuffer][temp]      % gets name: jobname-temp.tmp
\savebufferinfile[MyBuffer][temp.log] % gets name: temp.log
```

You can also stepwise fill such a buffer:

```
\definesavebuffer[slide]

\startslide
  \starttext
\stopslide
\startslide
  slide 1
\stopslide
text 1 \par
\startslide
```

```

    slide 2
\stopslide
text 2 \par
\startslide
    \stoptext
\stopslide

```

After this you will have a file `\jobname-slide.tex` that has the two lines wrapped as text. You can set up a ‘save buffer’ to use a different filename (with the `file` key), a different prefix using `prefix` and you can set up a directory. A different name is set with the `list` key.

You can assign content to a buffer with a somewhat clumsy interface where we use the delimiter `\endbuffer`. The only restriction is that this delimiter cannot be part of the content:

```
\setbuffer[name]here comes some text\endbuffer
```

For more details and obscure commands that are used in other commands you can peek into the source.

Using buffers in the `cld` interface is tricky because of the `catcode` magick that is involved but there are setters and getters:

<b>function</b>	<b>arguments</b>
<code>buffers.assign</code>	<code>name, content [,catcodes]</code>
<code>buffers.erase</code>	<code>name</code>
<code>buffers.prepend</code>	<code>name, content</code>
<code>buffers.append</code>	<code>name, content</code>
<code>buffers.exists</code>	<code>name</code>
<code>buffers.empty</code>	<code>name</code>
<code>buffers.getcontent</code>	<code>name</code>
<code>buffers.getlines</code>	<code>name</code>

There are a few more helpers that are used in other (low level) commands. Their functionality might adapt to their usage there. The `context.startbuffer` and `context.stopbuffer` are somewhat differently defined than regular `cld` commands.

## 18.8 Setups

A setup is basically a macro but is stored and accessed in a namespace separated from ordinary macros. One important characteristic is that inside setups newlines are ignored.

```
\startsetups MySetupA
  This is line 1
  and this is line 2
\stopsetups
```

```
\setup{MySetupA}
```

**This is line 1 and this is line 2**

A simple way out is to add a comment character preceded by a space. Instead you can also use `\space`:

```
\startsetups [MySetupB]
  This is line 1 %
  and this is line 2\space
  while here we have line 3
\stopsetups
```

```
\setup[MySetupB]
```

**This is line 1 and this is line 2 while here we have line 3**

You can use square brackets instead of space delimited names in definitions and also in calling up a (list of) setup(s). The `\directsetup` command takes a single setup name and is therefore more efficient.

Setups are basically simple macros although there is some magic involved that comes from their usage in for instance xml where we pass an argument. That means we can do the following:

```
\startsetups MySetupC
  before#1after
\stopsetups
```

```
\setupwithargument{MySetupC}{ {\em and} }
```

**before *and* after**

Because a setup is a macro, the body is a linked list of tokens where each token takes 8 bytes of memory, so MySetupC has 12 tokens that take 96 bytes of memory (plus some overhead related to macro management).

## 18.9 xml

Discussing xml is outside the scope of this document but it is worth mentioning that once an xml tree is read is, the content is stored in strings and can be filtered into T<sub>E</sub>X, where it is interpreted as if coming from files (in this case Lua strings). If needed the content can be interpreted as T<sub>E</sub>X input.

## 18.10 Lua

As mentioned already, output from Lua is stored and when a Lua call finishes it ends up on the so called input stack. Every time the engine needs a token it will fetch from the input stack and the top of the stack can represent a file, token list or Lua output. Interpreting bytes from files or Lua strings results in tokens. As a side note: Lua output can also be already tokenized, because we can actually write tokens and nodes from Lua, but that's more an implementation detail that makes the Lua input stack entries a bit more complex. It is normally not something users will do when they use Lua in their documents.

## 18.11 Protection

When you define macros there is the danger of overloading some defined by the system. Best use CamelCase so that you stay away from clashes. You can enable some checking:

```
\enablenonstandard[overloadmode=warning]
```

or when you want to quit on a clash:

```
\enablenonstandard[overloadmode=error]
```

When these trackers are enabled you can get around the check with:

```
\pushoverloadmode
...
\popoverloadmode
```

But delay that till you're sure that redefining is okay.

## 18.11 Colofon

Author           Hans Hagen  
ConT<sub>E</sub>Xt        2025.07.04 21:26  
LuaMetaT<sub>E</sub>X    2.11.07 | 20250705  
Support         www.pragma-ade.com  
                  contextgarden.net  
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## 19 Accuracy

low level

TEX

accuracy

## Contents

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### 19.1 Introduction

*This is work in progress, uncorrected.*

When you look at T<sub>E</sub>X and MetaPost output the accuracy of the rendering stands out, unless of course you do a sloppy job on design and interfere badly with the system. Much has to do with the fact that calculations are very precise, especially given the time when T<sub>E</sub>X was written. Because T<sub>E</sub>X doesn't rely on (at that time non-portable) floating point calculations, it does all with 32 bit integers, except in the backend where glue calculations are used for finalizing the glue values. It all changed a bit when we added Lua because there we mix integers and doubles but in practice it works out okay.

When looking at floating point (and posits) one can end up in discussions about which one is better, what the flaws for each are, etc. Here we're only interested in the fact that posits are more accurate in the ranges where T<sub>E</sub>X and MetaPost operate, as well as the fact that we only have 32 bits for floats in T<sub>E</sub>X, unless we patch more heavily. So, it is also very much about storage.

When you work with dimensions like points, they get converted to an integer number (the sp unit) and from that it's just integer calculations. The maximum dimension is 16383.99998pt, which already shows a rounding issue. Of course when one goes precise for sure there is some loss, but on the average we're okay. So, in the next example the two last rows are equivalent:

```
.1pt 0.1pt 6554sp
.2pt 0.2pt 13107sp
.3pt 0.3pt 19661sp
.1pt + .2pt 0.3pt 19661sp
```

When we're at the Lua end things are different, there numbers are mapped onto 64 bit floating point variables (doubles) and not all numbers map well. This is what we get when we work with doubles in Lua:

```
.1          0.1
.2          0.2
```



```

      .3                0.3
.1 + .2  0.300000000000000004

```

The serialization looks as if all is okay but when we test for equality there is a problem:

```

      .3 == .3  true
.1 + .2 == .3  false

```

This means that a test like this can give false positives or negatives unless one tests the difference against the accuracy (in MetaPost we have the `eps` variable for that). In  $\text{T}_{\text{E}}\text{X}$  clipping of the decimal fraction influences equality.

```

      \iflua{ .3 == .3 } \else N\fi  different
\iflua{ .1 + .2 == .3 } \else N\fi  different

```

The serialization above misguides us because the number of digits displayed is limited. Actually, when we would compare serialized strings the equality holds, definitely within the accuracy of  $\text{T}_{\text{E}}\text{X}$ . But here is reality:

```

      .3                .1 + .2
%0.10g  0.3            0.3
%0.17g  0.29999999999999999  0.300000000000000004
%0.20g  0.2999999999999999889  0.300000000000000004441
%0.25g  0.2999999999999999888977698  0.3000000000000000044408921

```

The above examples use 0.1, 0.2 and 0.3 and on a 32 bit float that actually works out okay, but `LuaMeta $\text{T}_{\text{E}}\text{X}$`  is 64 bit. Is this really important in practice? There are indeed cases where we are bitten by this. At the Lua end we seldom test for equality on calculated values but it might impact check for less or greater then. At the  $\text{T}_{\text{E}}\text{X}$  end there are a few cases where we have issues but these also relate to the limited precision. It is not uncommon to loose a few scaled points so that has to be taken into account then. So how can we deal with this? In the next section(s) an alternative approach is discussed. It is not so much the solution for all problems but who knows.

## 19.2 Posits

The next table shows the same as what we started with but with a different serialization.

```

      .1                0.1
      .2                0.2
      .3  0.300000001
.1 + .2  0.300000001

```

And here we get equality in both cases:

```
.3 == .3 true
.1 + .2 == .3 true
```

The next table shows what we actually are dealing with. The `\if`-test is not a primitive but provided by `ConTeXt`.

```
\ifpositunum{ .3 == .3 } Y\else N\fi equal
\ifpositunum{ .1 + .2 == .3 } Y\else N\fi equal
```

And what happens when we do more complex calculations:

```
math .sin(0.1 + 0.2) == math .sin(0.3) false
posit.sin(0.1 + 0.2) == posit.sin(0.3) true
```

Of course other numbers might work out differently! I just took the simple tests that came to mind.

So what are these posits? Here it's enough to know that they are a different way to store numbers with fractions. They still can loose precision but a bit less on smaller values and often we have relative small values in `TEX`. Here are some links:

<https://www.johngustafson.net/pdfs/BeatingFloatingPoint.pdf>  
<https://posithub.org/conga/2019/docs/14/1130-FlorentDeDinechin.pdf>

There are better explanations out there than I can provide (if at all). When I first read about these unums (a review of the 2015 book “The End of Error Unum Computing”) I was intrigued and when in 2023 I read something about it in relation to RISC-V I decided to just add this playground for the users. After all we also have decimal support. And interval based solutions might actually be good for MetaPost, so that is why we have it as extra number model. There we need to keep in mind that MetaPost in non scaled models also apply some of the range checking and clipping that happens in scaled (these magick 4096 tricks).

For now it is enough to know that it's an alternative for floats that *could* work better in some cases but not all. The presentation mentioned above gives some examples of physics constants where 32 posits are not good enough for encoding the extremely large or small constants, but for  $\pi$  it's all fine.<sup>26</sup> In double mode we actually have quite good precision compared to 32 bit posits but with 32 bit floats we gain some. Very small

<sup>26</sup> Are 64 bit posits actually being worked on in softposit? There are some commented sections. We also need to patch some unions to make it compile as C.

numbers and very large numbers are less precise, but around 1 we gain: the next value after 1 is 1.0000001 for a float and 1.000000008 for a posit (both 32 bit). So, currently for MetaPost there is no real gain but if we'd add posits to T<sub>E</sub>X we could gain some because there a halfword (used for storing data) is 32 bit.

But how about T<sub>E</sub>X? Per April 2023 the LuaMetaT<sub>E</sub>X engine has native support for floats (this in addition to Lua based floats that we already had in ConT<sub>E</sub>Xt). How that works can be demonstrated with some examples. The float related commands are similar to those for numbers and dimensions: `\floatdef`, `\float`, `\floatexpr`, `\iffloat`, `\ifzerofloat` and `\ifintervalfloat`. That means that we also have them as registers. The `\positdef` primitive is similar to `\dimensiondef`. When a float (posit) is seen in a dimension context it will be interpreted as points, and in an integer context it will be a rounded number. As with other registers we have a `\newfloat` macro. The `\advance`, `\multiply` and `\divide` primitives accept floats.

```
\scratchdimen=1.23456pt
\scratchfloat=1.23456
```

We now use these two variables in an example:

```
\setbox0\hbox to \scratchdimen {x}\the\wd0
\scratchdimen \dimexpr \scratchdimen * 2\relax
\setbox0\hbox to \scratchdimen {x}\the\wd0
\advance \scratchdimen \scratchdimen
\setbox0\hbox to \scratchdimen {x}\the\wd0
\multiply\scratchdimen by 2
\setbox0\hbox to \scratchdimen {x}\the\wd0
```

```
1.23456pt
2.46912pt
4.93823pt
9.87646pt
```

When we use floats we get this:

```
\setbox0\hbox to \scratchfloat {x}\the\wd0
\scratchfloat \floatexpr \scratchfloat * 2\relax
\setbox0\hbox to \scratchfloat {x}\the\wd0
\advance \scratchfloat \scratchfloat
\setbox0\hbox to \scratchfloat {x}\the\wd0
\multiply\scratchfloat by 2
\setbox0\hbox to \scratchfloat {x}\the\wd0
```

## Posits

1.23456pt  
 2.46912pt  
 4.93823pt  
 9.87648pt

So which approach is more accurate? At first sight you might think that the dimensions are better because in the last two lines they indeed duplicate. However, the next example shows that with dimensions we lost some between steps.

```

\scratchfloat \floatexpr \scratchfloat * 2\relax \the\scratchfloat
\advance \scratchfloat \scratchfloat \the\scratchfloat
\multiply\scratchfloat by 2 \the\scratchfloat

```

1.2345599979162216187  
 2.4691199958324432373  
 4.9382399916648864746  
 9.8764799833297729492

One problem with accuracy is that it can build up. So when one eventually does some comparison the expectations can be wrong.

```

\dimen0=1.2345pt
\dimen2=1.2345pt

```

```

\ifdim \dimen0=\dimen2 S\else D\fi \space +0sp: [dim]
\ifinterval\dim0sp\dimen0 \dimen2 0\else D\fi \space +0sp: [0sp]

```

```

\advance\dimen2 1sp

```

```

\ifdim \dimen0=\dimen2 S\else D\fi \space +1sp: [dim]
\ifinterval\dim 1sp \dimen0 \dimen2 0\else D\fi \space +1sp: [1sp]
\ifinterval\dim 1sp \dimen2 \dimen0 0\else D\fi \space +1sp: [1sp]
\ifinterval\dim 2sp \dimen0 \dimen2 0\else D\fi \space +1sp: [2sp]
\ifinterval\dim 2sp \dimen2 \dimen0 0\else D\fi \space +1sp: [2sp]

```

```

\advance\dimen2 1sp

```

```

\ifinterval\dim 1sp \dimen0\dimen2 0\else D\fi \space +2sp: [1sp]
\ifinterval\dim 1sp \dimen2\dimen0 0\else D\fi \space +2sp: [1sp]
\ifinterval\dim 5sp \dimen0\dimen2 0\else D\fi \space +2sp: [5sp]
\ifinterval\dim 5sp \dimen2\dimen0 0\else D\fi \space +2sp: [5sp]

```

Here we show a test for overlap in values, the same can be done with integer numbers (counts) and floats. This interval checking is an experiment and we'll see it if gets used.

```
S +0sp: [dim] 0 +0sp: [0sp]
D +1sp: [dim] 0 +1sp: [1sp] 0 +1sp: [1sp] 0 +1sp: [2sp] 0 +1sp: [2sp]
D +2sp: [1sp] D +2sp: [1sp] 0 +2sp: [5sp] 0 +2sp: [5sp]
```

There are also `\ifintervalfloat` and `\ifintervalnum`. Because I have worked around these few scaled point rounding issues for decades, it might actually take some time before we see the interval tests being used in `ConTEXt`. After all, there is no reason to touch somewhat tricky mechanism without reason (read: users complaining).

To come back to posits, just to be clear, we use 32 bit posits and not 32 bit floats, which we could have but that way we gain some accuracy because less bits are used by default for the exponential.

In `ConTEXt` we also provide a bunch of pseudo primitives. These take one float: `\pfsin`, `\pfcos`, `\pftan`, `\pfasin`, `\pfacos`, `\pfatan`, `\pfsinh`, `\pfcosh`, `\pftanh`, `\pfasinh`, `\pfacosh`, `\pfatanh`, `\pfsqrt`, `\pfllog`, `\pfexp`, `\pfceil`, `\pffloor`, `\pfround`, `\pfabs`, `\pfrad` and `\pfdeg`, whiel these expect two floats: `\pfatantwo`, `\pfpow`, `\pfmod` and `\pfrem`.

## 19.3 MetaPost

In addition to the instances `metafun` (double in `LMTX`), `scaledfun`, `doublefun`, `decimalfun` we now also have `positfun`. Because we currently use 32 bit posits in the new number system there is no real gain over the already present 64 bit doubles. When 64 bit posits show up we might move on to that.

## 19.4 Lua

We support posits in Lua too. Here we need to create a posit user data object. The usual metatable magick kicks in:

```
local p = posit.new(123.456)
local q = posit.new(789.123)
local r = p + q
```

Here we just mention what is currently interface. The management functions are: `new`, `copy`, `tostring`, `tonumber`, `integer`, `rounded`, `toposit` and `fromposit`. The usual operators are also supported: `+`, `-`, `*`, `/`, `^`, as well as the binary `|`, `&`, `~`, `<<` and `>>`. We can

compare with `==`, `>=`, `<=` and `~=`. The more verbose `bor`, `bxor`, `band`, `shift`, `rotate` are there too.

There is a subset of math provided: `min`, `max`, `abs`, `conj`, `modf`, `acos`, `asin`, `atan`, `ceil`, `cos`, `exp`, `exp2`, `floor`, `log`, `log10`, `log1p`, `log2`, `logb`, `pow`, `round`, `sin`, `sqrt` and `tan`. Somewhat special are `NaN` and `NaR`.

Currently integer division (`//`) and modulo (`%`) are not available, but that might happen at some time.

## 19.4 Colofon

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## 20 Balancing

low level

TEX

balancing



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### 20.1 Introduction

*This is work in progress as per end 2024 these mechanisms are still in flux. We expect them to be stable around the ConT<sub>E</sub>Xt meeting in 2025. The text is not corrected, so feel free to comment.*

This manual is about a new (sort of fundamental) feature that got added to LuaMetaT<sub>E</sub>X when we started upgrading column sets. In T<sub>E</sub>X we have a par builder that does a multi-pass optimization where it considers various solutions based on tolerance, penalties, demerits etc. The page builder on the other hand is forward looking and backtracks to a previous break when there is an overflow. The balancing mechanism discussed here is basically a page builder operating like the par builder: it looks at the whole picture.

In order to make this a useful mechanism the engine also permits intercepting the main vertical list, so we start by introducing this.

### 20.2 Intercepting the MVL

When content gets processed it's added to a list. We can be in horizontal mode or vertical mode (let's forget about math mode). In vertical mode we can be in a box context (say `\vbox`) or in what is called the main vertical list: the one that makes the page. But what is page? When T<sub>E</sub>X has collected enough to match the criteria set by `\pagegoal` which starts out as `\vsize`, it will call the so called output routine which basically is expanding the `\output` token list. That routine had do so something with the box that has the collected material. It can become a page, likely with the content wrapped in a page body with headers and footers and such, but it can also be stored for later assembly, for instance in multiple columns, or after some analysis fed back into the main vertical list.

For various mechanisms it matters if they are used inside a contained boxed environment or in the more liberal main vertical list (from now on called mvl). That's why we can intercept the mvl and use it later. Intercepting works as follows:

```
\beginmvl 1
various content
\endmvl
```

```
\beginmvl 2
various content
\endmvl
```

When at some point you want this content, you can do this:

```
\setbox\scratchboxone\flushmvl 2
\setbox\scratchboxtwo\flushmvl 1
```

and then do whatever is needed. You can see what goes on with:

```
\tracingmvl 1
```

There is not much more to say other than that this is the way to operate on content as if it were added to the page which can be different from collecting something in a vertical box. Think of various callbacks that can differ for the mvl and a box.

The `\beginmvl` primitive takes a number or a set of keywords, as in:

```
\beginmvl
  index 1
  options \numexpr "01 + "04\relax
\relax
```

There is of course some possible interference with mechanism that check the page properties like `\pagegoal`. If needed one can check this:

```
\ifcase\mvlcurrentlyactive
  % main mvl
\or
  % first one
\else
  % other ones
\fi
```

Possible applications of this mechanism are the mentioned columns and parallel, independent, streams. However for that we need to be able to manipulate the collected content. Actually, the next manipulator preceded the capturing, because we first wanted to make sure that what we had in mind made sense.


The `\beginmvl` also accepts keywords. You can specify an `index` (an integer), a `prevdepth` (dimensions) and options (an integer bitset). Possible option bit related values are:

```
0x1 ignore prevdepth \ignoreprevdepthmvloptioncode
0x2 no prevdepth \noprevdepthmvloptioncode
0x4 discard top \discartopmvloptioncode
0x8 discard bottom \discartbottommvloptioncode
```

Here the last column is a numeric alias available in `ConTEXt`. More options are likely to show up. When we eventually will balance these lists the routine will deal with the discardables (like glue) but one can also remove them via the options.

```
\beginmvl
  index      1
  prevdepth  0pt
  options    \discartopmvloptioncode
\relax
\scratchdimen\prevdepth
\dontleavehmode
\quad\the\mvlcurrentlyactive\quad\the\scratchdimen
\quad\blackrule[height=\strutht,depth=\strutdp,color=darkred]
\endmvl

\ruledhbox {\llap{1\quad}}\flushmvl 1}
```

1 1 0.0pt 

```
\beginmvl
  index      2
  options    \numexpr
             \ignoreprevdepthmvloptioncode
             + \discartopmvloptioncode
  \relax
\relax
\scratchdimen\prevdepth
\dontleavehmode
\quad\the\mvlcurrentlyactive\quad\the\scratchdimen
\quad\blackrule[height=\strutht,depth=\strutdp,color=darkred]
```

```
\endmvl
```

```
\ruledhbox {\llap{2\quad}\flushmvl 2}
```

```
2 | 2. -1000.0pt. ■
```

```
\beginmvl 3 % when no keywords are used we expect a number
```

```
\scratchdimen\prevdepth
```

```
\dontleavehmode
```

```
\quad\the\mvlcurrentlyactive\quad\the\scratchdimen
```

```
\quad\blackrule[height=\strutht,depth=\strutdp,color=darkred]
```

```
\endmvl
```

```
\ruledhbox {\llap{3\quad}\flushmvl 3}
```

```
3 | 3. 0.0pt. ■
```

```
\beginmvl index 4 options 1
```

```
\scratchdimen\prevdepth
```

```
\dontleavehmode
```

```
\quad\the\mvlcurrentlyactive\quad\the\scratchdimen
```

```
\quad\blackrule[height=\strutht,depth=\strutdp,color=darkred]
```

```
\endmvl
```

```
\ruledhbox {\llap{4\quad}\flushmvl 4}
```

```
4 | 4. -1000.0pt. ■
```

## 20.3 Balancing

Balancing is not referring to balancing columns but to ‘a result that looks well balanced’. Just like we want lines in a paragraph to look consistent with each other, something that is reflected in the (adjacent) demerits, we want the same with vertical split of pieces. For this purpose we took elements of the par builders to construct a (page) snippet builder. Here are some highlights:

- Instead of a pretolerance, tolerance and emergency pass we only enable the last two. In the par builder the pretolerance pass is the one without hyphenation.
- We seriously considered vertical discretionaries but eventually rejected the idea: we just don't expect users to go through the trouble of adding lots of split related pre, post and replace content. It's not hard to support it but in the end it also interfered with other demands that we had. We kept the code around for a while but then

removed it. To mention one complication: if we add some new node we also need to intercept it in various callbacks that we already have in place in ConT<sub>E</sub>Xt. As with horizontal discretionaries, we then need to go into the components and sometimes even need to make decisions what can not yet be made.

- As with the par builder, T<sub>E</sub>X will happily produce an overfull box when no solution is possible that fits the constraints. In a paragraph there are plenty spaces (with stretch) and discretionaries (with components that vary in width) which enlarges the solution space. In vertical material there is less possible so there an emergency pass really makes sense: better be underful than overful.
- In many cases there is no stretch available. There are also widow, club, shape and orphan penalties that can limit the solution space.
- When we look at splitting pages (and boxes) we see (split) top skip kick in. This is something that we need to provide one way or the other. And as we have to do that, we can as well provide support for bottom skip. A horizontal analogue is protrusion, something that also has to be taken into account in a rather dynamic way, at the beginning or end of the currently analyzed line.
- There is no equivalent of hanging indentation but a shape makes sense. Here the shape defines heights, top and bottom skips and maybe more in the future. For that reason we use a keyword driven shape.
- Because we have so called par passes, it made sense to have something similar for balancing. This gives is the opportunity to experiment with various variables that drive the process.
- For those who read what we wrote about the par builder, it will not come as surprise that we also added extensive tracing and a callback for intercepting the results. This makes it possible to show the same detailed output as we can do for par passes.

It's about time for some examples but before we come to that it is good to roughly explain how the page builder works. When the page builder is triggered it will take elements from the contributions list and add them to the page. When doing that it keeps track of the height and depth as contributed by boxes and rules. Because it will discard glue and kerns it does some checking there. An important feature is that the depth is added in a next iteration. The routine also needs to look at inserts. The variables `\pagegoal` (original `\vsize` minus accumulated insert heights) and `\pagetotal` are compared and when we run over the target height the accumulated stretch and shrink in glue (when present) will be used to determine how bad this break is. If it is too bad, the previous best break will be taken. Penalties can make a possible break

more or less attractive. When the output routine gets a split of page, the total is not reliable because we can have backtracked to the previous break. In LuaMetaT<sub>E</sub>X we have some more variables, like `\pagelastheight`, that give a better estimate of what we got.

In order to make the first lines align properly relative to the top of the page there is a variable `\topskip`. The height of the first line is at least that amount. The correction is calculated when the first contribution happens: a box or rule.

When we look at the balancer it is good to keep in mind that where the page builder stepwise adds and checks, the balancer looks at the whole picture. The page builder does a decent job but is less sophisticated than the par builder. There is a badness calculation, penalties are looked at, glue is taken into account but there are no demerits.

We want the balancer to work well with column sets that are very much grid based. But in getting there we had some hurdles to take. Because the algorithm (like the par builder) happily results in overfull boxes unless emergency stretch is set, pages can overflow. When there is no stretch and/or shrink using emergency stretch can give an underfull page.

The way out of this is to have non destructive trial passes and decrease the number of lines. Of course we can get short pages but when for instance it concerns a section title that gets moved this is no big deal. In a similar fashion splitting a multi-line formula is also okay.

- Collect the content in an mvl list and after that's done put the result in a box.
- Set up a balance shape that specifies the slots in in columns (normally a column is just a blob of text).
- Perform a trial balance run. As soon as an overfull page is seen, adapt the balance shape and do a new trial run.
- When we're fine, either because we reached the end without overfull column or by passing the set deadcycles value, quit the trial process and balance the original list using the most recent balance shape.
- Flush the result by fetching the topmost from the result split collection and feed it into the page flow. The boxed pseudo page will happily trigger the output routine that in turn construct the final page.

At some point we decided to support multiple mvl streams and therefore changed the last mentioned step. Because we store the whole column set we can as well also store

the assembled page bodies. This way we can flush different streams into the same result.

- Flush the result by fetching the topmost from the result split collection and feed it into the page flow. Do this for every saved (mvl) stream.
- When we're done, the boxed pseudo pages will be flushed as pages. In the process, for every page we identify marks.

We are now ready to look at some examples. Here we also show what balance shapes do. These basically describe a sequence of slots to be filled. The last specification is used when we exceed the number of defined slots. These are just examples of simple situations, for real applications more code is needed.

We start with some content in a box. This can of course be a flushed mvl but here we just set it directly:

```
\setbox\scratchboxone\vbox\bgroup
  \hsize.30\hsize
  \samplefile{tufte}
\egroup
```

We will split this box in columns. If you are familiar with T<sub>E</sub>X you might know that a paragraph of text can follow a shape defined by `\parshape`. In a similar way as lines are split by width, we can split a vertical list by height. For that we define a balance shape:

```
\balanceshape 3
  vsize      12\lineheight
  topskip    \strutht
  bottomskip \strutdp
next
  vsize      5\lineheight
  topskip    \strutht
  bottomskip \strutdp
next
  vsize      8\lineheight
  topskip    \strutht
  bottomskip \strutdp
\relax
```

```
\setbox\scratchboxtwo\vbalance\scratchboxone
```

Contrary to a `\parshape`, a `\balanceshape` is not wiped after the work is done. It also expects keys and values. As with `\parpasses` each step is separated by next. This makes it an extensible mechanism. Finally we will split the box according to this shape:

```
\hbox \bgroup
  \localcontrolledendless {%
    \ifvoid\scratchboxtwo
      \expandafter\quitloop
    \else
      \setbox\scratchbox\ruledhbox\bgroup
        \vbalancedbox\scratchboxtwo
      \egroup
      \vbox to 12\lineheight \bgroup
        \box\scratchbox
        \vfill
      \egroup
      \hskip1em
    \fi
  }\unskip
\egroup
```

The result is shown here:

We thrive in information-thick worlds because of our marvelous and everyday capacity to select, edit, single out, structure, highlight, group, pair, merge, harmonize, synthesize, focus, organize, condense, reduce, boil down, choose, categorize, catalog, classify, list, abstract, scan, look into, idealize, isolate, discriminate, distinguish, screen, pigeonhole, pick
---

over, sort, integrate, blend, inspect, filter, lump, skip, smooth, chunk, average, approximate, cluster, aggregate, outline, summarize, itemize, review,
--

dip into, flip through, browse, glance into, leaf through, skim, refine, enumerate, glean, synopsisize, winnow the wheat from the chaff and separate the sheep from the goats.
--

Like the par builder we can end up with overfull boxes but we can deal with that by using trial runs.

```
\setbox\scratchboxtwo\vbalance\scratchboxone trial
```

In that case the result is made from empty boxes so the original is not disturbed. Here we show an overflow, so in the first resulting box you can compare the height with the



requested one and when it's larger you can decide to decrease the first height in the shape and try again.

Many readers will skim over formulas on their first reading of your exposition. Therefore, your sentences should flow smoothly when all but the simplest formulas are replaced by "blah" or some other grunting noise.

test

Many readers will skim over formulas on their first reading of your exposition. Therefore, your sentences should flow smoothly when all but the sim-

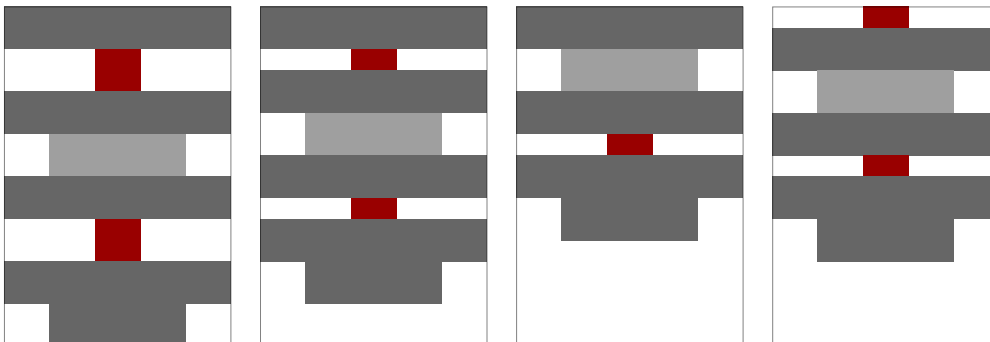
plest formulas are replaced by "blah" or some other grunting noise.

Of course that involves some juggling of the shape but after all we have Lua at our disposal so in the end it's all quite doable.

	<b>real</b>	<b>target</b>
1	167.8961pt	156.95874pt
2	65.39948pt	65.39948pt
3	49.17705pt	104.63916pt

Because the balancer can produce what otherwise the page builder produces, we need to handle the equivalent of top skip which is what the already shown top keyword takes care of. This means that the current slice (think current line in the par builder) has to take that into account. This can be compared to the left- and right protrusion in the par builder. When we typeset on a grid we have an additional demand.

When we surround (for instance a formula) with halfline spacing, we eventually have to return on the grid. One complication is that when we are in grid mode and use half line vertical spacing, we can end up in a situation where the initial half line space is on a previous page. That means that we need to use a larger top skip. This is not something that we want to burden the balancer with but we have ways to trick it into taking that compensation into account.



## Balancing

However, when we split in the middle of that segment, we can end up with a half line skip in a next slot because  $\text{T}_{\text{E}}\text{X}$  will remove glue at the edge. So we end up with what we see in the third sequence above. We deal with that in a somewhat special way: a box as a discardable field which value will be taken into account as additional top value. That field is set and reset by glue options  $\text{O}\times\text{20}$  and  $\text{O}\times\text{40}$  that can be manipulated in Lua as part of some spacing model. Here we suffice by mentioning that it makes sure that (as in the fourth blob above) at the top we have a half line spacing.

## 20.4 Forcing breaks

Because the initial application of balancing was in column sets, we also need the ability to goto a next slot (step in a shape), column (possibly more steps), page (depending on the page state), and spread (for instance if we are double-sided). For this we use `\balanceboundary`. It takes two values and when the boundary node triggers a callback in the builder these are passed along with a shape identifier and current shape slot. That callback can then signal back that we need to try a break here with a given penalty. Assuming that at the Lua end we know at which slot we have a slot, column, page or spread break. Multiple slots can be skipped by multiple boundaries. There is one pitfall: we need something in a slot in order to break at all, so one ends up with for instance:

```
\balanceboundary 3 1\relax
\vskip\zeropoint
\balanceboundary 3 0\relax
\vskip\zeropoint
\balanceboundary 3 0\relax
```

Here the 3 is just some value that the callback can use to determine its action (like goto a next page) and the second value provides a detail. Of course all depends on the intended usage. By using a callback we can force breaks while not burdening the engine with some hard coded solution. For example, in  $\text{ConT}_{\text{E}}\text{Xt}$  we used these (the values are these from experiments and might change:

first	second	action	user interface
1	1 or 0	goto next spread (1 initial, 0 follow up)	<code>\page[spread]</code>
2	1 or 0	goto next page (idem)	<code>\page</code>
3	1 or 0	goto next column (idem)	<code>\column</code>
4	1 or 0	goto next slot (idem)	<code>\column[slot]</code>
5	n	next slot when more than n lines	<code>\testroom[5]</code>
6	s	next slot when more than s scaled points	<code>\testroom[80pt]</code>

## 20.5 Marks

It is possible to synchronize the marks with those in the results of balanced segments with a few Lua helpers that do the same as the page builder does at the start of a page, while packaging the page and when wrapping it up. So, instead of split marks we can have real marks.

## 20.6 Inserts

Before we go into detail, we want to point out that when implementing a (balancing) mechanism as introduced above, decisions have to be made. In traditional  $\text{T}_{\text{E}}\text{X}$  there is for instance an approach to inserts that involves splitting them over pages. In our case that is a bit harder to do but there are ways to deal with it. When deciding on an approach it helps that we know a bit what situations occur and where we can put some constraints. One can argue that solutions should be very generic because (for instance) a publisher has some specific demands but in practice those are not our audience. In decades of developing  $\text{LuaT}_{\text{E}}\text{X}$  and  $\text{LuaMetaT}_{\text{E}}\text{X}$  it's ( $\text{ConT}_{\text{E}}\text{Xt}$ ) user demands and challenges that drives what gets implemented. Publishers, their suppliers, and large scale (commercial) users are pretty silent when it comes to development (and supporting it) while users communicate via meetings and mailing lists. Also, rendering of documents that have notes are often typeset kind of traditional.

Users on the other hand have come up with demands for columns, typesetting on the grid, multiple notes, balancing, and parallel content streams. The picture we get from that makes us confident that what we provide is generally enough and as users understand the issues at hand (maybe as side effect of struggling with solutions) it's not that hard to explain why constraints are in place. It makes more sense to have a limited reliable mechanism that deals with the kind of (foot)notes that known users need than to cook up some complex mechanism that caters potential specific demands by potential users. Of course we have our own challenges to deal with, even if the resulting features will probably not be used that often. So here are the criteria that make sense:

- We can assume a reasonable amount of notes.
- These are normally small with no (vertical) whitespace.
- Notes taking multiple lines may split.
- But we need to obey widow and club penalties.
- There can be math formulas but mostly inline.
- We need to keep them close to where they are referred from.

But,

- We can ignore complex conflicting demands.
- As long as we get some result, we're fine.
- So users have to check what comes out.
- We don't assume fully automated unattended usage.

And of course:

- Performance should be acceptable.
- User interfaces should be intuitive.
- Memory consumption should be reasonable.

We have users who use multiple note classes so that also has to be handled but again we don't need to come up with solutions that solve all possible demands. We can assume that when a book is published that needs them, the author will operate within the constraints.

We mentioned footnotes being handled by the page builder so how about them in these balanced slots? Given the above remarks, we assume sane usage, so for instance columns that have a single slot with possibly fixed content at the top or bottom (and maybe as part of the stream). The balancer handles notes by taking their height into account and when a result is used one can request the embedded inserts and deal with them. Again this is very macro package dependent. Among the features dealt with are space above and between a set of notes, which means that we need to identify the first and successive notes in a class. Given how the routine works, this is a dynamic feature of a line: the amount of space needed depends on how many inserts are within a slot. When we did some extreme tests with several classes of notes and multiple per column we saw runtime increasing because instead of a few passes we got a few hundred. In an extreme case of 800 passes to balance the result we noticed over four million checks for note related spacing. We could bring that down to one tenth so in the end we are still slower but less noticeable. Here are the helper primitives for inserts:

```
<state> = \boxinserts <box>
<box>    = \vbalancedinsert <box> <class>
<state> = \boxinserts <box>
```

A (foot)note implementation is very macro package dependent so the next example is just that: an example of using the available primitive. We start by populating a mvl with a sample text and a single footnote.

```
\begingroup
  \forgetall
  \beginmvl
```

```

index 5
options \numexpr
  \ignoreprevdepthmvloptioncode
  + \discardtopmvloptioncode
\relax
\relax
\hsize .4tw
Line 1 \par Line 2 \footnote {Note 1} \par Line 3 \par
Line 4 \footnote {Note 2} \par Line 5 \par Line 6 \par
\endmvl
\endgroup

```

We fetch the footnote number, which is one of many possible defined inserts

```

\cdef\currentnote{footnote}%
\scratchcounter\currentnoteinsertionnumber

```

The quick and dirty balancer uses a simple shape of 5 lines with normal strut properties. From the balanced result we take two columns. We test if there is an insert and take action when there is. Here we just filter the footnotes but there can of course be more. We overlay these notes over (under) the column that has them. So we work per column.

```

\beginngroup
  \setbox\scratchboxone\flushmvl 5
  \balanceshape 1
    vsize      5lh
    topskip    1sh
    bottomskip 1sd
  \relax
  \setbox\scratchboxtwo\vbalance\scratchboxone
  \ruledhbox \bgroup
    \localcontrolledrepeat 2 {
      \ifnum\currentloopiterator > 1
        \hskip2\emwidth
      \fi
      \setbox\scratchboxthree\vbancedbox\scratchboxtwo \relax
      \ifnum\boxinserts\scratchboxthree > 3
        \setbox\scratchboxfour\vbancedinsert
          \scratchboxthree\scratchcounter
        \wd\scratchboxfour 0pt
        \box\scratchboxfour

```

```

\fi
\box\scratchboxthree
}\unskip
\egroup
\endgroup

```

The result is:

Line 1	Line 4 <sup>28</sup>
Line 2 <sup>27</sup>	Line 5
Line 3	Line 6
<sup>27</sup> Note 1	<sup>28</sup> Note 2

As we progressed we realized that the ‘balancer’ used in column sets can also be used for single columns and we can even support a mix of single and multi columns. There is however a problem: within a mvl we can deal with spacing but we can't do that reliably across mvl's and especially when we cross a page it becomes hard to identify if some (vertical) spacing is needed; we don't want it at the bottom or top of a page. This feature is too experimental to be discussed right now.

We assumed reasonable notes to be used but even if a user tries to keep notes small and avoid too many, there are cases where they might look like a paragraph and when there are more in a row, it might be that a column overflows. This is why we have some support for split notes. This is accomplished by two additional commands:

```

\setbox\scratchboxone\vbalance\scratchboxone\relax
\vbalanceddeinsert\scratchboxone\relax

```

Here we convert inserts in such a way that they are taken into account by the balancer so that multi-slot optimization takes place. Afterwards, when we loop over the result we can reconstruct the inserts:

```

\setbox\scratchboxtwo\vbancedbox\scratchboxone
\vbancedreinsert\scratchboxtwo\relax

```

Among the reasons that these are explicit actions, is that we want to experiment but also be able to see the effect by selectively enabling it. You can get better results by forcing depth correction.

```

\setbox\scratchboxone\vbalance\scratchboxone
\vbalanceddeinsert\scratchboxone forceddepth\relax

```

This will use the depth as defined by `\insertlinedepth` which is an insert class specific parameter, but discussing details of inserts is not what we do here. The reason for using a `\relax` in the above examples is that we want to stress that when keywords are involved, you need to prevent look-ahead, especially when an `\if...` or expandable loop follows, which is not uncommon when we balance.

It is possible to define top and bottom inserts but of course these need to be filtered and placed at the  $\TeX$  end, so this is macro package specific. Here we just mention that it is possible to set `\insertstretch` and `\insertshrink` which will be taken into account. However, this can result in overlap so if indeed stretch or shrink is applied, the `handle_uinsert` callback should be used for bringing what actually gets inserted to the right dimensions. For now we consider this an experimental feature.

## 20.7 Discardables

This is a preliminary explanation.

```

\begingroup
  \beginmvl
    index 5
    options \numexpr
      \ignoreprevdepthmvl\optioncode
      + \discardtopmvl\optioncode
    \relax
  \relax
  \hsize .4tw
  \par
  \vskip0pt
  {\darkred \hrule discardable height 1sh depth 1sd width 1em}
  \par
  % we need the strut because the rule obscures it .. todo
  \dorecurse{8}{\strut Line #1 \par}
  \vskip\zeropoint
  {\darkblue \hrule discardable height 1sh depth 1sd width 1em}
  \par
  \endmvl
\endgroup

\setbox\scratchboxone\flushmvl 5
\balanceshape 1
  vsize      5lh

```

```

topskip    1sh % see comment above
bottomskip 1sd
options    3

```

```
\relax
```

```
\setbox\scratchboxtwo\vbalance\scratchboxone\relax % lookahead
```

```
\hpack \bgroup
```

```
\localcontrolledrepeat 3 {
```

```
\ifvoid\scratchboxtwo\else
```

```
\setbox\scratchboxthree\vbancedbox\scratchboxtwo
```

```
\ifvoid\scratchboxthree\else
```

```
\dontleavehmode\llap{[\the\currentloopiterator]\quad}%
```

```
\ruledhpack{\box\scratchboxthree}\par
```


```
\fi
```

```
\hskip 4em
```

```
\fi
```

```
}\unskip
```

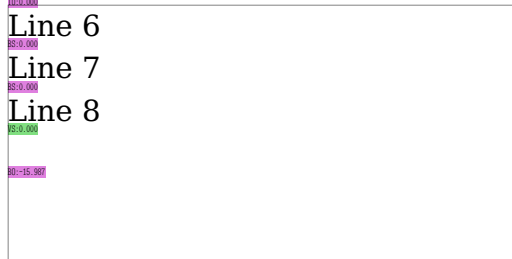
```
\egroup
```

[1] 

[2] 

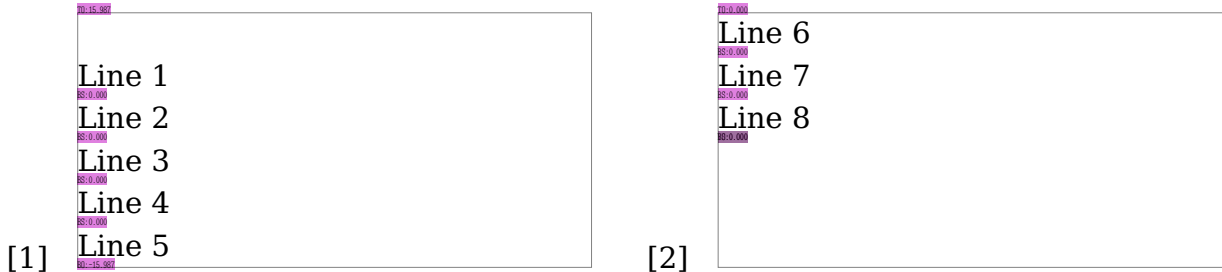
When at the top, the rule will be ignored and basically sticks out. When at the bottom the rule might end up in a zero dimension box. With `\vbanceddiscard\scratchboxtwo` they will become an `\nohrule`. Basically we're talking of optional content. The options bitset in the shape definition tells if we have a top (1) and/ or bottom (2), here we have both (3) but in for instance column sets it depends.

[1] 

[2] 



Here we actually still have the rule but marked as invisible. So, topskip has a negative amount. In the next case the remove keyword makes the rule go away in which case we also adapt the topskip accordingly.



You need to juggle a bit with skips and penalties to get this working as you like. Instead of rules you can also use boxes, for example before:

```
\vskip\zeropoint
\ruledvbox discardable {\hpack{\strut BEFORE}}
\par
```

and after:

```
\forgetall \par \vskip\zeropoint
\ruledvbox discardable {\hpack{\strut AFTER}}%
\penalty\minusone % !
\par
```

It currently is a playground so it might (and probably will) evolve. Although it was also made for a specific issue it might have other usage.

## 20.8 Penalties

*todo*

```
\showmakeup[vpenalty,line]
\balancefinalpenalties 6 10000 9000 8000 7000 6000 5000\relax
\balancevsize 5\lineheight
\setbox\scratchbox\vbox\bgroup
  \dorecurse{1}{\samplefile{tufte}\footnote{!}\par}
\egroup
\vbalance\scratchbox
```

## 20.9 Passes

In LuaMetaT<sub>E</sub>X the par builder has been extended with additional features (like orphan, toddler and twin control) and the ability to define and apply multiple passes over the paragraph to get the best result. The balancer has a similar feature: `\balancepasses`. As with `\parpasses` we have an infrastructure for tracing.

```
% threshold
% tolerance
% looseness
% adjdemerits
% originalstretch
% emergencystretch
% emergencyfactor
% emergencypercentage
```

## 20.9 Colofon

Author	Hans Hagen
ConT <sub>E</sub> Xt	2025.07.04 21:26
LuaMetaT <sub>E</sub> X	2.11.07   20250705
Support	<a href="http://www.pragma-ade.com">www.pragma-ade.com</a>
	<a href="http://contextgarden.net">contextgarden.net</a>
	<a href="mailto:ntg-context@ntg.nl">ntg-context@ntg.nl</a>

## **21 Lines**

low level

TEX

lines

## Contents

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### 21.1 Introduction

There is no doubt that  $\text{T}_{\text{E}}\text{X}$  does an amazing job of “breaking paragraphs into lines” where a paragraph is a sequence of words in the input separated by spaces or its equivalents (single line endings turned space). The best descriptions of how that is done can be found in Don Knuths “The  $\text{T}_{\text{E}}\text{X}$  Book”, “ $\text{T}_{\text{E}}\text{X}$  The Program” and “Digital Typography”. Reading and rereading the relevant portions of those texts is a good exercise in humility.

That said, whatever follows here builds upon what Knuth gave us and in no way we pretend to do better than that. It started out as a side track of improving rendering math in combination with more control over breaking inline math. It pretty much about having fun with the par builder but in the end can also help make your results look better. This is especially true for proze.

Trying to describe the inner working of the par builder makes no sense. Not only is it kind of complex, riddled with magic constants and heuristics, but there is a good chance for us to talk nonsense thanks to misunderstanding. However, some curious aspects will be brought up. Consider what follows a somewhat naive approach and whatever goes wrong, blame the authors, not  $\text{T}_{\text{E}}\text{X}$ .

If you're one of those reader who love to complain about the bad manuals, you can stop reading here. There is plenty said in the mentioned books but you can also consult Viktor Eijkhouts excellent “ $\text{T}_{\text{E}}\text{X}$  by Topic” (just search the web for how the get these books). If you're curious and in for some adventure, keep reading.

## 21.2 Warning

This is a first version. What is described here will stay but is still experimental and how it evolves also depends on what demands we get from the users. We have defined some experimental setups in ConT<sub>E</sub>Xt. We will try to improve the explanations in ways that (we hope) makes clear what happens deep down but that takes time. These might change depending on feedback. We assume that we're in granular mode:

```
\setupalign[granular]
```

We will explain below what that means, but let us already now make clear that this will likely become the default! As far as we can see, due to the larger solution space, the inter-word spacing is more even but that also means that some paragraphs can become one line less or more.

## 21.3 Constructing paragraphs

There are several concepts at work when T<sub>E</sub>X breaks a paragraph into lines. Here we assume that we talk about text: words separated by spaces. We also assume that the text starts at the left edge and nicely runs till the right edge, with the exception of the last line.

- The spaces between words can stretch or shrink. We don't want that to be too inconsistent (visible) between two lines. This is where the terms loose and tight come into play.
- Words can be hyphenated but we don't want that to happen too often. We also discourage neighboring lines to have hyphens. Hyphenating the (pre) final line is also sort of bad.
- We definitely don't want words to stick out in the margin. If we have to choose, stretching is preferred over shrinking. If spaces become too small words, start to blur.
- If needed glyphs can stretch or shrink a little in order to get rid of excessive spacing. But we really want to keep it minimal, and avoid it when possible. Usually we permit more stretch than shrink. Not all scripts (and fonts for that matter) might work well with this feature.
- As a last resort we can stretch spaces so that we get rid of any still sticking out word. When T<sub>E</sub>X reports an overfull box (often a line) you have to pay attention!

When T<sub>E</sub>X decides where to break and when to finish doing so it uses a system of penalties and demerits and at some point makes decisions with regards to how bad a break-point (and eventually a paragraph) is. The penalties are normally relatively small unless we really want to penalize. When T<sub>E</sub>X is in the process of breaking a paragraph it calculates badness values for each line. This can be seen as a measure on how bad looking a line is; a badness of zero is good, but the larger the badness becomes, the worse the line is.

Here we shortly summarize the parameters that play a role in calculating what T<sub>E</sub>X calls the costs of breaking a line at some point: it's a combination of weighting penalties as well as over- or undershooting the line with, where the amount (dimension) and kind of (fillers) stretch and shrink determine the final verdict.

```
\ruledhbox to 20 ts{left \hss right}
\ruledhbox to 40 ts{left \hss right}
\ruledhbox to 5 ts{left \hss right}
\ruledhbox to 5 ts{left      right}
\ruledhbox to 5 es{%
  left
  \hskip 1ts plus 0.5ts\relax
  middle
  \hskip 1ts plus 1.5ts\relax
  right%
}
```

These boxes show a bit what happens with spacing that can stretch or shrink. The first three cases are not bad because it's what we ask for with the wildcard `\hss`.<sup>29</sup>

```
left ..... right
left ..... right
left right
left right
left ..... middle ..... right
```

T<sub>E</sub>X will run over each paragraph at most three times. On each such run, it will choose different breakpoints, calculate badness of each possible line, combine that with eventual penalties, and calculate a certain demerit value for each possible paragraph. It creates a set of solutions as it progresses, discards the worse cases so far and eventually ends on what it thinks is best.

<sup>29</sup> We use this opportunity to promote the new `ts` and `es` units.

The process is primarily controlled by these parameters:

- `\pretolerance`: This number determines the success of the first, not hyphenated pass. Often the value is set to the plain `TEX` value of 100. If `TEX` finds a possible division of a paragraph such that no line has a badness higher than `\pretolerance`, the algorithm quits here and that line is chosen.
- `\tolerance`: This number determines the success of the second, hyphenated pass. Often the value is set to the plain `TEX` value of 200.
- `\emergencystretch`: This dimension kicks in when the second pass is not successful. In `ConTEXt` we often set it to `2\bodyfontsize`.

When we are (in `ConTEXt` speak) tolerant, we have a value of 3000, while `verytolerant` bumps it to 4500. These are pretty large values compared to the default 100 and 200 that seem to cover most cases well, especially when we have short words, a reasonable width and lots of opportunities for hyphenation. Keep in mind that a macro package has to default to values that make sense for the average case.

We now come to the other relevant parameters. You need to keep in mind that the demerits are made from penalty values that get squared which is why parameters with demerits in their name have high values: a penalty of 50 squared has to relate to a demerit of 5000, so we might have  $2500 + 5000$  at some point.

The formula (most often) used to calculate the demerits  $d$  is

$$d = (l + b + p)^2 + e$$

Here  $l$  is the `\linepenalty`, set to 10 in plain,  $b$  is the badness of the line, and  $p$  is the penalty of the current break (for example, added by hyphenation, or by breaking an inline formula). The  $e$  stands for extra non-local demerits, that do not depend on only the current line, like the `\doublehyphendemerits` that is added if two lines in a row are hyphenated.

The badness reflects how the natural linewidth relates to the target width and uses a cubic function. A badness of zero is of course optimal, but a badness of 99 is pretty bad. A magic threshold is 12 (around that value a line is considered decent). If you look at the formula above you can now understand why the line penalty defaults to the low value of 10.



- `\hyphenpenalty`: When a breakpoint occurs at a discretionary this one gets added. In LuaMetaTeX we store penalties in the discretionary nodes but user defined `\discretionary`'s can carry dedicated penalties. This value is set to 50, which is not that much. Large values reduce the solution space so best keep this one reasonable.
- `\linepenalty`: Normally this is set to 10 and it is the baseline for a breakpoint. This is again a small value compared to for instance the penalties that you find in inline math. There we need some breakpoints and after binary and relation symbols such an opportunity is created. The specific penalties are normally 500 and 700. One has to keep in mind, as shown in the formula above, that the penalties are not acting on a linear scale when the demerits are calculated. Math spacing and penalty control is discussed in the (upcoming) math manual.
- `\doublehyphendemerits`: Because it is considered bad to have two hyphens in a row this is often set pretty high, many thousands. These are treated as demerits (so outside of the squared part of the above formula).
- `\finalhyphendemerits`: The final (pre last) line having a hyphen is also considered bad. The last line is handled differently anyway, just because it gets normally flushed left.
- `\adjdemerits`: lines get rated in terms of being loose, decent, tight, etc. When two lines have a different rating we bump the total demerits.
- `\looseness`: it is possible to force less or more lines but to what extent this request is honored depends on for instance the possible (emergency) stretch in the spaces (or any glue for that matter). `

It is worth noticing that you can set `\lastlinefit` such that the spaces in the last line will be comparable to those in the preceding line. This is a feature that  $\varepsilon$ -TeX brought us. Anyways, keep in mind normally penalties are either small, or when we want to be tough, pretty high. Demerits are often relatively large.

The next one is a flag that triggers expansion (or compression) of glyphs to kick in. Those get added to the available stretch and/or shrink of a line:

- `\adjustspacing`: Its value determines if expansion kicks in: glyphs basically get a stretch and shrink value, something that helps filling our lines. We only have zero, two and three (and not the pdfTeX value of two): three means 'only glyphs' and two means 'font kerns and glyphs'.

In LuaMetaTeX we also have:

- `\linebreakcriterion`: The normal distinction between loose, decent and tight in  $\TeX$  uses 12 for 0.5 and 99 for about 1.0, but because we have more granularity (.25) we can set four values instead. The default of zero ("0C0C0C63) then becomes "020C2A63. When set that way the default `\adjdemerits` has to be halved 5000 so that we compare the more granular distances. Don't worry if you 'don't get it', hardly any user will change these values. One can think of the 100 squared becomes a 10000 (at least this helps relating these numbers) and 10000 is pretty bad in  $\TeX$ s perception.
- `\adjustspacingstep`: When set this one is are used instead of the font bound value which permits local control without defining a new font instance.
- `\adjustspacingstretch`: idem.
- `\adjustspacingshrink`: idem.
- `\orphanpenalty`: This penalty will be injected before the last word of a paragraph.
- `\orphanpenalties`: Alternatively a series of penalties can be defined. This primitive expects a count followed by that number of penalties. These will be injected starting from the end.

The shape of a paragraph is determined by `\hangindent`, `\hangafter`, `\parshape` and `\parindent`. The width is controlled by `\hsize`, `\leftskip`, `\rightskip`. In addition there are `\parinitleftskip`, `\parinitrightskip`, `\parfillleftskip` and `\parfillrightskip` that control first and last lines.

We also have these:

- `\linebreakpasses`: When set to one, the currently set `\parpasses` will be applied.
- `\parpasses`: This primitive defined a set of sub passes that kick in when the second pass is finished. This basically opens up the par builder. It is still experimental and will be improved based upon user feedback. Although it is a side effect of improving the breaking of extensive mixes of math and text, it is also quite useful for text only (think novels).

In the next sections we will explain how these can improve the look and feel of what you typeset.

## 21.4 Subpasses

In  $\text{Lua}\TeX$  and therefore also in  $\text{LuaMeta}\TeX$  a paragraph is constructed in steps:

### Subpasses

- The list of nodes that makes the paragraph is hyphenated: words become a mixture of glyphs and discretionaries.
- That list is processed by a font handler that can remove, add or change glyphs depending on how glyphs interact. This depends on the language and scripts used.
- The result is fed into the par builder that applies up to three passes as mentioned before.

In traditional T<sub>E</sub>X these three actions are combined into one and the overhead is shared. In the split case the processing time gets distributed and in practice the last action is not the one that takes most time. This is why the mechanism that we discuss next has little impact on the run: calling the par builder a few times more seldom results in more runtime. This is why in we support so called sub passes between the second and third one.

Here is an example of a setup. We set a low tolerance for the first pass and second pass. We can do that because we don't need to play safe nor need to compromise.

```

\pretolerance 75
\tolerance 150
\parpasses 3
  threshold 0.025pt
  classes \indecentparpassclasses
  tolerance 150
next
  threshold 0.025pt
  classes \indecentparpassclasses
  tolerance 200
  emergencystretch .25\bodyfontsize
next
  threshold 0.025pt
  classes \indecentparpassclasses
  tolerance 200
  optional 1
  emergencystretch .5\bodyfontsize
\relax
\linebreakpasses 1

```

Because we want to retain performance we need to test efficiently if we really need the (here upto three) additional passes, so let's see how it is done. When a pass list is defined, and line break passes are enabled, the engine will check *after* the second

### Subpasses

pass if some more work is needed. For that it will do a quick analysis and calculate four values:

- overflow : the maximum value found, this is something really bad.
- underflow : the maximum value found, this is something we can live with.
- verdict : what is the worst badness of lines in this paragraph.
- classified : what classes are assigned to lines, think looseness, decent and tight.

There are two cases where the engine will continue with the applying passes: there is an overflow or there is a verdict (max badness) larger than zero. When we tested this on some large documents we noticed that this is nearly always true, but by checking we save a few unnecessary passes.

Next we test if a pass is really needed, and if not we check the next pass. When a pass is done, we pick up where we left, but we test for the overflow or badness every sub pass. The next checks make us run a pass:

- overfull exceeds threshold
- verdict exceeds badness
- classified overlaps classes

Here threshold, badness and classes are options in a pass section. Which test makes sense depends a bit on how T<sub>E</sub>X sees the result. Internally T<sub>E</sub>X uses numbers for its classification (0..5) but we map that onto a bitset because we want an overview:

	<b>indecent</b>	<b>almostdecent</b>	<b>loose</b>	<b>tight</b>
1 <b>veryloose</b>	+	+	+	
2 <b>loose</b>	+	+	+	
4 <b>semiloose</b>	+		+	
8 <b>decent</b>				
16 <b>semitight</b>	+			+
32 <b>tight</b>	+	+		+

The semiloose and semitight values are something LuaMetaT<sub>E</sub>X. In ConT<sub>E</sub>Xt we have these four variants predefined as `\indecentparpassclasses` and such.

The sections in a par pass setup are separated by `next`. For testing purposes you can add `skip` and `quit`. The `threshold` tests against the overfull value, the `badness` against the verdict and `classes` checks for overlap with encountered classes, the classification.

You can specify an identifier in the first segment that then will be used in tracing but it is also passed to callbacks that relate to this feature. Discussing these callback is outside the scope fo this wrapup.

## Subpasses

You need to keep in mind that parameters are not reset to their original values between two subpasses of a paragraph. We have `tolerance` and `emergencystretch` which are handy for simple setups. When we start with a small tolerance we often need to bump that one. The stretch is likely a last resort. The usual demerits can be set too: `doublehyphenemerits`, `finalhyphenemerits` and `adjdemerits`. We have `extrahyphenpenalty` that gets added to the penalty in a discretionary. You can also set `linepenalty` to a different value than it normally gets.

The looseness can be set but keep in mind that this only makes sense in very special cases. It's hard to be loose when there is not much stretch or shrink available. The `linebreakcriterion` parameter can best be left untouched and is mostly there for testing purposes.

The LuaMetaTeX specific `orphanpenalty` gets injected before the last word in a paragraph. High values can lead to overfull boxes but when used in text that hyphenate well or with languages that have short words it might work out well.

The next four parameters are related to expansion: `adjustspacing`, `adjustspacing-step`, `adjustspacingshrink` and `adjustspacingstretch`. Here we have several scenarios.

- Fonts are set up for expansion (in ConTeXt for instance with the quality specifier). When `hz` is then enabled it will always kick in.
- When we don't enable it, the par pass can do it by setting `adjustspacing` (to 3).
- When the other parameters are set these will overload the ones in the font, but used with the factors in there, so different characters get scaled differently. You can set the step to one to get more granular results.
- When expansion is *not* set on the font, setting the options in a pass will activate expansion but with the factors set to 1000. This means all characters are treated equal, which is less subtle.

When a font is not set up to use expansion, you can do something like this:

```
\parpasses      6
  classes          \indecentparpassclasses
  threshold        0.025pt
  tolerance        250
  extrahyphenpenalty  50
  orphanpenalty    5000
% font driven
```

## Subpasses

```

next ifadjustspacing
  threshold      0.025pt
  classes        \tightparpassclasses
  tolerance      300
  adjustspacing  3
  orphanpenalty 5000
next ifadjustspacing
  threshold      0.025pt
  tolerance      350
  adjustspacing  3
  adjustspacingstep  1
  adjustspacingshrink  20
  adjustspacingstretch 40
  orphanpenalty 5000
  emergencystretch .25\bodyfontsize
% otherwise, factors 1000
next
  threshold      0.025pt
  classes        \tightparpassclasses
  tolerance      300
  adjustspacing  3
  adjustspacingstep  1
  adjustspacingshrink  10
  adjustspacingstretch 15
  orphanpenalty 5000
next
  threshold      0.025pt
  tolerance      350
  adjustspacing  3
  adjustspacingstep  1
  adjustspacingshrink  20
  adjustspacingstretch 40
  orphanpenalty 5000
  emergencystretch .25\bodyfontsize
% whatever
next
  threshold      0.025pt
  tolerance      3000
  orphanpenalty 5000
  emergencystretch .25\bodyfontsize

```

```
\relax
```

With `ifadjustspacing` you ignore steps that expect the font to be setup, so you don't waste time if that is not the case.

There is also a callback parameter but that one is experimental and used for special purposes and testing. We don't expect users to mess with that.

A really special feature is optional content. Here we use as example a quote from Digital Typography:

```
Many readers will skim over formulas on their first reading of your
exposition. Therefore, your sentences should flow smoothly when all but
the simplest formulas are replaced by \quotation {blah} or some other
\optionalword {1} {grunting }noise.
```

Here the grunting (with embedded space) is considered optional. When you set `\linebreakoptional` to 1 this word will be typeset. However, when you set the pass parameter `linebreakoptional` to 0 it will be skipped. There can be multiple optional words with different numbers. The numbers are actually bits in a bit set so plenty is possible. However, normally these two values are enough, if used at all.

## 21.5 Definitions

The description above is rather low level and in practice users will use a bit higher level interface. Also, in practice only a subset of the parameters makes sense in general usage. It is not that easy to decide on what parameter subset will work out well but it can be fun to play with variants. After all, this is also what  $\TeX$  is about: look, feel and fun.

Some users praise the ability of recent  $\TeX$  engines to provide expansion and protrusion. This feature is a bit demanding because not only does it add to runtime (although in  $\text{LuaMeta}\TeX$  that normally can be neglected), it also makes the output files larger. Some find it hard to admit, but it even can result in bad looking documents when applied with extremes.

The traditional (MkIV) way to set up expansion is to add this to the top of the document, or at least before fonts get loaded.

```
\scratchcounter 1
\bgroup
\advance\scratchcounter 1
```

```
\egroup
\the\scratchcounter
```

and later on to enable it with:

```
\setupalign[hz]
```

However, par passes make it possible to be more selective. Take the following two definitions:

```
\startsetups align:pass:quality:1
  \pretolerance 50
  \tolerance 150
  \parpasses 6
    identifier \parpassidentifier{quality:1}
    threshold 0.025pt
    tolerance 175
  next
    threshold 0.025pt
    tolerance 200
  next
    threshold 0.025pt
    tolerance 250
  next
    classes \almostdecentparpassclasses
    tolerance 300
    emergencystretch .25\bodyfontsize
  next ifadjustspacing
    classes \indecentparpassclasses
    tolerance 300
    adjustspacing 3
    emergencystretch .25\bodyfontsize
  next
    threshold 0.025pt
    tolerance 3000
    emergencystretch 2\bodyfontsize
  \relax
\stopsetups

\startsetups align:pass:quality:2
  \pretolerance 50
```



```

\tolerance      150
\parpasses      5
  identifier      \parpassidentifier{quality:2}
  threshold      0.025pt
  tolerance      175
next
  threshold      0.025pt
  tolerance      200
next
  threshold      0.025pt
  tolerance      250
next ifadjustspacing
  classes        \indecentparpassclasses
  tolerance      300
  adjustspacing  3
  emergencystretch .25\bodyfontsize
next
  threshold      0.025pt
  tolerance      3000
  emergencystretch 2\bodyfontsize
\relax
\stopsetups

```

You can now enable one of these:

```
\setupalignpass[quality:1]
```

The result is shown in figure 21.1 where you can see that expansion is applied selectively; you have to zoom in to see where.

## 21.6 Tracing

There are several ways to see what goes on. The engine has a tracing option that is set with `\tracingpasses`. Setting it to 1 reports the passes on the console, and a value of 2 also gives some details.

There is also a tracker, `paragraphs.passes` that can be enabled. This gives a bit more information:

```

\enabletrackers[paragraphs.passes]
\enabletrackers[paragraphs.passes=summary]

```



```
\enabletrackers[paragraphs.passes=details]
```

If you want to see where expansion kicks in, you can use:

```
\showmakeup[expansion]
```

This is just one of the options, spaces, penalties, glue are are useful when you play with passes, but if you are really into the low level details, this is what you want:

```
\startnarrower[5*right]
```

```
\startshowbreakpoints[option=margin,offset=\dimexpr{.5\emwidth-\rightskip}]
```

```
\samplefile{tufte}
```

```
\stopshowbreakpoints
```

```
\stopnarrower
```

We thrive in information-thick worlds because of our marvelous **and** everyday capacity to select, edit, single out, structure, highlight, group, pair, merge, harmonize, synthesize, focus, organize, condense, reduce, boil down, choose, categorize, catalog, classify, list, abstract, scan, look into, idealize, isolate, discriminate, distinguish, screen, pigeonhole, pick over, sort, integrate, blend, inspect, filter, lump, skip, smooth, chunk, average, approximate, cluster, aggregate, outline, summarize, itemize, review, dip into, flip through, browse, glance into, leaf through, skim, refine, enumerate, glean, synopsisize, winnow the wheat from the chaff, and separate the sheep from the goats.

You can see the chosen solutions with

```
\showbreakpoints[n=1]
```

```

1 1 1 0 187 48809 veryloose glue
2 1 2 1 0 58909 decent glue
3 1 3 2 4 59105 decent glue
4 1 4 3 0 59205 decent glue
5 2 5 4 6 59461 decent glue
6 1 6 5 43 62270 loose glue
  1 7 5 17 62690 tight disc
7 2 8 6 9 62631 decent glue
  9 7 0 75290 decent disc
8 1 10 8 195 64152 loose glue
  1 11 8 2 65615 decent disc
9 12 10 23 65241 loose glue
  13 11 0 65715 decent glue

12 10 8 6 5 4 3 2 1
13 11 8 6 5 4 3 2 1

```

```

pass      : 3  demerits : 65341
subpass   : T  looseness : 0
subpasses : 0

```

When we started playing with the par builder in the perspective of math, we side tracked

and ended up with a feature that can be used in controlled situations. Currently we only have a low level ConT<sub>E</sub>Xt interface for this (see figure 21.2).

<p>The Earth, as a habitat for animal life, is in old age and has a fatal illness. Several, in fact. It would be happening whether humans had ever evolved or not. But our presence is like the effect of an old-age patient who smokes many packs of cigarettes per day—and we humans are the cigarettes.</p> <p><code>\tracinglousiness 1 \lousiness 0</code></p>	<p>The Earth, as a habitat for animal life, is in old age and has a fatal illness. Several, in fact. It would be happening whether humans had ever evolved or not. But our presence is like the effect of an old-age patient who smokes many packs of cigarettes per day—and we humans are the cigarettes.</p> <p><code>\lousiness 1 11 0</code></p>	<p>The Earth, as a habitat for animal life, is in old age and has a fatal illness. Several, in fact. It would be happening whether humans had ever evolved or not. But our presence is like the effect of an old-age patient who smokes many packs of cigarettes per day—and we humans are the cigarettes.</p> <p><code>\silliness 11</code></p>
---	--	--

**Figure 21.2** Influencing the way T<sub>E</sub>X breaks lines applied to `ward.tex`.

## 21.7 Criterion

The granular alignment option will configure the `linebreakcriterion` to work with 0.25 steps instead of 0.50 steps which means that successive lines can become a bit closer in spacing. There is no real impact on performance because testing happens anyway. In figure 21.3 you see some examples, where in some it indeed makes a difference.

## 21.8 Examples

*The ConT<sub>E</sub>Xt distribution comes with a few test setups: `spac-imp-tests.mkxl`. Once we have found a suitable set of values and sample texts we might discuss them here.*

*Currently we provide the following predefined passes that you can enable with `\setupalignpass`: `decent`, `quality`, `test1`, `test2`, `test3`, `test4`, `test5`. We hope that users are willing to test these.*

## 21.9 Pages

While the par builder does multiple passes, the page builder is a single pass progressive routine. Every time something gets added to the (so called) main vertical list the page state gets updated and when the page overflows what has been collected gets passed to the output routine. It is to a large extent driven by glue (with stretch and shrink) and penalties and when content (boxes) is added the process is somewhat complicated by inserts as these needs to be taken into account too.

You can get pages that run from top to bottom by adding stretch between lines but by default in ConT<sub>E</sub>Xt we prefer to fill up the bottom with white space.

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**Figure 21.3** More granular interline criteria.

It can be hard to make decisions at the  $\text{T}_{\text{E}}\text{X}$  end around a potential page break because in order to get an idea how much space is left, one needs to trigger the page builder which can have side effects.

Penalties play an important role and because these are used to control for instance widows and clubs high values can lead to underfull pages so if we want to influence that we need to cheat. For this we have three experimental mechanisms:

- tweaking the page goal: `\pageextragoal`
- initializing the state quantities: `\initialpageskip`
- adapting the state quantities as we go: `\additionalpageskip`

The first tweak is for me to play with, and when a widow or club is seen the extra amount can kick in. This feature is likely to be replaced by a more configurable one.

The second tweak lets the empty page start out with some given height, stretch and shrink. This variable is persistent over pages. This is not true for the third tweak: it

kicks in when the page gets initialized or as we go, but after it has been applied the value is reset. That makes it a feature like \looseness. We could combine these into one (because one can set up a persistent one in the macro package at well defined spots) but having an initial one also nicely can compensate the usual topskip glue hackery with a more natural control option.

Adapting the layout (within the regular text area) is done with \setpage\looseness an demonstrated in figure 21.4 and figure 21.5. Possible parameters are lines, height, stretch and shrink. You can also directly specify the number of lines. The other two features are not (yet) interfaced.

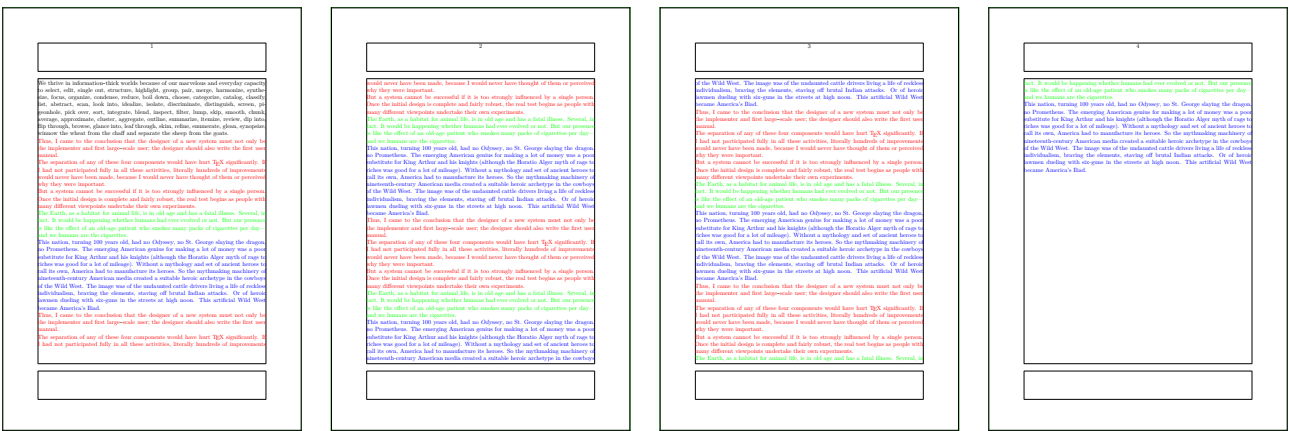


Figure 21.4 Cheating with page dimensions: [`\lines=2`].

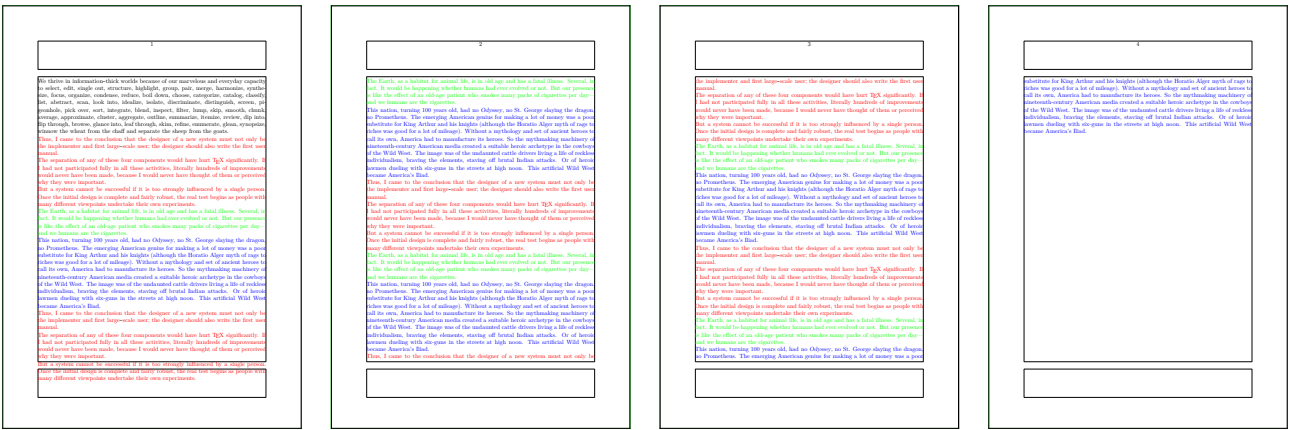


Figure 21.5 Cheating with page dimensions: [`-3`].

It is not that trivial to fulfill the wide range of user demands but over time the \setalign commands has gotten plenty of features. Getting for instance windows and clubs right in the kind of mixed usage that is common in ConTeXt is not always easy. One can experiment with scenarios (also to get some understanding of matters) but none is probably perfect (unless one does something close to manual tweaking). There is also the butterfly effect: a change here might trigger an issue there.

The examples in figure 21.6, 21.7 and 21.8 scale vertically in order to fill up the text area; the vZ parameter is set with setuplayout. In the example the widow and club penalties are set to 10000. In these examples we have enabled the layout.vz trackers that shows a small black rule indicating the amount of stretch.

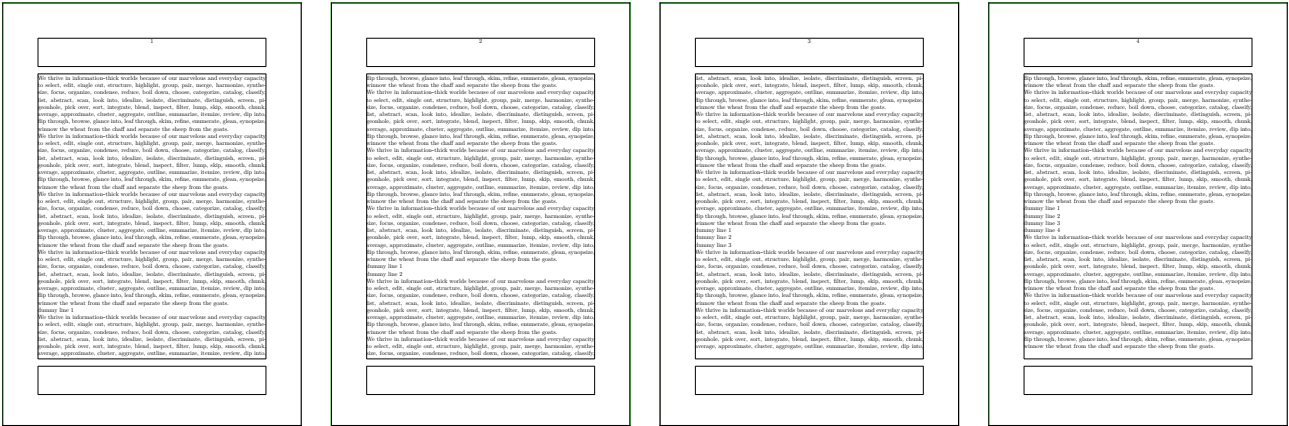


Figure 21.6 Cheating with vertical expansion: [vz=no].

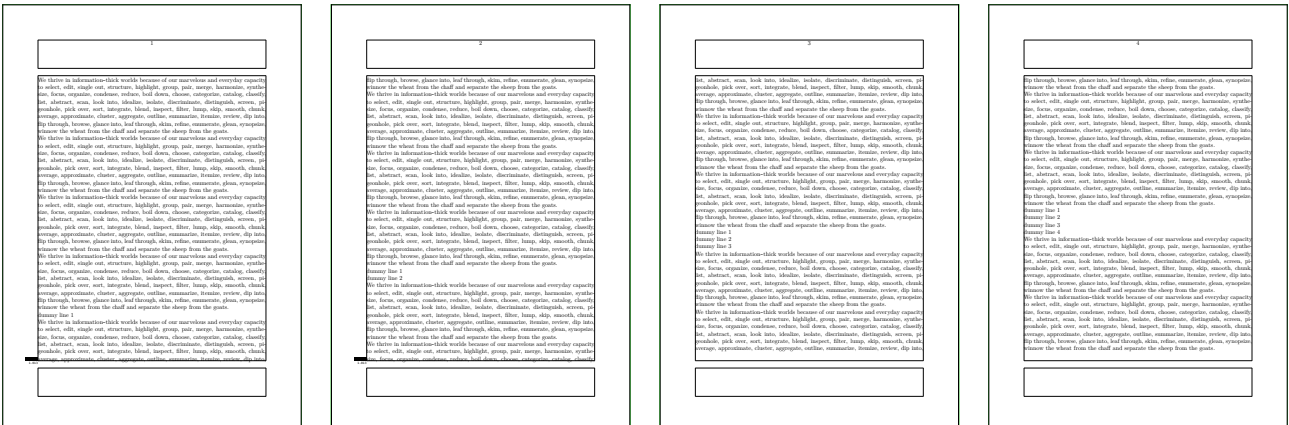


Figure 21.7 Cheating with vertical expansion: [vz=yes].

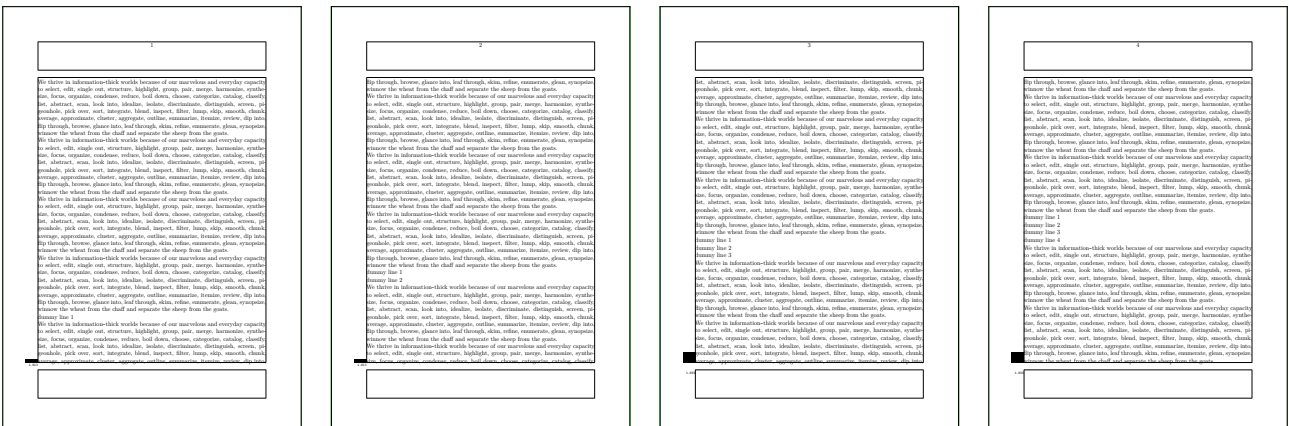


Figure 21.8 Cheating with vertical expansion: [vz=2].

There are a few other tweaks but these one can wonder about these. We can add stretch and shrink to the baseline skip, something that can also be triggered with the 'spread' option to `\setupalign`, assuming that also height is given). An alternative is to permit an extra line and accept a visual overflow, assuming that the layout is set up to make sure that the footer line doesn't overlap. None of this guarantees that a whole document with plenty of graphics and special constructs will come out well, but for text only it might work okay. Figures 21.9, 21.10 and 21.11 show some of this.



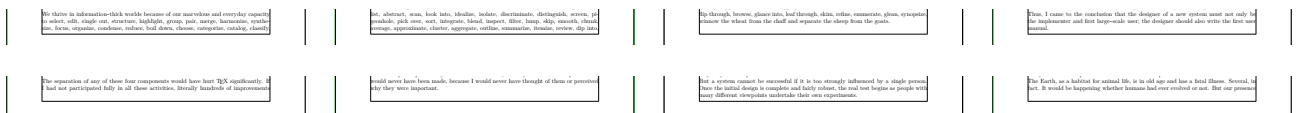


Figure 21.9 Cheating: just high penalties.

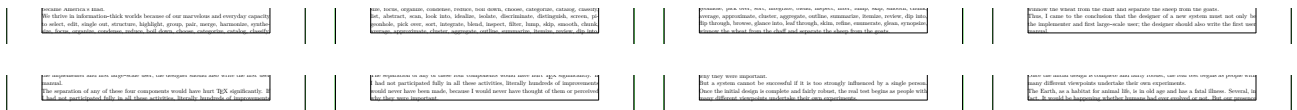


Figure 21.10 Cheating: \baselineskip 1\baselineskip plus 1pt minus .1pt.

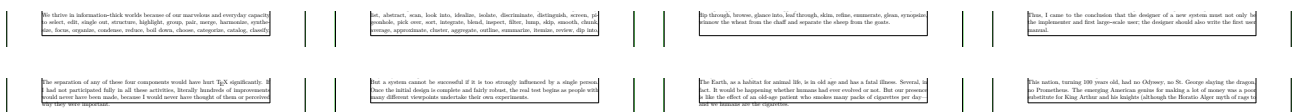


Figure 21.11 Cheating: \pageextragoal\lineheight.

## 21.10 Profiles

You can have a paragraph with lines that exceed the maximum height and/or depth or where spaces end up in a way that create so called rivers. Rivers are more a curiosity than an annoyance because any attempt to avoid them is likely to result in a worse looking result. The unequal line distances can be annoying too but these can be avoided when bringing lines closer together doesn't lead to clashes. This can be done without reformatting the paragraph by passing the profile option to `\setupalign`. It comes at the cost of a little more runtime and (as far as we observed) it kicks in seldom, for instance when inline math is used that has super- or subscripts, radicals, fractions or other slightly higher constructs.

## 21.10 Colofon

Author            Hans Hagen & Mikael Sundqvist  
ConT<sub>E</sub>Xt        2025.07.04 21:26  
LuaMetaT<sub>E</sub>X    2.11.07 | 20250705  
Support         www.pragma-ade.com  
                   contextgarden.net  
                   ntg-context@ntg.nl

## 22 Debugging

low level

TEX

debugging

## Contents

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### 22.1 Introduction

Below there will be some examples of how you can see what  $\text{T}_{\text{E}}\text{X}$  is doing. We start with some verbose logging but then move on to the more visual features. We occasionally point to some features present in the LuaMeta $\text{T}_{\text{E}}\text{X}$  engine. More details about what is possible can be found in documents in the Con $\text{T}_{\text{E}}\text{X}$ t distribution, for instance the ‘lowlevel’ manuals.

Typesetting involves par building, page building, inserts (footnotes, floats), vertical adjusters (stuff before and after the current line), marks (used for running headers and footers), alignments (to build tables), math, local boxes (left and right of lines), hyphenation, font handling, and more and each has its own specific ways of tracing, either provided by the engine, or by Con $\text{T}_{\text{E}}\text{X}$ t itself. You can run `context --trackers` to get a list of what Con $\text{T}_{\text{E}}\text{X}$ t can do, as it lists most of them. But we start with the language, where tokens play an important role.

### 22.2 Token lists

There are two main types of linked lists in  $\text{T}_{\text{E}}\text{X}$ : token lists and node lists. Token lists relate to the language and node lists collect (to be) typeset content and are used for several stack based structures. Both are efficiently memory managed by the engine. Token lists have only forward links, but node lists link in both directions, at least in Lua $\text{T}_{\text{E}}\text{X}$  and LuaMeta $\text{T}_{\text{E}}\text{X}$ .

When you define a macro, like the following, you get a token list:

```
\def\test#1{\bgroup\bf#1\egroup}
```

Internally the `\test` macro has carry the argument part and the body, and each is encoded as a number plus a pointer to the next token.

---

**control sequence: test**

---

9600	19	49	match	argument 1
601693	20	0	end match	
594835	1	123	left brace	bgroup
597786	143	0	protected call	bf
594054	21	1	parameter reference	
597573	2	125	right brace	egroup

---

Here the first (large) number is a memory location that holds two 4 byte integers per token: the so called info part codes the command and sub command, the two smaller numbers in the table, and a link part that points to the next memory location, here the next row. The last columns provide details. A character like ‘a’ is one token, but a control sequence like `\foo` is also one token because every control sequence gets a number. So, both take eight bytes of memory which is why a format file can become large and memory consumption grows the more macros you use.

In the body of the above `\test` macro we used `\bf` so let's see how that looks:

---

**permanent protected control sequence: bf**

---

604	137	24	if test	ifmmode
605	131	0	expand after	expandafter
606	143	0	protected call	mathbf
607	137	3	if test	else
608	131	0	expand after	expandafter
609	143	0	protected call	normalbf
610	137	2	if test	fi

---

Here the numbers are much lower which is an indication that they are likely in the format. They are also ordered, which is a side effect of LuaMetaTeX making sure that the token lists stored in the format file keep their tokens close together in memory which could potentially be a bit faster. But, when we are in a production run, the tokens come from the pool of freed or additionally allocated tokens:

```
\tolerant\permanent\protected\def\test[#1]#:#2%
```

```
{\iftok{#1}{sl}\bs\else\bf\fi#2}}
```

Gives us:

---

**permanent tolerant protected control sequence: test**

---

620151	12	91	other char	[	U+0005B	
50130	19	49	match			argument 1
611550	12	93	other char	]	U+0005D	
620147	19	58	match			argument :
601902	19	50	match			argument 2
599394	20	0	end match			
<hr/>						
599880	1	123	left brace			
599478	137	29	if test			iftok
615712	1	123	left brace			
489419	21	1	parameter reference			
611548	2	125	right brace			
599132	1	123	left brace			
579682	11	115	letter	s	U+00073	
620210	11	108	letter	l	U+0006C	
600896	2	125	right brace			
593965	143	0	protected call			bs
603764	137	3	if test			else
603803	143	0	protected call			bf
22619	137	2	if test			fi
611702	21	2	parameter reference			
598668	2	125	right brace			

---

If you are familiar with  $\TeX$  and spend some time looking at this you will start recognizing entries. For instance 11 115 translates to letter s because 11 is the so called command code of letters (also its `\catcode`) and the s has utf8 value 115. The LuaMeta $\TeX$  specific `\iftok` conditional has command code 135 and sub code 29. Internally these are called cmd and chr codes because in many cases it's characters that are the sub commands.

There is more to tell about these commands and the way macros are defined, for instance tolerant here means that we can omit the the first argument (between brackets) in which case we pick up after the `#:`. With protected we indicate that the macro will not expand in for instance an `\edef` and permanent marks the macro as one that a user cannot redefine (assuming that overload protection is enabled). The extended macro argument parsing features and macro overload protection are something specific to LuaMeta $\TeX$ .

These introspective tables can be generated with:

```
\luatokenable\test
```

after loading the module `system-tokens`. The reason for having a module and not a built-in tracer is that users seldom want to do this. Instead they might use `\showluatokens\test` that just reports something similar to the console and/or log file.

There is much more to tell but most users have no need to look into these details unless they are curious about what  $\TeX$  does. In that case using `tracingall` and inspecting the log file can be revealing too, but be prepared for huge files. In `LuaMeta $\TeX$`  we have tried to improve these traces a bit but that's of course subjective and even then logs can become huge. But even if one doesn't understand all that is shown, it gives an impression about how much work  $\TeX$  is actually doing.

## 22.3 Node lists

A node list is what you get from input that is (to be) typeset. There are several ways to see what node lists are produced but these are all very verbose. Take for instance:

```
\setbox\scratchbox\hbox{test \bf test}
```

```
\showboxhere\scratchbox
```

This gives us:

```
\hlist[box][color=1,colormodel=1,mathintervals=1], width 47.8457pt, height 7.48193pt, depth
0.15576pt, direction l2r, state 1
.\list
..\glyph[unset][color=1,colormodel=1], protected, wd 4.42041pt, ht 7.48193pt, dp 0.15576pt, language
(n=1,l=2,r=3), hyphenationmode "79F3F, options "80, font <1: DejaVuSerif @ 11.0pt>, glyph U+0074
..\glyph[unset][color=1,colormodel=1], protected, wd 6.50977pt, ht 5.86523pt, dp 0.15576pt, language
(n=1,l=2,r=3), hyphenationmode "79F3F, options "80, font <1: DejaVuSerif @ 11.0pt>, glyph U+0065
..\glyph[unset][color=1,colormodel=1], protected, wd 5.64502pt, ht 5.86523pt, dp 0.15576pt, language
(n=1,l=2,r=3), hyphenationmode "79F3F, options "80, font <1: DejaVuSerif @ 11.0pt>, glyph U+0073
..\glyph[unset][color=1,colormodel=1], protected, wd 4.42041pt, ht 7.48193pt, dp 0.15576pt, language
(n=1,l=2,r=3), hyphenationmode "79F3F, options "80, font <1: DejaVuSerif @ 11.0pt>, glyph U+0074
..\glue[spaceskip][color=1,colormodel=1] 3.49658pt plus 1.74829pt minus 1.16553pt, font 1
..\glyph[unset][color=1,colormodel=1], protected, wd 5.08105pt, ht 7.48193pt, dp 0.15576pt, language
(n=1,l=2,r=3), hyphenationmode "79F3F, options "80, font <10: DejaVuSerif-Bold @ 11.0pt>, glyph
U+0074
..\glyph[unset][color=1,colormodel=1], protected, wd 6.99854pt, ht 5.86523pt, dp 0.15576pt, language
(n=1,l=2,r=3), hyphenationmode "79F3F, options "80, font <10: DejaVuSerif-Bold @ 11.0pt>, glyph
U+0065
```

### Node lists



```

..\glyph[unset][color=1,colormodel=1], protected, wd 6.19287pt, ht 5.86523pt, dp 0.15576pt, language
  (n=1,l=2,r=3), hyphenationmode "79F3F, options "80, font <10: DejaVuSerif-Bold @ 11.0pt>, glyph
  U+0073
..\glyph[unset][color=1,colormodel=1], protected, wd 5.08105pt, ht 7.48193pt, dp 0.15576pt, language
  (n=1,l=2,r=3), hyphenationmode "79F3F, options "80, font <10: DejaVuSerif-Bold @ 11.0pt>, glyph
  U+0074

```

The periods indicate the nesting level and the slash in front of the initial field is mostly a historic curiosity because there are no `\hlist` and `\glue` primitives, but actually there is in `LuaMetaTEX` a `\glyph` primitive but that one definitely doesn't want the shown arguments.

That said, here we have a horizontal list where the list field points to a glyph that itself points to a next one. The space became a glue node. In `LuaTEX` and even more in `LuaMetaTEX` all nodes have or get a subtype assigned that indicates what we're dealing with. This is shown between the first pair of brackets. Then there are attributes, between the second pair of brackets, which actually is also a (sparse) linked list. Here we have two attributes set, the color, where the number points to some stored color specification, and the (here somewhat redundant) color space. The names of these attributes are macro package dependent because attributes are just a combination of a number and value. The engine itself doesn't do anything with them; it is the Lua code you plug in that can do something useful based on the values.

It will be clear that watching a complete page, with many nested boxes, rules, glyphs, discretionaries, glues, kerns, penalties, boundaries etc quickly becomes a challenge which is why we have other means to see what we get so let's move on to that now.

## 22.4 Visual debugging

In the early days of `ConTEXt`, in the mid 90's of the previous century, one of the first presentations at an `ntg` meeting was about visual debugging. This feature was achieved by overloading the primitives that make boxes, add glue, inject penalties and kerns, etc. It actually worked quite well, although in some cases, for instance where boxes have to be unboxed, one has to disable it. I remember some puzzlement among the audience about the fact that indeed these primitives could be overloaded without too many side effects. It will be no surprise that this feature has been carried on to later versions, and in `ConTEXt MkIV` it was implemented in a different (less intrusive) way and it got gradually extended.

```
\showmakeup \hbox{test \bf test}
```

This gives us a framed horizontal box, with some text and a space glue:

### Visual debugging

# H test test

Of course not all information is well visible simply because it can be overlaid by what follows, but one gets the idea. Also, when you have a layer capable pdf viewer you can turn on and off categories, so you can decide to only show glue. You can also do that immediately, with `\showmakeup[glue]`.

There is a lot of granularity: `hbox`, `vbox`, `vtop`, `kern`, `glue`, `penalty`, `fontkern`, `strut`, `whatsit`, `glyph`, `simple`, `simplehbox`, `simplevbox`, `simplevtop`, `user`, `math`, `italic`, `origin`, `discretionary`, `expansion`, `line`, `space`, `depth`, `marginkern`, `mathkern`, `dir`, `par`, `mathglue`, `mark`, `insert`, `boundary`, the more selective `vkern`, `hkern`, `vglue`, `hglue`, `vpenalty` and `hpenalty`, as well as some presets like `boxes`, `makeup` and all.

When we have:

```
\showmakeup \framed[align=normal]{\samplefile{ward}}
```

we get:

And that is why exploring this with a layers enabled pdf viewer can be of help. Alternatively a more selective use of `\showmakeup` makes sense, like

```
\showmakeup[line,space] \framed[align=normal]{\samplefile{ward}}
```

Here we only see lines, regular spaces and spaces that are determined by the space factor that is driven by punctuation.

We can typeset the previous example with these settings:

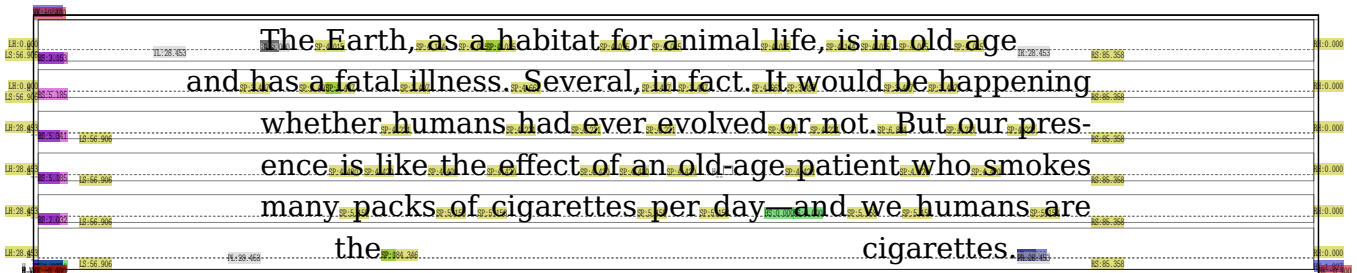
```
\leftskip      2cm
\rightskip     3cm
\hangindent    1cm
```

```

\hangafter      2
\parfillrightskip 1cm
\parfillleftskip 1cm % new
\parinitrightskip 1cm % new
\parinitleftskip 1cm % new
\parindent      2cm % different

```

This time we get:



Looking at this kind of output only makes sense on screen where you can zoom in but what we want to demonstrate here is that in LuaMetaTeX we have not only a bit more control over the paragraph (indicated by comments) but also that we always have the related glue present. The reason is that we then have a more predictable packaged line when we look at one from the Lua end. Where TeX normally moves the final line content left or right via either glue or the shifts property of a box, here we always use the glue. We call this normalization. Keep in mind that TeX was not designed (implemented) with exposing its internals in mind, but for LuaTeX and LuaMetaTeX we have to take care of that.

Another characteristic is that the paragraph stores these (and many more) properties in the so called initial par node so that they work well in situations where grouping would interfere with our objectives. As with all extensions, these are things that can be configured in detail but they are enabled in ConTeXt by default.

## 22.5 Math

Math is a good example where this kind of tracing helps development. Here is an example:

```
\im { \showmakeup y = \sqrt {2x + 4} }
```

Scaled up we get:

$$y = \sqrt{2x + 4}$$

Instead of showing everything we can again be more selective:

```
\im {
  \showmakeup[mathglue,glyph]
  y = \sqrt {2x + 4}
}
```

Here we not only limit ourselves to math glue, but also enable showing the bounding boxes of glyphs.

$$y = \sqrt{2x + 4}$$

This example also shows that in LuaMetaTeX we have more classes than in a traditional TeX engine. For instance, radicals have their own class as do digits. The radical class is an engine one, the digit class is a user defined class. You can set up the spacing between each class depending on the style. For the record: this is just one of the many extensions to the math engine and when extensions are being developed it helps to have this kind of tracing. Take for instance the next example, where we have multiple indexes (indicated by `__`) on a nucleus, that get separated by a little so called continuation spacing.

```
\im {
  \showmakeup[mathglue,glyph]
  y = \sqrt {x__1__a {\darkred +} x__1__b}
}
```

$$y = \sqrt{x_1 a + x_1 b}$$

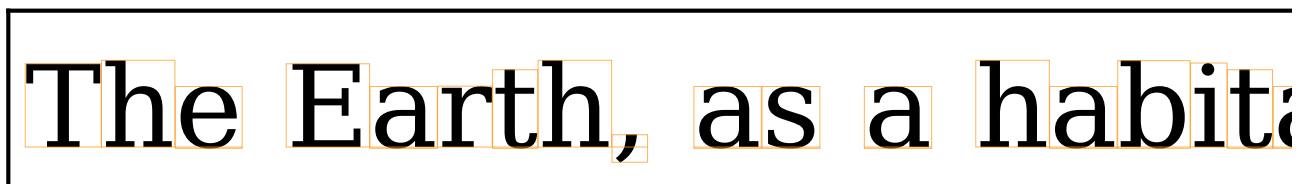
Here the variable class is used for alphabetic characters and some more, contrary to the more traditional (often engine assigned) ordinary class that is now used for the left-overs.

## 22.6 Fonts

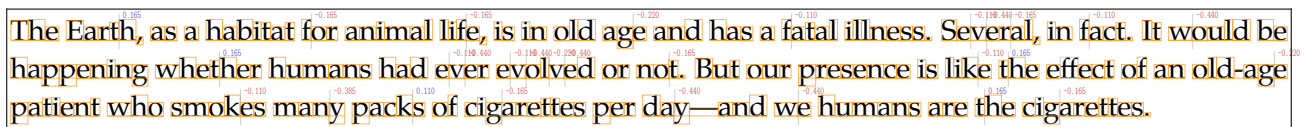
Some of the mentioned tracing has shortcuts, for instance `\showglyphs`. Here we show the same sample paragraph as before:

```
\showglyphs
\showfontkerns
\framed[align=normal]{\samplefile{ward}}
```

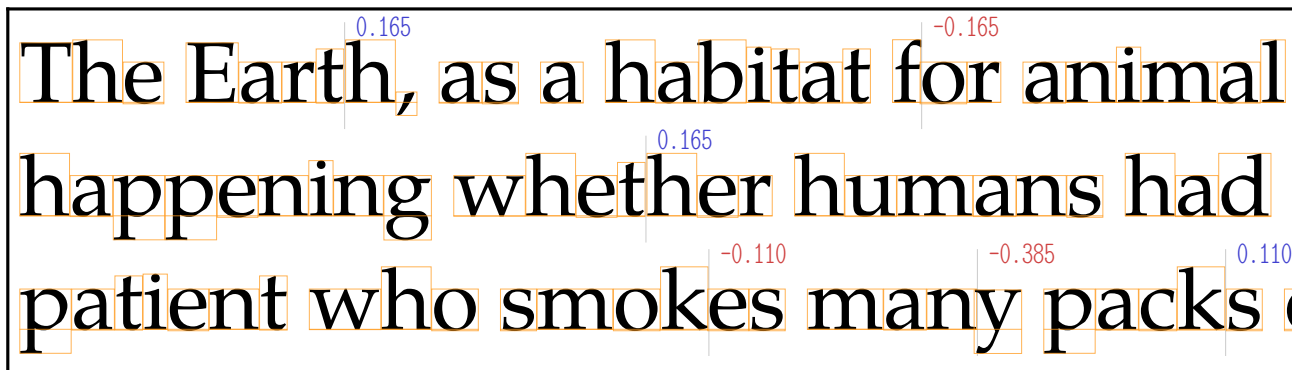
Here is the upper left corner of the result:



What font kerns we get depends on the font, here we use pagella:

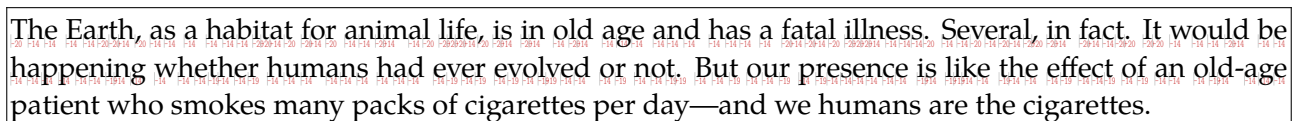


If we zoom in the kerns are more visible:



And here is another one:

```
\showfontexpansion
\framed[align={normal,hz}]{\samplefile{ward}}
```



or blown up:

The Earth, as a habitat for animal  
 happening whether humans had  
 patient who smokes many packs

The last line (normally) doesn't need expansion, unless we want it to be compatible with preceding lines, space-wise. So when we do this:

```
\showfontexpansion
\framed[align={normal,hz,fit}]{\samplefile{ward}}
```

the `fit` directives result in somewhat different results:

The Earth, as a habitat for animal  
 happening whether humans had  
 patient who smokes many packs

As with other visual tracers you can get some insight in how  $\text{T}_{\text{E}}\text{X}$  turns your input into a typeset result.

## 22.7 Overflow

By default the engine is a bit stressed to make paragraphs fit well. This means that we can get overflowing lines. Because there is a threshold only visible overflow is reported. If you want a visual clue, you can do this:

```
\enabletrackers[builders.hpack.overflow]
```

With:

```
\ruledvbox{\hsize 3cm test test test test test test test test}
```

We get:

```
test test test test
test test test test
```

The red bar indicates a potential problem. We can also get an underflow, as demonstrated here:

```
\ruledvbox {
  \setupalign[verytolerant,stretch]
  \hsize 3cm test test test test test test test test
}
```

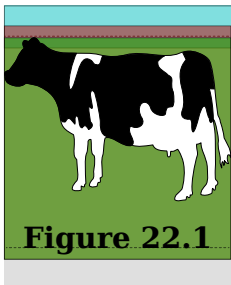
Now we get a blue bar that indicates that we have a bit more stretch than is considered optimal:

```
test test test
test test test
test test
```

Especially in automated flows it makes sense to increase the tolerance and permit stretch. Only when the strict attempt fails that will kick in.

## 22.8 Side floats

Some mechanisms are way more complex than a user might expect from the result. An example is the placement of float and especially side floats.



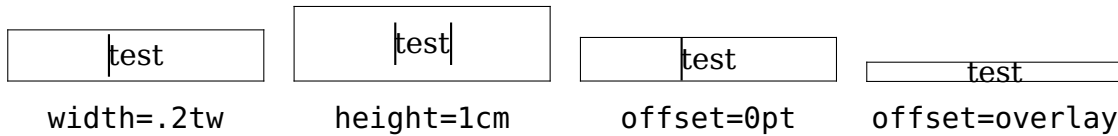
Not only do we have to make sure that the spacing before such a float is as good and consistent as possible, we also need the progression to work out well, that is: the number of lines that we need to indent.

For that we need to estimate the space needed, look at the amount of space before and after the float, check if it will fit and move to the next page if needed. That all involves dealing with interline spacing, interparagraph spacing, spacing at the top of a page, permitted slack at the bottom of page, the depth of the preceding lines, and so on. The tracer shows some of the corrections involved but leave it to the user to imagine what it relates to; the previous sentence gives some clues. This tracker is enabled with:

```
\enabletrackers[floats.anchoring]
```

## 22.9 Struts

We now come to one of the most important trackers, `\showstruts`, and a few examples shows why:



Here in all cases we've set the width to 20 percent of the text width (tw is an example of a plugged in dimension). In many places ConT<sub>E</sub>Xt adds struts in order to enforce proper spacing so when spacing is not what you expect, enabling this tracker can help you figure out why.

## 22.10 Features

Compared to the time when T<sub>E</sub>X showed up the current fonts are more complicated, especially because features go beyond only ligaturing and kerning. But even ligaturing can be different, because some fonts use kerning and replacement instead of a new character. Pagella uses a multiple to single replacement:

```
font      22: texgyrepagella-regular.otf @ 10.0pt

features  [basic: kern=yes, liga=yes, mark=yes, mkmk=yes, script=dflt]
          [extra: analyze=yes, autolanguage=position, autoscript=position,
          checkautoeffect=yes, checkmarks=yes, checkmissing=yes,
          compoundhyphen=yes, curs=yes, devanagari=yes, dummies=yes,
          expansion=quality, extensions=yes, extrafeatures=yes,
          extraprivates=yes, fixdot=yes, indic=auto, itlc=yes,
          mathrules=yes, mode=node, spacekern=yes,
          textcontrol=collapsehyphens, replaceapostrophe, visualspace=yes,
          wipemath=yes]

step 1    effe.fietsen U+65:e U+66:f [ pre: U+2D:␣ ] U+66:f U+65:e [glue] U+66:f
          U+69:i U+65:e U+74:f [ pre: U+2D:␣ ] U+73:s U+65:e U+6E:m

          feature 'liga', type 'gsub_ligature', lookup 's_s_9', replacing
          U+66 (f) upto U+66 (f) by ligature U+FB00 (f_f) case 2

step 2    effe.fietsen U+65:e [ pre: U+66:f U+2D:␣ post: U+66:f replace:
          U+FB00:ff ] U+65:e [glue] U+66:f U+69:i U+65:e U+74:f [ pre: U+2D:␣ ]
          U+73:s U+65:e U+6E:m
```



feature 'liga', type 'gsub\_ligature', lookup 's\_s\_10',  
replacing U+66 (f) upto U+69 (i) by ligature U+FB01 (f\_i)  
case 2

**step 3** `effe fietsen` U+65:e [ pre: U+66:f U+2D:␣ post: U+66:f replace:  
U+FB00:ff ] U+65:e [glue] U+FB01:ffi U+65:e U+74:t [ pre: U+2D:␣ ]  
U+73:s U+65:e U+6E:m

feature 'kern', type 'gpos\_pair', lookup 'p\_s\_0', inserting  
move -0.14992pt between U+66 (f) and U+65 (e) as  
postinjections

**result** `effe fietsen` U+65:e [ pre: U+66:f U+2D:␣ post: U+66:f [kern] replace:  
U+FB00:ff ] U+65:e [glue] U+FB01:ffi U+65:e U+74:t [ pre: U+2D:␣ ]  
U+73:s U+65:e U+6E:m

Not all features listed here are provided by the font (only the four character ones) because we're using  $\TeX$  which, it being  $\TeX$ , means that we have plenty more ways to mess around with additional features: it's all about detailed control. But what you see here are the steps taken: the font handler loops over the list of glyphs and here we see the intermediate results when something has changed. There can be way more loops than in this simple case.

With Cambria we get a single replacement combined with kerning:

**font** 23: cambria.ttc @ 10.0pt

**features** [basic: kern=yes, liga=yes, mark=yes, mkmk=yes, script=latn]  
[extra: analyze=yes, autolanguage=position, autoscript=position,  
checkautoeffect=yes, checkmarks=yes, checkmissing=yes,  
compoundhyphen=yes, curs=yes, devanagari=yes, dummies=yes,  
expansion=quality, extensions=yes, extrafeatures=yes,  
extraprivates=yes, fixdot=yes, indic=auto, itlc=yes,  
mathrules=yes, mode=node, spacekern=yes,  
textcontrol=collapsehyphens, replaceapostrophe, visualspace=yes,  
wipemath=yes]

**step 1** `effe fietsen` U+65:e U+66:f [ pre: U+2D:␣ ] U+66:f U+65:e [glue] U+66:f  
U+69:i U+65:e U+74:t [ pre: U+2D:␣ ] U+73:s U+65:e U+6E:m

feature 'liga', type 'gsub\_contextchain', chain lookup  
's\_s\_38', replacing single U+66 by U+F016C

**step 2** `effe fietsen` U+65:e U+66:f [ pre: U+2D:␣ ] U+66:f U+65:e [glue]  
U+F016C:f U+69:i U+65:e U+74:t [ pre: U+2D:␣ ] U+73:s U+65:e U+6E:m

feature 'kern', type 'gpos\_pair', merged lookup 'p\_s\_0',  
inserting move -0.12207pt between U+66 and U+65 as injections

**result**    `effe fietsen` U+65:e U+66:f [ pre: U+2D:␣ ] U+66:f [kern] U+65:e [glue]  
U+F016C:f U+69:i U+65:e U+74:t [ pre: U+2D:␣ ] U+73:s U+65:e U+6E:m

One complication is that hyphenation kicks in which means that whatever we do has to take the pre, post and replacement bits into account combined with what comes before and after. Especially for complex scripts this tracker can be illustrative but even then only for those who like to see what fonts do and/or when they add additional features runtime.

## 22.11 Profiling

There are some features in ConT<sub>E</sub>Xt that are nice but only useful in some situations. An example is profiling. It is something you will not turn on by default, if only because of the overhead it brings. The next two paragraphs (using Pagella) show the effect.

The command `\binom` is the standard notation for binomial coefficients and is preferred over `\choose`, which is an older macro that has limited compatibility with newer packages and font encodings:  $|A| = \binom{N}{k}^2$ . Additionally, `\binom` uses proper spacing and size for the binomial symbol. In conclusion, it is recommended to use `\binom` instead of `\choose` in T<sub>E</sub>X for typesetting binomial coefficients for better compatibility and uniform appearance.

The previous paragraph is what comes out by default, while the next one used these settings plus an additional `\enabletrackers[profiling.lines.show]`.

The command `\binom` is the standard notation for binomial coefficients and is preferred over `\choose`, which is an older macro that has limited compatibility with newer packages and font encodings:  $|A| = \binom{N}{k}^2$ . Additionally, `\binom` uses proper spacing and size for the binomial symbol. In conclusion, it is recommended to use `\binom` instead of `\choose` in T<sub>E</sub>X for typesetting binomial coefficients for better compatibility and uniform appearance.

This feature will bring lines together when there is no clash and is mostly of use when a lot of inline math is used. However, when this variant of profiling (we have an older one too) is enabled on a 300 page math book with thousands of formulas, only in a few places it demonstrated effect; it was hardly needed anyway. So, sometimes tracing shows what makes sense or not.

## 22.12 Par builder

Here is a sample paragraph from Knuths “Digital Typography”:

**15.** (This procedure maintains four integers  $(A, B, C, D)$  with the invariant meaning that “our remaining job is to output the continued fraction for  $(Ay + B)/(Cy + D)$ , where  $y$  is the input yet to come.”) Initially set  $j \leftarrow k \leftarrow 0$ ,  $(A, B, C, D) \leftarrow (a, b, c, d)$ ; then input  $x_j$  and set  $(A, B, C, D) \leftarrow (Ax_j + B, A, Cx_j + D, C)$ ,  $j \leftarrow j + 1$ , one or more times until  $C + D$  has the same sign as  $C$ . (When  $j > 1$  and the input has not terminated, we know that  $1 < y < \infty$ ; and when  $C + D$  has the same sign as  $C$  we know therefore that  $(Ay + B)/(Cy + D)$  lies between  $(A + B)/(C + D)$  and  $A/C$ .) Now comes the general step: If no integer lies strictly between  $(A + B)/(C + D)$  and  $A/C$ , output  $X_k \leftarrow \lfloor A/C \rfloor$  and set  $(A, B, C, D) \leftarrow (C, D, A - X_k C, B - X_k D)$ ,  $k \leftarrow k + 1$ ; otherwise input  $x_j$  and set  $(A, B, C, D) \leftarrow (Ax_j + B, A, Cx_j + D, C)$ ,  $j \leftarrow j + 1$ . The general step is repeated ad infinitum. However, if at any time the *final*  $x_j$  is input, the algorithm immediately switches gears: It outputs the continued fraction for  $(Ax_j + B)/(Cx_j + D)$ , using Euclid’s algorithm, and terminates.

There are indicators with tiny numbers that indicate the possible breakpoints and we can see what the verdict is:

1 2 1 0 10	400	decent	glue	1 26 19 28	259433	tight	penalty	1 51 47 99	9640	decent	glue
3 2 0 70	6400	tight	glue	6 2 27 21 73	7943	decent	glue	52 48 13	9545	loose	glue
2 2 3 1 69	6641	loose	glue	28 23 0	12610	decent	glue	53 47 13	9416	decent	glue
4 2 5	29125	decent	glue	2 29 21 0	8647	tight	glue	54 49 59	2729967	loose	glue
5 1 5	28084	tight	glue	30 25 39	1113011	loose	penalty	2 55 49 12	2725690	decent	glue
3 6 2 0	6500	decent	math	2 31 24 39	1112825	decent	penalty	2 56 49 0	2725306	decent	glue
1 7 2 41	9001	tight	glue	1 32 26 3	1362102	decent	penalty	57 49 66	2730982	tight	glue
3 2 8 3 2	6785	decent	glue	7 1 33 27 95	18968	loose	glue	11 1 58 51 8	9964	decent	glue
1 9 6 67	12429	loose	glue	1 34 29 3	8816	decent	glue	59 56 55	2729531	loose	glue
10 3 67	18490	tight	glue	35 27 3	10647	tight	glue	60 55 55	2725811	decent	glue
11 6 4	29100	decent	glue	2 36 29 15	18251	tight	glue	61 56 71	2725406	decent	glue
1 12 7 0	9101	decent	math	37 31 0	1112925	decent	glue	62 55 71	2733971	tight	glue
2 13 6 0	7589	tight	math	38 31 32	1114589	tight	glue	12 63 58 61	15005	tight	glue
4 1 14 8 1	6906	decent	glue	1 39 32 0	1362202	decent	glue				
1 15 9 90	12529	decent	glue	8 2 40 33 1	269089	decent	penalty	59 56 49 44 39 32 26 19 13 6 2			
1 16 8 90	12410	tight	glue	2 41 34 0	8916	decent	glue	60 55 49 44 39 32 26 19 13 6 2			
2 17 12 22	10125	loose	glue	42 36 41	18395	decent	glue	61 56 49 44 39 32 26 19 13 6 2			
18 13 0	497689	decent	penalty	43 36 25	18980	tight	glue	62 55 49 44 39 32 26 19 13 6 2			
2 19 13 10	7989	decent	math	1 44 39 92	1622606	tight	penalty	63 58 51 47 41 34 29 21 14 8 3 1			
5 20 15 22	13553	loose	glue	9 1 45 40 8	291913	decent	glue	pass : 1 demerits : 15105			
2 21 14 22	7747	tight	glue	46 40 0	269189	decent	math	subpass : P looseness : 0			
22 17 64	15601	loose	glue	2 47 41 10	9316	decent	glue	subpasses : 0			
1 23 16 64	12510	decent	glue	1 48 41 20	9016	decent	glue				
1 24 17 20	10225	decent	glue	4 49 44 0	2725206	decent	penalty				
1 25 19 1	8110	decent	glue	10 50 45 41	294514	loose	glue				

The last lines in the last column show the route that the result takes. Without going into details, here is what we did:

```
\startshowbreakpoints
  \samplefile{math-knuth-dt}
\stopshowbreakpoints

\showbreakpoints
```

This kind of tracing is part of a mechanism that makes it possible to influence the choice by choosing a specific preferred breakpoint but that is something the average user is unlikely to do. The main reason why we have this kind of trackers is that when developing the new multi-step par builder feature we wanted to see what exactly it did influence. That mechanism uses an LuaMetaT<sub>E</sub>X feature where we can plug in additional passes using the `\parpasses` primitive that can add different strategies that are tried until criteria for over- and underfull thresholds and/or badness are met. Each step can set the relevant parameters differently, including expansion, which actually makes for more efficient output and better runtime when that features is not needed to get better results.

## 22.13 More

There are many more visual trackers, for instance `layout.vz` for when you enabled vertical expansion, `typesetters.suspects` for identifying possible issues in the input like invisible spaces. Trackers like `nodes.destinations` and `nodes.references` will show the areas used by these mechanisms. There are also trackers for positions, (cjk and other), script handling, rubies, tagging, italic correction, breakpoints and so on. The examples in the previous sections illustrate what to expect and when to use a specific mechanism knowing this might trigger you to check if a tracker exists. Often the test suite has examples of usage.

## 22.13 Colofon

Author	Hans Hagen & Mikael Sundqvist
ConT <sub>E</sub> Xt	2025.07.04 21:26
LuaMetaT <sub>E</sub> X	2.11.07   20250705
Support	<a href="http://www.pragma-ade.com">www.pragma-ade.com</a> <a href="http://contextgarden.net">contextgarden.net</a> <a href="mailto:ntg-context@ntg.nl">ntg-context@ntg.nl</a>

**23 Pages**

low level

TEX

pages

## Contents

23.1 Introduction	360
23.2 Rows becoming columns	360

### 23.1 Introduction

There are several builder in the engine: paragraphs, math, alignments, boxes and if course pages. But where a paragraph is kind of complete and can be injected on a line by line basis, a page is less finished. When enough content is collected the result so far is handled over to the output routine. Calling it a routine is somewhat confusing because it is not really a routine, it's the token list `\output` that gets expanded and what in there is supposed to something with the result, like adding inserts (footnotes, moved around graphics aka floats, etc.), adding headers and footers, possibly using marks, and finally wrapping up and shipping out.

The engine primarily offers a single column page so two or more columns are done by using tricks, like typesetting on a double height and splitting the result. If columns need to be balanced some extra work has to be done, and it's definitely non trivial when we have more that just text.

In this chapter we will discuss and collect some mechanisms that deal with pages or operate at the outer vertical level. We might discuss some primitive but more likely you will see various solutions based on `TEX` macros and Lua magic.

*This is work in progress.*

### 23.2 Rows becoming columns

*This is an experimental mechanism. We need to check/decide how to deal with penalties. We also need to do more checking.*

Conceptually this is a bit strange feature but useful nevertheless. There are several multi-column mechanisms in `ConTEXt` and each is made for a specific kind of usage. You can, to some extent, consider `tabulate` to produce columns too, however it demands a bit of handy work. Say that you have this:

```
\starttabulate[|l|l|]
\NC 1 \NC one \NC \NR
\NC 2 \NC two \NC \NR
\NC 3 \NC three \NC \NR
```

```
\NC 4 \NC four \NC \NR
\NC 5 \NC five \NC \NR
\stoptabulate
```

but you don't want to waste space. So you might want:

```
1 one                4 four
2 two                5 five
3 three
```

or maybe even this:

```
1 one                3 three                5 five
2 two                4 four
```

but still wants to code like this:

```
\starttabulate[|l|l|]
\NC 1 \NC one \NC \NR
\NC 2 \NC two \NC \NR
\NC 3 \NC three \NC \NR
\NC 4 \NC four \NC \NR
\NC 5 \NC five \NC \NR
\stoptabulate
```

You can do this:

```
\startcolumns[n=3]
\getbuffer
\stopcolumns
```

The (mixed) columns mechanism used here normally works ok but because of the way columns are packaged they don't work well with for instance 'vz'. Page columns do a better job but don't mix with single columns that well. Another solution is this:

```
\startrows[n=3,before=\blank,after=\blank]
\getbuffer
\stoprows
```

Here the result is collected in a vertical box, post processed and flushed line by line. We need to explicitly handle the before and after spacing here because it gets discarded (if added at all). When a slice of the box is part of the shipped out page the cells are swapped so that instead of going horizontal we go vertical. Compare the original



1 one	2 two	3 three
4 four	5 five	

with the swapped one:

1 one	3 three	5 five
2 two	4 four	

This is not really a manual but let's mention a few configuration options. The `n` parameter controls the number of columns. In order to support swapping this mechanism adds empty pseudo cells for as far as needed. By default the order is vertical but one can set it to horizontal instead. In the next example we have set height to `2\strutht` and depth to `2\strutdp`:

1 one	3 three	5 five
2 two	4 four	

When you set height and depth to max all cells will get these dimensions from the tallest cell. Compare:

1 $y = x + 1$	3 $y = \sqrt{x^2} + 1$
2 $y = x^2 + 1$	4 $y = \frac{1}{x^2} + 1$

with:

1 $y = x + 1$	3 $y = \sqrt{x^2} + 1$
2 $y = x^2 + 1$	4 $y = \frac{1}{x^2} + 1$

In the examples with `tabulate` we honor the original dimensions but you can also set the width, combined with a distance. Instead of a dimension the width parameter can be set to `fit`.

In case one wonders, of	have lines (or cells in ta-	for all problems but when
course regular columns	bles) that can be swapped.	used in situations where
can be used, but this is	For as far as possible foot-	one knows what goes in, if
an alternative that actu-	notes are supported but of	is quite powerful anyway.
ally gives you balancing	course floats are not.	It also has a relatively sim-
for free, but of course	So, this rows based mech-	ple implementation.
with the limitation that we	anism is not the solution	

### Rows becoming columns

In the previous rendering we have set the width as mentioned but also set `align to verytolerant`, `stretch` so that we don't overflow lines. The before and after parameters are set to `\blank`.

## 23.2 Colofon

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